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## PRM, The Language. Version O.2

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# PRM — The Language

Version 0.1.99

This document is a draft. The PRM language specification may evolve. The current prmc interpretation may differ from the present specification.

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# 1 A PRM Introduction

PRM stands for “Programming with Refinement and Modules”. It is an open-source programming language that has a simple straightforward style and can usually be picked up quickly, particularly by anyone who has programmed before. It is object-oriented but allows a procedural style. `prmc` is a PRM compiler that produces efficient machine language executables.

The PRM website: <http://www.lirmm.fr/~privat/prm>

PRM is a language with a high degree of structure: it is statically typed and it allows programmers to easily produce readable source code. However it has two qualities that are mainly found in dynamically typed languages:

- Concise but Clear Syntax. The syntax of the PRM language is clear and simple, without verbosity. A PRM program looks like a program written with a modern scripting language like RUBY or PYTHON—in fact, the syntax mainly comes from RUBY. The PRM syntax makes difficult to have one day an *Obfuscated PRM Code Contest*.
- Small but Powerful Core. Even if the language proposes many features, the number of different core mechanisms of the language is small. It means that the language mechanisms are clean and easy to learn. It also means that in order to provide a clear syntax with a small core, the PRM language makes a great use of syntactic sugar. Many concise pieces of syntax are strictly equivalent to a more verbose one that directly uses the core mechanism. Therefore, the semantic of these small pieces of syntax can be deduced from the semantic of the more verbose one.

## 1.1 Three Simple Examples

Before starting with the full description of the PRM syntax and semantic, here can be found three simple PRM programs. One can have an idea of the light syntax of the language.

◁ → The objective here is not to show the most valuable PRM characteristics but to give an idea of the syntax that will help to understand the full PRM specification.

### 1.1.1 Hello World

One of the simplest programs:

```
print("Hello ", "World.\n")
```

This program simply outputs the text “Hello World.”

### 1.1.2 Variables and Strings

A program with two local variables:

```
let s: String      # 's' is a local variable
s := "Hello"      # Assignment with a literal strings
s.append(" World") # Concatenate two string
let a := s + "."   # 'a' is another local variable
println(a)        # Output 'a'
```

This program also outputs the text “Hello World.”

### 1.1.3 Subprograms

A program with the definition of a function and a procedure:

```
def square(i: Int): Int
# Return  $i^2$ 
do
    return i * i
end

def main
do
    println(square(5))
    println(square(-1))
end

main
```

This program outputs the numbers 25 and 1.

## 1.2 The PRM Syntax: A First Impression

A number of points can be seen about PRM clean style and modern syntax:

- ??: comments
  - Comments are introduced by the # symbol and run to the end of line.
- ??: line structure
  - Semicolons are superfluous—but these may optionally be placed between constructs.
- ??: statement block
  - Blocks of statements start with the keyword `do` and are terminated by the keyword `end`.
- ??: procedure and function call
  - Calls use the usual notation `foo(args)`, and without arguments there is no need of parenthesis. Method invocation<sup>1</sup> on objects uses the dotted notation `x.foo(args)`.
- ??: `def`
  - Procedures and functions are defined with the `def` keyword.
- ??: `let`
  - Local variables are declared with the `let` keyword, and the static type may be inferred.
- ??: types
  - Types in signatures (`def`) and in local variables definitions (`let`) use the PASCAL column notation “`x: Foo`”.
- ??: names
  - Type names start with an uppercase (`Int`, `String`) while the other names—variables, functions, procedures, etc.—start with a lowercase (`print`, `sqr`, `s`).
- ??: assignment
  - Assignments use the `:=` notation.
- ??: literals
  - Literals strings are enclosed within double quote.

<sup>1</sup>In the current documentation, we use the terms “method invocation”. In other OO languages, one can find the same idea under the terms “method call” or “send of message”.

### 1.3 Outline

The present manual is divided into three parts:

- Section ?? is about object oriented programming: classes, properties, inheritance, etc.
- Section ?? is about modules management, class refinement, and the base library (base modules and base classes).
- Section ?? is about the base language: names, types, expressions, and statements.

## 2 Object-Oriented Programming

PRM is a pure object-oriented language. It means that:

- Each manipulated value is an object.
- Each object is an instance of a class.
- Each subprogram is a method defined in a class.
- Each method invocation corresponds to a message sending—or to a late binding i.e. depends on the dynamic type of the receiver.

→ ??: procedural style

However, it is possible to program software in PRM without explicitly defining classes and even subprograms: we call **procedural style** this way of programming. Procedural style is useful for simple programs or for teaching <sup>2</sup>. The current section focuses on “pure” object-oriented programming. The truth about procedural style is explained in a latter section.

### 2.1 Class Definition

A class represents entities, the attributes of those entities and the operations that those entities can perform. Classes can represent real world entities in a model, or more artificial artefacts which occur only in computer programs.

**class** Classes are defined with the usual **class** keyword. Here is an example of a simple PRM class definition:

```
class Car
end
```

→ ??: object creation

A class represents all objects of that type. For instance in the real world we have one concept of **Car**, but there are many instances of **Car**.

### 2.2 A Word about Class Specialisation

Classes are elements of a *specialisation* hierarchy. A class can have superclasses (i.e. class more general) and subclasses (i.e. class more specific).

→ ??: Any

The root of the hierarchy is a class named **Any**<sup>3</sup>. It means that **Any** is the most general class, and that each other class specialises the **Any** class.

<sup>2</sup> PRM is used to teach algorithmic to students in first year of computer science.

<sup>3</sup> The **Any** class corresponds to the **Object** class of the some languages like JAVA.

→ ??: class specialisation

We talk more about specialisation, and especially inheritance, in a latter section.

## 2.3 Properties: Attributes and Methods

Each class has a set of properties which represents the attributes and the operations of its instances. Operations on an object typically alter its state changing the values of one or more of its attributes. In PRM, such operations are known as *procedures*. Computations that return a value to a query about the state of an object are known as *functions*. Functions and procedures are together known as *methods*.

### 2.3.1 Definition of Properties

→ ??: names      **def** In a class, the definition of a property starts with the **def** keyword followed by the name of the property.

Attribute Definition      Attributes are identified with the first letter of their name, a @ character<sup>4</sup>.  
→ ??: type      It can be pronounced *at* and stands for *attribute*. Attributes must have a static type.

```
class Car
  def @speed: Int
    # The current speed of the car

  def @color: String
    # The color of the car
end
```

Method Definition      Methods need a signature and a body. A signature is composed of some parameters (possibly none) and a return type for functions—procedures do not have return types. Bodies are blocks of statements. A function must return its result with a **return** statement.

→ ??: statement block  
→ ??: return

```
class Foo
  def bar
    # a procedure without parameters
  do
    print("bar")
  end

  def baz: Int
    # a function without parameters
  do
    return 5
  end

  def foobar(i: Int): Int
    # a function with a parameter
  do
    return i + 1
  end

  def foobar(i: Int, j: Int): Int
```

<sup>4</sup>The @ for attributes comes from RUBY.

```

    # a function with two parameters
    do
        return i + j + 1
    end
end

```

### Some Remarks

- Since attributes start with an @ character, methods and attributes can therefore share the “same name”:

```

class Foo
    def @foo: Int
    def foo: Int
    do
        ...
    end
end

```

→ ??: accessors

This characteristic is often used for attribute accessors.

- Methods can share the same name if they have a different number of parameters. Like ERLANG or to a lesser extent, SMALLTALK or SELF.

→ ??: implicit parameter value

This characteristic is used for implicit parameter values.

- There is no static overriding: two methods with the same name and the same number of parameters cannot coexist in a same class, even if they have distinct signatures.

→ ??: covariant typing policy

This is because PRM follows a covariant typing policy.

→ ??: procedural style

- Methods can be defined outside class with the same syntax. Obviously, it is not possible with attributes.

### 2.3.2 Access to attributes

Objects directly access their attributes by their names:

```

class Foo
    def @bar: Int
    # an attribute

    def baz
    do
        @bar := 5 # write access
        print(@bar) # read access
        # output '5'
    end
end

```

→ ??: assignment

Attribute write accesses follow the assignment rules.

→ ??: exported attribute

Except in some special cases, attributes can only be accessed by the objects that own them.

### 2.3.3 Invocation of Methods

Methods are usually invoked with the common dotted notation `x.foo(args)` where `x` is the *receiver*, `foo` the name of a method and `args` some arguments. If there is no arguments, parentheses are optional—providing superfluous ones may provoke a warning during compilation.

```
class Foo
  def bar
  do
    print("bar")
  end

  def baz: Int
  do
    return 5
  end

  def foobar(a: Int)
  do
    print("foobar", a)
  end

  def foobaz(a: Int): Int
  do
    return a + 1
  end
end
...
let x: Foo
...
x.bar           # Output 'bar'
print(x.baz)    # Output '5'
x.foobar(6)     # Output 'foobar6'
print(x.foobar(7)) # Output '8'
```

`self`, the Current Receiver If the receiver is the current receiver, called `self` in PRM<sup>5</sup>, it can be implicit. Therefore, `self.foo(args)` is equivalent with `foo(args)`.

```
class Foo
  def bar
  do
    print("bar")
  end

  def baz
  do
    bar
    self.bar
  end
end
...
let x: Foo
...
x.bar # Output 'bar'
```

<sup>5</sup>In C++ and JAVA, the current receiver is called `this` ; in Eiffel it called `Current`.

```
x.baz # Output 'barbar'
```

→ ??: visibility

It is important to distinct invocation on `self` with other invocations since visibility does not apply on invocation on `self`.

Some methods that have a special name are not invoked with the usual dotted syntax. These methods are *operators* and *assignment procedures*.

## 2.3.4 Operators

Operators are methods often used for mathematical operations. There are three kinds of operators: infix operators, prefix operators and bracket operators.

**Infix Operator** Infix operators are: `+`, `-`, `*`, `/`, `%`, `=`, `!=`, `<`, `>`, `<=`, `>=`, `<=>`, `<<`, and `>>`. They are methods with one parameter and should be defined and invoked as follow:

```
# Definition signature
class C
  def -(o: T): U
  ...
end

# Invocation: x is the receiver and y the argument
let x: C
let y: T
let r: U
r := x - y
```

**Prefix Operator** Prefix operators are: `+` and `-`. They are methods without parameters and should be defined and invoked as follow:

```
# Definition signature
class C
  def -: U
  ...
end

# Invocation: x is the receiver
let x: C
let r: U
r := -x
```

**Bracket Operator** Bracket operators are methods mainly used for indexed access (arrays for instance). They are defined and invoked as follow:

```
# Definition signature
class C
  def [](o: T, p: U): V
  # Example with two parameters
  ...
end

# Invocation:
# * x is the receiver, y and z are the arguments
let c: C
```

```

let y: T
let z: U
let r: V
r := x[y, z]

```

↳ → Two remarks about operators:

→ ??: Booleans

- Invocation of operators requires an explicit receiver.
- Some statements looks like operations but are not. For instance the Boolean's pseudo-operators: `and`, `or`, `not`.

### 2.3.5 Assignment procedures

→ ??: assignment

Their names are ended with `:=`. They must have at least one parameter and no return value. They are invoked with a syntax that looks like the assignment statement and follow its rules.

```

# Definition signature
class C
  def foo:=(o: T)
    # Example of the simple form, i.e. with one parameter
    ...

  def foo:=(o: T, p: U, q: V)
    # Example with three parameters
    ...
end

# Invocation
let x: C
let y: T
let z: U
let k: V

# * x is the receiver, y is the argument
x.foo := y

# * x is the receiver, y is the first argument,
# z the second one, and k the third one
x.foo(y, z) := k

```

Bracket Assignment  
Procedure

There are also bracket assignment procedures:

```

# Definition signature
# * Example with three parameters
class C
  def []:=(o: T, p: U, q: V)
    ...
end

# Invocation
# * x is the receiver, y is the first argument,
# z the second one, and k the third one
let x: C
let y: T
let z: U

```

```
let k: V
x[y, z] := k
```

◁ → Some remarks about assignment procedures:

→ ??: accessors

- They also exist in RUBY in the simplest form (i.e. with exactly one parameter).
- With one parameter, they are mainly used to write attributes accessors.
- With more than one parameter, they are mainly used with indexed access when different kinds of indexes exist.
- Assignment procedure is different with the user-defined C++ assignment operator. In C++, “x.a = 5” may correspond to the invocation of the `operator=` method on the attribute `a` of the `x` object. In PRM “x.a := 5” corresponds to the invocation of the `a:=` method on the `x` object.

### 2.3.6 Implicit Parameter Value

◁ → The current prmc compiler does not yet implement this part of the specification.

PRM can yield a kind of implicit argument value:

```
class C
  def foo(a: Int, b: Int := 5, c: Int)
  do
    ...
  end
  ...
end
```

However, implicit argument is only syntactic sugar since the last code example is strictly equivalent to:

```
class C
  def foo(a: Int, b: Int, c: Int)
  do
    ...
  end

  def foo(a: Int, c: Int)
  do
    foo(a, 5, c)
  end
end
```

Multiple Implicit Parameter Values

A method can have multiple implicit parameter values:

```
class C
  def foo(a: Int := 5, b: Int := 6)
  do
    ...
  end
```

```
end
```

In order to avoid ambiguities, the first parameters are less implicit than the last parameters. Therefore the two following listings are equivalent:

```
class C
  def foo(a: Int, b: Int)
  do
    ...
  end

  def foo(a: Int)
  do
    foo(a, 6)
  end

  def foo
  do
    foo(5, 6)
  end
end
```

Comparison with Other Languages

Implicit parameter values exist in many other languages like C++, JAVA 5.0 or RUBY. However, their semantics differ in the way that there is one method defined, and the implicit parameter values are integrated to the arguments when the method is invoked.

For instance, let “void foo(int a, int b = 5)” be a C++ method. The two expressions `foo(1, 5)` and `foo(1)` invoke this method with the same arguments, since 5 is implicitly added in the last expression.

With PRM and the equivalent `foo` method “def foo(a: Int, b: Int := 5)”, the two expressions `foo(1, 5)` and `foo(1)` invoke two distinct methods—respectively, `foo` with two parameters and `foo` with one parameter.

The PRM way offers two advantages:

- Implicit parameter values are only sugar syntax: it does not extend the PRM core mechanism.
- Each method is independent and can be independently redefined. Thus more flexibility is offered to the programmer.

→ ??: redefinition

### 2.3.7 Variable Argument Number

◁ → This part of the specification is not considered as stable and may change in future version.

In PRM, some methods can be invoked with an unbounded number of arguments.

Definition In those method definitions, one special parameter is declared as `t: T*` where `t` is the name of the parameter and `T` the type of arguments. In such definitions, the static type of `t` is `Array[T]`. Example:

→ ??: Array

```
class C
  def foo(a: Int, b: Int*, c: Int)
```

```

    # The static type of 'b' is 'Array[Int]'
    do
        print(a, "-", b.length, "-", c)
    end
end

```

Invocation In method invocation, the special parameter is associated with one or more arguments:

```

let c: C
c.foo(1, 2, 3) # 'Output 1-1-3'
c.foo(1, 2, 3, 4, 5, 6) # 'Output 1-4-6'
c.foo(1, 2) # compilation error, unknown foo method
# with two parameters

```

→ ??: print

Without surprise, the standard print method, used in all those examples, accepts multiple arguments. Its signature is `print(a: Any*)`.

Passing Array One can pass an array object instead of a list of elements with the `*a` notation— here, `*` is not an operator, it is just a notation:

```

let a: Array[Int]
let c: C
...
c.foo(1, *a, 3)

```

Passing array is often used to chain calls:

```

class C
  def printprint(a: Any*)
    # Remember, the static type of 'a' is 'Array[Any]'
    do
        print(*a)
        print(*a)
    end
end
let c: C
c.printprint("Hello") # Output 'HelloHello'

```

⊲ → In a class, cannot coexist:

- Two methods with the same name that both accept a variable number of arguments:

```

class Foo # ERROR !
  def bar(a: Int, b: Int*, c: Int)
    do ... end

  def bar(d: Int*)
    do ... end
end
let f: Foo
f.bar(1, 2, 3) # Ambiguous

```

- Two methods with the same name, one accepts a variable number of arguments, and the other has more parameters than the minimal number of the first:

```

class Foo
  def bar(a: Int, b: Int*)
    do ... end

  def bar(c: Int, d: Int)
    do ... end
end
let f: Foo
f.bar(1, 2) # Ambiguous

```

However, the following listing is OK:

```

class Foo # ERROR !
  def bar(a: Int, b: Int*)
    # bar1
    do ... end

  def bar(c: Int)
    # bar2
    do ... end
end
let f: Foo
f.bar           # Not ambiguous, it is an error
f.bar(1)       # Not ambiguous, it is bar2
f.bar(1, 2)    # Not ambiguous, it is bar1
f.bar(1, 2, 3) # Not ambiguous, it is bar1

```

## 2.4 Object Creation

**new** Objects are created with the special **new** statement:

```

new Car
new Car("red")
new Car.with_color("blue")

```

The point to note is *constructors* need to be declared in classes in order to allow them to be instantiated. In PRM, constructors are a little different from those of languages like C++ and JAVA; EIFFEL constructors are the closest.

**Constructor** Constructors are procedures defined in a class after the **constructor** keyword. More than one procedure can be defined as constructors.

In the following listing, the two **init** procedures and the **with\_color** one are constructors, but **paint** is a “normal” procedure—note that if a lot of code is duplicated it is only for the need of the illustration:

```

class Car
  def @color: String
    # The color of the car

  def paint(c: String)
    # Repaint the car
  do
    @color := col

```

```

        end
    constructor
        def init
        do
            @color := "black" # Mr. Ford?
        end

        def init(col: String)
        do
            @color := col
        end

        def with_color(col: String)
        do
            @color := col
        end
    end
end

```

◁ → The PRM naming convention is to reserve the methods named `init` or `with_something` to be constructor procedures.

→ ??: Visibility

It is important to notice that constructors look like “normal” procedures. The visibility section will show the truth about the `constructor` keyword and status of constructors.

Implicit Constructor

Implicitly, the constructor named `init` is called on object instantiation. Thus `new Car` is equivalent to `new Car.init`, and `new Car("Blue")` is equivalent to `new Car.init("Blue")`

◁ → Since both object creation and method invocation use a dot in their notation, some cases should be disambiguated:

- `new Foo.bar` is always considered as the instantiation of a `Foo` object with a constructor named `bar`.
- `(new Foo).bar` and `new Foo.init.bar` are the instantiation of a `Foo` object with a constructor named `init`; and followed by the invocation of a property named `bar` on this newly created object.
- `new Foo(5).bar` and `new Foo.init(5).bar` are the instantiation of a `Foo` object with a constructor named `init` with 5 as argument; and followed by the invocation of a property named `bar` on this newly created object.

#### 2.4.1 Abstract Classes

*Abstract classes* are classes that can not be instantiated. Classes that are not abstract are called *concrete classes*.

◁ → In PRM, Abstract classes are simply classes without constructors. In corollary, classes without constructor are abstract, therefore can not be instantiated.

Empty Constructor

In comparison with other languages, there are no default constructors since their use is marginal, even if they are the cause of many errors.

However, sometimes, programmers need to define concrete classes with empty constructors. In PRM they just have to explicitly do it:

```
class Foo
  constructor
    def init
    do
    end
  end
end
```

## 2.4.2 Garbage Collector

PRM has no delete operator. This is because, as many other modern languages, PRM is garbage collected. Garbage collection is known to completely cure the programming ills of dangling pointers and memory leaks. This greatly simplifies the programming effort by removing one of the largest bookkeeping headaches for programmers. Garbage collection has also proved to be very efficient in modern implementations.

## 2.5 Visibility

C++ and JAVA programmers might be wondering how to make methods public, protected and private. With PRM you have far more control: as in Eiffel, any set of methods can be exported to all, to none or to some specific classes. Thus you have the possibility of many shades of grey between public and private. You might want a method to be public to some specific classes, but private to others. Moreover, method visibility and constructors are related together in a nice original way.

◁ → Visibility is not related to method invocation on `self`. Therefore, properties are always accessible to the current receiver.

### 2.5.1 Method Visibility Blocks

Visibility is controlled by three keywords that delimit *visibility blocks*: `public`, `private`, and the already known `constructor`. Method defined after such a keyword belong to the corresponding visibility block.

A class definition can contain any number of blocks, in any order:

```
class Foo
  ...
  public
  ...
  public
  ...
  private
  ...
  public
  ...
  constructor
  ...
  private
  ...
end
```

◁ → It is recommended to regroup related methods with the same visibility in the same block. And it is also recommended to put two unrelated sets of methods in two different visibility blocks, even if they share the same visibility.

## 2.5.2 Public Method Visibility

Methods defined in a `public` block are exported and can be used by other classes. If the name of a class is added after the `public` keyword, methods are only exported to this class and to its subclasses. If there is not such a class name, methods are exported to any classes—in fact, they are exported to the `Any` class and to its subclasses.

→ ??: `Any`

Example:

```
class Car
...
public
  def speed: Int
  # Get the speed of the car
  do
    return @speed
  end
public Driver
  def stop
  # Stop the car
  do
    @speed := 0
  end
...
end
```

Let `c` be a variable statically typed by a `Car`. Here the function `speed` is exported to any class, therefore `c.speed` is valid in any class. The procedure `stop` is exported to the class `Driver` (and all its subclasses), therefore `c.stop` is only valid in the class `Driver` and in any subclasses of `Driver`.

Implicit Visibility Block The *implicit visibility block* (i.e., the visibility block above the first visibility keyword) is a `public` one. For example, the three following listings are equivalent:

```
class Foo
public Any
  def bar
  do
    print("baz")
  end
end
```

```
class Foo
public
  def bar
  do
    print("baz")
  end
```

```

end

class Foo
  def bar
  do
    print("baz")
  end
end

```

### 2.5.3 Private Method Visibility

Method defined in a `private` block are not exported. Therefore, private methods are only accessible to the current receiver.

```

class Driver
  def @car: Car
    # The driven car

  private
    def stop_car
      # Stop the driven car
    do
      @car.stop
    end
  end
end

```

Let `c` be a variable statically typed by a `Car`. The procedure `stop_car` is exported to nobody, therefore `c.stop_car` is valid nowhere. The only way to invoke such a method is to use the current receiver.

PRM private vs. C++ private  
 In PRM, private methods are usable only by self—it is an instance visibility.  
 In C++, private methods are usable only by instances of the current classes—it is a class visibility.

The following listing will try to illustrate the difference:

```

class Foo
public Foo
  def bar
  ...
private
  def baz
  ...
...
public
  def test
  do
    bar # OK, the receiver is self
    baz # OK, the receiver is self

    let f: Foo

    f.bar # OK, bar is public Foo
    # and I am Foo

    f.baz # Error, baz is private
  end
end

```



```

        def paint(col: String)
        do
            @color := col
        end
    constructor
        def with_color(col: String)
        do
            paint(@col)
        end
    end
end

```

### 2.5.5 Attribute Accessor

As in SMALLTALK, attributes are “private”: they can only be accessed by the objects that own them. Therefore, some methods should be defined in order to access attributes. Methods that play this role are called *accessors*.

Usually, there is the need of two accessors, one for the read access and one for the write access. In PRM, you can use the same name for the attribute and for the two accessors: the attribute is distinguished with the @ and the write accessor is usually an assignment procedure, therefore distinguished with the :=.

→ ??: assignment procedure

Example:

```

class Car
    def @speed: Int
        # Attribute

        def speed: Int
            # Read accessor
        do
            return @speed
        end

        def speed:=(s: Int)
            # Write accessor
        do
            @speed := s
        end
    end
    ...
end

```

In this example, let *c* be a *Car*. *c.speed* returns the value of the attribute @speed and *c.speed := 5* assigns 5 to the attribute @speed:

```

let c := new Car
c.speed := 5
print(c.speed) # Output '5'
c.speed := 10
print(c.speed) # Output '10'

```

Automatic Accessor The keywords `def_read` and `def_write` can be used to simplify the declaration of such accessors. On attribute definition, `def_read`

automatically generates a read accessor and `def_write` automatically generate a write accessor.

The following example is equivalent to the previous one.

```
class Car
  def @speed: Int def_read def_write
  ...
end
```

#### Automatic Accessor Visibility

Since `def_read` and `def_write` only correspond to syntax sugar, the visibility of automatic accessors is the one of the current visibility block.

```
class Car
public
  def public_price: Int
  do
    return @cost + @margin
  end

public CarSeller
  def @cost: Int def_read
  def @margin: Int def_read def_write
end
```

In this example, only a car seller can access the real price of a car.

#### Pseudo-accessor

Accessors are just a role playing by some methods. It is possible to define “pseudo-accessors”, i.e. methods that act like accessors from the user point of view. The following example defines two pairs of accessors on the speed attribute of a `Car` class but with different speed units, one in kilometre per hour and the other in miles per hour:

```
Class Car
  def @speed_kmph: Int def_read def_write
  # Speed in kmph

  def speed_mph: Int
  # Speed in mph
  do
    return @speed_kmph * 63 / 100
  end

  def speed_mph := (s: Int)
  # Speed in mph
  do
    @speed_kmph := s * 100 / 63
  end

  End
  ...
end
```

Thus, from a user point of view, it is not possible to distinguish the “true” accessor from the pseudo-accessor:

```
let c := new Car
c.speed_kmph := 80
print(c.speed_kmph, " ", s.speed_mph)
```

```
# Output '80 50'
c.speed_mph := 63
print(c.speed_kmph, " ", s.speed_mph)
# Output '100 63'
```

## 2.5.6 Exported Attribute

◁ → The specification is not considered as stable and may change. Moreover, the current `prmc` compiler do not yet implement it.

→ ??: attribute access Software engineering considers that attribute should be accessible only for the current receiver (`self`). However, in some exceptional case, attributes need to be directly accessed by different objects.

`export` The keyword `export` permits to change the visibility of an attribute. The visibility granted is the one of the current visibility block.

In the following example, the attribute `@baz` is visible in the class `Bar` and in all its subclasses:

```
class Foo
public Bar
    def @baz: Int
    export @baz
end
```

Exported Attribute Access Exported attributes are accessed with the dotted notation `x.@baz` where `x` is the receiver (i.e. the instance that owns the attribute) and `@baz` the name of the attribute.

→ ??: assignment Exported attributes can be used as an expression or as the left part of an assignment:

```
f.@baz := f.@baz + 1
```

## 2.6 Class Specialisation

Specialisation has three main uses:

- ??: property inheritance
  - Build new classes out of existing classes since classes inherit properties defined their superclasses.
- ??: visibility
  - Gain property visibility since properties exported to a class (`public` and `constructor`) are visible to their subclasses.
- ??: type
  - Permit subtyping since objects of a class can be used where objects of the superclasses are expected.

◁ → In many object-oriented languages, inheritance is mainly a way to reuse property already defined. The semantic of inheritance of the PRM language is a bit different since it strictly corresponds to the natural semantic of specialisation: *If A is a superclass of B then each instance of B is also an instance of A.* The three uses of specialisation are simply corollaries of this strict semantic. It also means that two uses of specialisation, frequent in some OO languages, are forbidden in PRM: inheritance of implementation and repeated inheritance.

**inherit** The `inherit` keyword is used to declare the superclass of the class. This keyword must be used before any property declarations.

The following listing is a very simple example of inheritance where a `Car` class is a subclass of a `Vehicle` class:

```
class Vehicle
end

class Car
inherit Vehicle
end
```

Multiple Class Specialisation With multiple specialisation, the `inherit` keyword is repeated:

```
class Drake
inherit Duck
inherit Male
end
```

Transitive Specialisation In PRM, transitive specialisation relation links are ignored. Therefore, the two following listings are equivalent:

```
class Ambulance
inherit Car
end
```

```
class Ambulance
inherit Car
inherit Vehicle
end
```

Moreover, the last one may produce a warning during compilation because of the superfluous `inherit Vehicle`.

### 2.6.1 Property Inheritance and Redefinition

Properties Inheritance Subclasses inherit the properties—attributes and methods—of their superclasses.

```
class Car
  def @color: String def_read

  def sound: String
  do
    return "vroom"
  end
end

class Convertible
inherit Car
  def @roof_is_open: Boolean
end
```

This example shows a superclass `Car` and a subclass `Convertible`. The `Convertible` class inherits the following properties: the attribute `@color`,

→ ??: accessors

the automatic `color` accessor, and the `sound` function. It also defines a new property, the attribute `@roof_is_open`.

Properties Redefinition

Subclasses can redefine some inherited properties by providing a new definition of a property.

The following example shows the redefinition of the `sound` function inherited from the `Car` class:

```
class Ambulance
inherit Car
  def sound: String
  do
    return "wo-wo"
  end
end
```

Precursor

With the Eiffel terminology, we says that the `sound` method of the `Car` class is a *precursor* of the `sound` method of the `Ambulance` class.

Global Property

In the previous example, the `sound` method of the `Car` class and the `sound` of the `Ambulance` class are two different methods. However, they belong to a “same property idea”, here the idea is something like “sound of cars”. We call *global property* this “same property idea”.

⊲ →

Global properties are introduced when its first property is defined. Example, the global property “sound of cars” is introduced in the `Car` class by the `sound` method.

In PRM, global properties are not strictly related to properties names. For example, in the following listing, the two properties `@height` belong to distinct global properties:

```
class Person
  def @height: Int # in cm
  def @weight: Int # in kg
end

class Button
# A button for a graphical user interface
  def @height: Int # in pixels
  def @width: Int # in pixels
end
```

⊲ →

This notion of *global property* is one of the PRM exclusivity. In great majority of other OO languages, the absence of this notion yields quantities of problems.

Attribute Redefinition

Obviously, redefinition is majority used for methods. It is also possible to redefine an attribute by specialising its static type:

```
class Car
  def @driver: Person
end

class PoliceCar
  def @driver: Policeman
end
```

→ ??: covariant typing policy This is because has a covariant typing policy.

### 2.6.2 Deferred Method

*Deferred methods* (called *pure virtual method* in C++) are methods without implementation. A deferred method is declared without a body, instead it has the `as deferred` keywords.

```
class Car
  def has_priority: Bool as deferred
end

class Ambulance
inherit Car
  def has_priority: Bool
  do
    return false
  end
end
```

Remarks:

→ ??: abstract classes

- Usually, classes that contain deferred methods are mainly abstract classes—i.e. do not have constructors.

→ ??: refinement

- Concrete classes with deferred methods can be useful with refinement.

### 2.6.3 Multiple Inheritance

When a class has only one super-class, inheritance and redefinition are quite intuitive mechanisms. PRM multiple inheritance mechanism is also intuitive.

Which Properties to Inherit? The inherited properties are the most specific ones—i.e. the properties **defined** in the most specific classes. This base behaviour is quite simple but slightly differs from the majority of OO languages.

Example:

```
class A
  def foo
  do
    print("fooA")
  end
end

class B
inherit A
end

class C
inherit A
  def foo
  do
    print("fooC")
  end
end
```

```

class BC
inherit B
inherit C
end

```

In the BC class, there are two potential inherited methods: `foo` defined in the A class, and `fooC` defined in the C class; the second is the most specific because C specialise A; therefore the BC class inherit the "foo" method defined in C.

Multiple Precursors A property can redefine more than one property inherited from super-classes:

```

class D
inherit A
  def foo
  do
    print("fooD")
  end
end

class CD
inherit C
inherit D
  def foo
  do
    print("fooCD2")
  end
end

```

Here, the `foo` method of the CD class has two precursors since it redefines the `foo` methods of the classes C and D.

Property Conflict

A property conflict occurs when the most specific property to inherit is not unique:

```

class CD2
inherit C
inherit D
end

```

The solution to avoid them is to redefine the conflicting property.

→ ??: deferred method

When all properties but one are deferred, the conflict is automatically resolved: the one that is not deferred is inherited.

Global Property Conflict

A global property conflict occurs when a class inherits homonym properties that belong to distinct global properties:

```

class O
  def foo
  do
    print("fooO")
  end
end

```

```

end

class A0
inherit A
inherit O
end

```

**rename** The solution to avoid them is to rename at least one of the two conflicting method with the **rename** keyword:

```

class A02
inherit A rename foo(0) as fooA
inherit O
constructor
    def init do end
end
let x := new A02
x.foo    # Output 'fooO'
x.fooA  # Output 'fooA'

```

↳ → When renaming methods, the first name has to precise between parentheses the number of parameters.

One can rename more than one property

```

class Y
    inherit X rename foo(0) as fooX,
                @bar as @barX, -(1) as minus
end

```

The PRM renaming differs from the Eiffel one and is slightly simpler and more coherent:

- One renaming per global property is enough, even if the global property comes from many super-classes:

```

class AC2 # WARNING: Superfluous renaming.
inherit A rename foo as foo2
inherit C rename foo as foo2
end

```

- A global property can not have two names in a same class:

```

class AC2 # ERROR: Multiple renaming.
inherit A rename foo as fooA
inherit C rename foo as fooC
end

```

- Two distinct global properties cannot be renamed to have the same name:

```

class P
    def bar
    do
        print("barP")
    end
end

```

```

end
class AP # ERROR: Global property conflict.
inherit A
inherit P rename bar as foo
end

```

## 2.6.4 Visibility

**Visibility Inheritance** The visibility `public` and `private` of method inherited. However, constructors are inherited as private methods. This is because constructors of a class are not adapted to its subclasses.

In the following listing, `Ambulance`, a subclass of a `Car` class inherit the `with_color` as a private method. Therefore, `Ambulance` defines `init`, a new specific constructor.

```

class Car
  def @color: String def_read

  constructor
    def with_color(col: String)
    do
      @color := col
    end
  end
end

class Ambulance
inherit Car
constructor
  def init
  do
    with_color("white")
    # OK, since with_color is inherited
  end
end
end

```

Here some uses of the two classes:

```

let c := new Car.with_color("black")
print(c.color) # Output 'black'
let a1 := new Ambulance
print(a1.color) # Output 'white'
let a2 := new Ambulance.with_color("blue") # Error!
# -> 'with_color' is not a constructor,
# it is a private method

```

**Visibility Redefinition** One can redefine the visibility of inherited method by redefining the method in the wanted visibility block. The visibility of inherited method can also be redefined without having to redefine the whole method.

**export** The `export` keyword enables to change the visibility of inherited method to the one of the current block:

```

class Car
constructor
  def init

```

```

        do
            @speed := 0
        end
    end

class Ambulance
inherit Car
constructor
    export init(0)
end

let c := new Ambulance # OK

```

↔ → As for renaming, the number of parameters has to be indicated between parentheses.

Multiple methods can be exported at the same time:

```

class Bar
inherit Foo
public Baz
    export foo(1), bar(0), baz:=(1), +(1)
end

```

## 2.6.5 Call to Super

**super** In a method redefinition, the programmer can refer to the previous property with the `super` keyword:

```

class Foo
    def foo(i: Int): Int
    do
        return i + 1
    end
end

class Bar
inherit Foo
    def foo(i: Int): Int
    do
        return super(i*2) * 2
    end
constructor
    def init do end
end

let b := new Bar
print(b.foo(2)) # Output '10'

```

**Implicit Super Arguments** Arguments of a super call are implicitly the parameters of the method. Therefore, the two following listings are equivalent:

```

...
    def foo(a: Int, b: String)
    do
        ...
    end

```

```

        super(a, b)
        ...
    end
...

```

```

...
    def foo(a: Int, b: String)
    do
        ...
        super
        ...
    end
...

```

Multiple Precursor When a method has more than one precursor, any call to `super` must be prefixed with a class name in order to remove the ambiguity. Such prefixes use the `::` notation:

```

class CD3
inherit C
inherit D
    def foo(a: Int)
        # This method redefines the ones
        # of the classes C and D
    do
        C::super(a+2)
        D::super(a-1)
    end
end

```

## 2.7 Genericity

Inheritance is one of the fundamental mechanisms for reuse; so is genericity. Genericity is also important in making programs type safe without resorting to type casts. Java 5.0 introduces genericity, in previous version, many type casts were needed to make up for this deficiency. C++ has genericity in the form of template classes. If you have had problems understanding C++ templates, don't worry, PRM's generic syntax is much easier, and more powerful, as it also allows generic parameters to be bounded; this is known as bounded genericity.

In PRM, genericity is mainly the one of the Eiffel language, please refers its specification to know more about genericity.

Generic Class Definition In order to use genericity, you create a *generic class* with formal generic parameters. In the following listing, `Pair` is a generic class with one formal parameter bounded by `Any` and `VehiclePark` is a generic class with one formal parameter bounded by `Vehicle`.

```

class Pair[E: Any]
end
class Vehicle
end
class VehiclePark[E: Vehicle]
end

```

```
class Car
inherit Vehicle
end
```

Generic Type Construction In program, the generic class can be used to construct many kinds of *generic types*:

```
let x: Pair[Int]           # x is a pair of integers
let y: Pair[Pair[String]] # y is a pair of pairs of strings
let z: VehiclePark[Car]   # z is a car-park
let t: VehiclePark[Int]   # Error since integers are not vehicles
```

Formal Generic Parameter Use Inside the class definition, the formal generic parameter can be used as a type:

```
class Pair[E: Any]
  def @first: E def_read def_write
  def @second: E def_read def_write

  def switch
  do
    let t: E
    t := @first
    @first := @second
    @second := t
  end

  def display
  do
    print(@first, " ", @second)
  end

  constructor
  def init(f: E, s: E)
  do
    @first := f
    @second := s
  end
end
```

Here some examples of use:

```
let pi := new Pair[Int].init(5, 4)
pi.display # Output '5 4'
pi.switch
pi.display # Output '4 5'

let ps := new Pair[String]("Hello", "Town")
ps.second := "World"
ps.display # Output 'Hello World'
```

Generic Types and Subtypes Genericity yields a kind of subtyping. For example, `VehiclePark[Car]` is a subtype of `VehiclePark[Vehicle]`:

```
let vehiclepark: VehiclePark[Vehicle]
let carpark: VehiclePark[Car]
```

```
vehicle := carpark # OK
```

→ ??: covariant typing policy

This is because PRM follows a covariant typing policy.

## 3 Modules

In the PRM language, classes are organised into *modules*. A module corresponds to a source file and contain class definitions. They are also the compilation units: each module can be separately compiled then linked to produce an executable.

Filename PRM source files should follow a strict naming scheme. They must be named `foo.prm` where `foo` is the name of the module.

### 3.1 Module Structure

A PRM source file (a module) is divided into four parts:

→ ??: module dependence

1. Importation of modules.

→ ??: class definition

2. Definition of classes.

→ ??: outside method

3. Definition outside classes of procedures and functions.

→ ??: main statements

4. Definition of the module main statements.

Each part is optional but the order must be respected.

### 3.2 Module Dependence

Modules can *depend* on others modules and *import* their classes. With analogy with the class terminology, we call *supermodules* of a module `m` the modules it depends on, and *submodules*<sup>6</sup> the modules that depend on `m`.

`import` The dependence between modules is declared with the `import` keyword followed by the name of the module:

```
import crypt
import http
```

⊲ → As in class hierarchy, cycles are forbidden in the module dependency! A module `m` cannot require itself nor require a module that requires `m`.

Implicit Dependence  
→ ??: `standard`

Implicitly, modules depend on the module `standard` that contains standard classes.

Class Conflict

There is a class conflict when a module imports two homonymous distinct classes.

Such a conflict can be resolved with the `rename` keyword:

```
import automobile # import Car
import tramway rename Car as Tramcar
```

→ ??: property renaming

Class renaming follows and property renaming follows the same rules.

<sup>6</sup> In some language, “submodules” refers to nested modules (i.e. modules defined into a modules). Since there is no module nesting in PRM, there are no ambiguities.

### 3.3 Class Refinement

PRM modules can extend imported classes, this is called *class refinement*:

```
# File m1.prm
class Foo
  def bar
  do
    print("before")
  end
constructor
  def init do end
end
```

```
# File m2.prm
import m1
class Foo
  def bar
  do
    print("after")
  end
end

(new Foo).bar # Output 'after'
```

Properties: Definition and Redefinition    The main usage of refinement is to add new properties (methods and attributes) or to redefine them.

◀ →    It's important to note that refinement is not specialisation: if you specialise a class, you have two classes, if you refine a class, you still have one class.

Class refinement is one of the greatest PRM features. It improves the reusability of OO software since it provides an answer to the separation of concern problem: a module can adapt existing classes to new concerns. Class refinement is clearly not a new OO feature and exists in many dynamically typed languages (like RUBY or LISP) and in some statically typed language (like OBJECTIVE-C).

Multiple Refinement    Refinement can be also combined without difficulties:

```
# File m3.prm
import m1
class Baz
  def @foo: Foo

  def bar
  do
    @foo.bar
  end
constructor
  def init
  do
    @foo := new Foo
  end
end
```

```
# File m4.prm
import m2
import m3
(new Baz).bar # Output 'after'
```

Class refinement works for any classes; even with build-in ones:

```
class Int
  def fib: Int
  # Fibonacci numbers
  # The inefficient recursive algorithm
  do
    if self <= 0 then
      return 0
    elsif self <= 2 then
      return 1
    else
      return (self - 1).fib +
             (self - 2).fib
    end
  end
end

print(6.fib) # Output '8'
```

Addition of Superclasses

It is also possible to refine a class by giving it new superclasses. This kind of refinement is quite rare, even in dynamically typed languages:

```
class Foo
  def foo
  do
    print(self, "-foo")
  end
end

class String
inherit Foo
end

"Hello".foo # OK since String inherit the foo method
# Output 'Hello-foo'

let f: Foo
f := "World" # OK since Strings are now Foos
f.foo # Output 'World-foo'
```

Refinement, Specialisation  
and Multiple Inheritance

### 3.4 Procedural style

For property inheritance, refinement behaves like specialisation.

From a programmer point of view, PRM can be used as a procedural language:

- Some methods, like `print` seems to not have receiver.
- Procedures and functions can be defined outside classes.

- The main statements can be written outside procedures and functions.

However, PRM is a pure object-oriented language: procedures and functions are always methods, and statements always belong to method bodies.

### 3.4.1 Method Without Receiver

→ `?: self`

Since each method invocation has a receiver, it means that each `print("Hello World")` use `self` as receiver.

In fact, `print` is a private method defined in `Any`, thus inherited as a private method in any other classes:

```
print("Hello ")      # Output 'Hello '
self.print("World") # Output 'World'
5.print(".")        # Error
# -> 'print' is a private method
```

### 3.4.2 Method Definition outside Classes

Procedures and functions defined outside classes are implicitly defined as private methods of the class `Any`. Therefore, as `print`, they can be used everywhere.

◁ → Methods defined outside classes correspond to an implicit refinement of the class `Any`. Example: the two following listings are equivalent:

```
def foo
do
    print("hello world")
end
```

```
class Any
private
    def foo
    do
        print("hello world")
    end
end
```

### 3.4.3 Module Main Statement

The module main statements belong to the body of an implicit private `main` procedure defined in an implicit `Sys` class.

→ `?: Sys class`

The `main` method of the `Sys` class corresponds to the entry point of programs.

◁ → Main statements correspond to an implicit refinement of the class `Sys`. The two following listings are equivalent:

```
print("Hello world")
```

is syntactically equivalent to

```

class Sys
private
    def main
    do
        print("Hello world")
    end
end

```

### 3.5 Base Modules

The PRM standard library is made of 6 modules:

→ ??: base classes

**standard** The standard module is the implicitly required by other modules. It contains all the necessary base classes.

**net** implicitly depends on **standard**.

It specialises some IO classes in order to provide networking socket communication.

**http** depends on **net**.

It defines classes related to basic HTTP communication.

**exec** implicitly depends on **standard**.

It defines classes used to execute arbitrary commands of the shell system.

**sdl** implicitly depends on **standard**.

It defines wrapper classes around the Simple DirectMedia Layer C library<sup>7</sup>. It is primarily defined to experiment the feasibility of wrapping C libraries.

**opts** implicitly depends on **standard**.

It defines classes related to the parse and the analysis of command line options.

◀ → This base module is the only one developed by someone else: Floréal Morandat.

#### 3.5.1 Standard Module

In fact, the **standard** module is not the root of the module hierarchy. **standard** is an empty module that only require base classes from some “super-standard” modules.

◀ → This section is mainly an illustration to how module hierarchy and class refinement can be used to develop modular applications—i.e. where modules are clearly “concern” units.

The PRM **standard** module depends on 11 super-modules:

<sup>7</sup> <http://www.libsdl.org/>

kernel	depends on nothing.  It is the root of the module hierarchy and minimally defines the most basic classes like <code>Any</code> , <code>Sys</code> , <code>Int</code> , <code>Float</code> , <code>Bool</code> or <code>Char</code> .
math	depends on <code>kernel</code> .  It refines the <code>Int</code> and <code>Float</code> classes with useful mathematical methods; for instance trigonometry. In a near future it will also define some class like complexes or big numbers.
abstract_collection	depends on <code>kernel</code> .  It defines the PRM collection abstract generic class hierarchy. From the general <code>Collection[E: Any]</code> and <code>Iterator[E: Any]</code> classes to more specific like abstract set or abstract maps (i.e. associative arrays).
range	depends on <code>abstract_collection</code> .  It defines the <code>Range</code> class and related ones.
array	depends on <code>abstract_collection</code> .  It defines the <code>Array</code> class and related ones. It also specialises many abstract collection classes into concrete ones implemented with arrays— <code>ArrayMap</code> , <code>ArraySet</code> , etc.
sorter	depends on <code>array</code> .  It refines the classes of the array module by the addition of sort method.
list	depends on <code>abstract_collection</code> .  It specialises some collection classes—i.e. like the <code>array</code> module but with linked lists.
string	depends on <code>array</code> .  It defines the <code>String</code> class and related ones. It also refines many classes by adding a <code>to_s</code> method used to convert any objects to a human readable representation.
hash	depends on <code>string</code> .  It refines classes with a <code>hash</code> function and implements some hashes collections— <code>HashMap</code> , <code>HashSet</code> , etc.
io	depends on <code>string</code> .  It define input/output related classes likes <code>File</code> . It also refines the <code>Any</code> by adding the <code>print</code> method.
string_search	depends on <code>string</code> .  This module is about string searching and matching. It also implements the Boyer-Moore fast string searching algorithm.

### 3.6 Base Classes

Kinds of Classes Classes can be classified into 4 categories:

- Primitive Classes: They correspond to the primitive values of the computer. They are mainly defined in the `kernel` module. Primitive classes and their super-classes can have some restrictions.
- Built-in Classes: They are known by the compiler since they have literals representation or are needed in some particular statements or expressions. Obviously, they include primitive classes.
- Base Classes: They are the classes defined or imported by the `standard` module. They include Built-in classes.
- User Classes: They are the other classes and are defined by PRM programmers.
- Literal Value Some built-in classes have literal value: programmers can create objects without explicitly instantiate them. The table ?? summarises build-in classes and gives example of literals values.

	examples	PRM types
integers	51, -85	Int
floats	5.5, -.05, 8.0	Float
characters	'a', 'n'	Char
strings	"hello!", "I"	String
Booleans	true, false	Bool
range	[1..5], ['a'..'b']	Range[Int], Range[Char]
arrays	[1,5,6], ['a','b','c']	Array[Int], Array[Char]
void	nil	None

Table 1: The Basic Types

### 3.6.1 Any

The `Any` built-in class is the root of the class hierarchy.

Here some notable properties that will be inherited or redefined in other classes:

```

Class Any
public # Equality tests
  def ==(a: Any): Bool
    # The identity equality
    # Return 'true' if 'self' and 'a'
    # are the same object
    # False otherwise
    # /\ This method cannot be redefined

  def !=(a: Any): Bool
    # Return 'not self == a'
    # /\ This method cannot be redefined

  def =(a: Any): Bool
    # The value identity
    # Return 'true' if 'self' and 'a'
    # have the same 'contents'

  def !=(a: Any): Bool
    # Return 'not self = a'

```

```

public # String
  def to_s: String
    # Convert 'self' to a human readable form

private # Basic IO
  def print(a: Any*)
    # For each argument, output the
    # human readable form ('to_s')
end

```

### 3.6.2 Int

The `Int` primitive class represents internal machine integers. Literals are sequence of digits.

→ `??`: operator

Notable `Int` properties are their operators. Many of them cannot be redefined.

### 3.6.3 Float

The primitive `Float` class represents internal machine float numbers.

Literals are sequence of digits a dot and another sequence of digit.

### 3.6.4 Character

The `Char` primitive class represents characters.

Literals are delimited with single quotes. Characters can include escaped sequence—Table `??`.

Escape sequence	Meaning
<code>\n</code>	ASCII Linefeed (LF)
<code>\r</code>	ASCII Carriage Return (CR)
<code>\t</code>	ASCII Horizontal Tab (TAB)
<code>\0</code>	ASCII Nul (NUL)
<code>\\</code>	Backslash (\)
in strings only...	
<code>\"</code>	Double quote (")
<code>\#</code>	Hash (#)

Table 2: Char and String Escape Sequences

### 3.6.5 String of Characters

The `String` built-in class represents strings of characters—i.e. pieces of text.

Literals are delimited with double quotes. As with character, some escape sequence can be used to represent some characters—Table `??`.

Extended String Literal String literal can contain embedded expression using `#{}8`:

<sup>8</sup>The `#{}8` notation comes from RUBY

```

print("4 + 8 = #{4+8}")
# Output '4 + 8 = 12'

let h := "Hello"
let w := "World"
let hw := "#{h} #{w}"
print(hw)
# Output 'Hello World'

```

Programmers are encouraged to use extended strings because they are better for internationalisation since only entire sentences should be presented to the translator. Therefore having one string `bj1bj2"Replace #o with #o?"` is obviously better than having the three strings `"Replace ", " with ", and "?"`.

### 3.6.6 Boolean

The `Bool` primitive class represents the two Booleans `true` and `false`.

Pseudo-operators `and`, `or` and `not` There is also three special operators on Booleans that are not operators of the `Bool` class: `and`, `or`, and `not`. These keyword have a special status since `and`, `or` are lazy and will return the value, as soon as falsity (`and`) or truth (`or`) is established.

Booleans are used in many places, especially:

- ??
  - Conditional statement.
- ??
  - While loop.
- ??
  - Check statement.

### 3.6.7 Collection

The `Collection` built-in generic abstract class is the root of the collection class hierarchy.

Here some notable deferred method:

```

class Collection[E: Any]
  def is_empty: Bool
  # Is there no item in the collection ?

  def length: Int
  # Number of items in the collection.

  def has(item: E): Bool
  # Is 'item' in the collection ?
  # Comparisons are done with =

  def iterator: Iterator[E]
  # Get a new iterator on the collection.
end

```

### 3.6.8 Iterator

The `Iterator` built-in generic abstract class is mainly used with collections. Instances of the `Iterator` class generates a series of elements, one at a time.

Here some notable deferred method:

```
class Iterator[E: Any]
  def item: E
  # The current item.

  def next
  # Jump to the next item.

  def is_ok: Bool
  # Is there a current item ?
end
```

### 3.6.9 Array

→ ??: genericity

The `Array` built-in generic class is a subclass of `Collection`. It is also the preferred representation form of collections of items.

**Array Literal** Literals representation use brackets and elements are separated with comma. Example:

```
let ai := [5, 4, 6, 1]
print(ai.length)      # Output '4'
print(ai.has(5))      # Output 'true'
print(ai.has(9))      # Output 'false'
print(ai)              # Output '5461'

let ai := ["Hello", " ", "World"]
print(ai.length)      # Output '3'
print(ai)              # Output 'Hello World'
```

**Static Type of Array Literal** Since `Array` is generic, the type of literal should be computed. Literals expressions are valid if and only if there is a unique more general static type. The static type of the literal expression is build with this type.

```
let ai := [4, 5]      # more general type: Int
                    # Therefore ai is an Array[Int]
let ae := [4, '5']   # more general types: Int and Char
                    # Therefore, it is an error

let x: Any
let aa := [4, '5', x] # more general type: Any
                    # Therefore ai is an Array[Any]
```

**Array Constructor** When literals are invalid, one can use the `init` array constructor therefore explicitly precise the desired static type:

```
let ae := new Array[Any](4, '5') # OK
```

→ ??: multiple arguments

Arrays are used in multiple argument procedures. The `Array init` constructor uses multiple arguments.

### 3.6.10 Range

The built-in a generic class `Range` is for `Discrete` elements. It represents intervals between a first element and a last element.

There are two kinds of ranges, inclusive ranges and exclusive ranges.

**Inclusive Range** They include the first element, the last element and each element between them. Their literals use the `[a..b]` notation (`a` and `b` can be any expression): Literals are valid if first element and the last element have the same static types.

```
let x := [1..5] # x is Range[Int]
print(x.length) # Output '5'
print(x.has(0)) # Output 'false'
print(x.has(1)) # Output 'true'
print(x.has(4)) # Output 'true'
print(x.has(5)) # Output 'true'
print(x.has(6)) # Output 'false'
```

**Exclusive Range** Like inclusive range but they exclude the last element. Their literals use the `[a..b[` notation:

```
let x := [1..5[ # x is Range[Int]
print(x.length) # Output '4'
print(x.has(5)) # Output 'false'
```

→ ??: for loop

Ranges are often used in for loops.

### 3.6.11 None

`None` is the absurd class, it is the class that specialise each other except primitives classes. It is not a “real” class since it does not have a definition. It is also the only class that cannot be specialised or refined.

`nil` `None` has only one instance `nil`, often called the *void object*. `nil` correspond to the `null` constant of JAVA or the `Void` object of EIFFEL.

Each method invocation on `nil` will fail. The only exception are equality operators `=`, `==`, `!=`, and `!==`. Therefore, it is frequent to verify if potential receiver is `nil` before sending a message:

```
def safe_array_length(a: Array[Int]): Int
do
  if a = nil then
    return 0
  else
    return a.length
  end
end
```

### 3.6.12 Sys

The `Sys` built-in class

**Program Start** When the program starts, it instantiates the `Sys` class then invoke the `init`

procedure.

**Sys Definition** The `Sys` class is defined in the `kernel` module as follow:

```
class Sys
private
  def init
    # The entry point of the program
    do
      init_begin
      main
    end

    def init_begin
      # Initialisation of library objects
      do
      end

      def main
        # The main part of the program
        do
        end
      end
    end
end
```

The `init_begin` procedure is used to initialise library object. For instance, the module `io` build the standard IO file objects in `init_begin`

→ ??: procedural style

The `main` procedure corresponds to the main statements of the module.

## 4 The Base Language

This section is dedicated to the base language programming.

### 4.1 Source Structure

PRM is a line-oriented language. PRM statements are terminated at the end of a line unless the statement is obviously incomplete—for example if the last token on a line is an operator or comma. A semicolon can be used to separate multiple expressions on a line.

**Comment** Comments start with ‘#’ and run to the end of the physical line. They are ignored during compilation. Currently, there are no multi-line comments.

```
# One statement, one line
a := 1

# Two statements, one line
b := 2; c := 3

# One statement, two lines
d := 4 + 5 +
    6 + 7

# Two statements, two lines
# But the second one clearly do nothing
# and may provoke a warning during compilation
e := 8 + 9
```

## 4.2 Name

PRM names are used to refer to variables, properties (methods and attributes), classes, and modules. The first character of a name helps PRM to distinguish its intended use.

**Reserved Names** Certain names, listed in Table ??, are reserved and should not be used as variable, property, class, or module names.

and	as	check	class	constructor	def	def_attr
def_read	deferred	do	else	elsif	end	extern
false	for	if	import	in	inherit	intern
isa	let	let	new	nil	not	not
once	or	private	public	rename	return	self
then	true	until	while			

Table 3: Reserved Names

In these descriptions, a lowercase letter means the characters “a” though “z”, as well as “\_”, the underscore. An uppercase letter means “A” though “Z”, and digit means “0” through “9”. Name characters mean any combination of upper- and lowercase letters and digits.

A local variable name, a method name, or a module name consists of a lowercase letter followed by name characters. Examples: `foo`, `foo_bar_baz`, `_x`.

A class name starts with an uppercase letter followed by name. Examples: `Int`, `Any`.

An attribute name starts with an “at” sign (“@”) followed by a lowercase letter, followed by any name characters. Examples: `@name`, `@x`, `@_`.

## 4.3 Type

PRM is a statically typed language, it’s mean that “things” should have a static type.

**Type Annotation** Type annotation use the PASCAL notation style: `thing: Type`. Such type annotations are used in the following places:

- ??
  - Local variable declarations.
- ??
  - Method signatures—for parameters, and in functions for the return type.
- ??
  - Generic classes—for formal generic parameters.

**Type Language** A type is:

- ??
  - A non generic class—e.g. `Int`, `Car...`
- ??
  - A generic type—e.g. `Array[Int]`, `Iterator[Car]`
- ??
  - A formal generic parameter.

Covariant Type Policy As Eiffel, PRM uses a covariant typing policy.

The covariant typing policy allows the programmer to redefine properties with a more specific signature:

```
class Food
end
class Grass
inherit Food
end
class Animal
  def eat(f: Food)
    ...
end
class Cow
inherit Animal
  def eat(g: Grass)
end
```

However, such a typing policy is unsafe. In concrete term it means that in some case, type error may occur at runtime and stop the program execution<sup>9</sup>.

## 4.4 Expression

There are ten kinds of expressions:

- ?? • The current receiver `self`.
- ?? • Literal values.
- ?? • Function invocation.
- ?? • Variable read.
- ?? • Attribute read.
- ?? • Exported attribute read.
- ?? • Object creation.
- ?? • Boolean pseudo operators.
- ?? • Type checks.
- ?? • Once expressions.

### 4.4.1 Type Checks

Type check can be used to tests if an object is an instance of a giver class (or an instance of a subclass). However, since PRM uses bounded genericity and a covariant typing policy, there is a very few need of such type checks.

→ ??: assertions; ??: The majority of their use corresponds to assertions and to assignment attempts

`isa` Type checks can be explicitly performed with the `isa` keyword:

---

<sup>9</sup> In a future version, runtime type error will raise an exception.

```

let a: Any
a := 5
print(a isa Int) # Output 'true'
print(a isa Any) # Output 'true'
print(a isa Bool) # Output 'false'

```

#### 4.4.2 Once Expression

↔ This part of the specification is not stable and may change in a future version.

This expression is constituted by the `once` keyword followed by another expression called *sub-expression*:

```

let x := once "Message"
let y := once new Car

```

The semantic of the `once` expression is to evaluate the sub-expression only one time during the execution of the program. Successive evaluations of the `once` expression will return the first evaluated value.

Once expressions are mainly used to create singletons and to perform some local optimisation—it is often used with literals string and arrays.

Examples:

```

def only_one(i: Int): Int
do
    return once i
end
print(only_one(1)) # Output '1'
print(only_one(2)) # Output '1'
print(only_one(3)) # Output '1'

```

```

class Person
    ...
end

def immortal: Person
# There can be only one
do
    return once new Person
end

```

The `once` expression is a generalisation of the Eiffel `once` keyword.

#### 4.5 Statement

There are nine different statements in PRM: statement block, local variable declaration, assignment, procedure invocation, conditionals, while loop, for loop, return and check.

Statements are always defined in a statement block or belong to the main module statements.

### 4.5.1 Statement Block

Blocks of statements often start with the `do` keyword and are ended with the `end` keyword:

```
# outside
do
    # inside
    do
        # more inside
    end
    # inside again
end
# outside again
```

→ `??`: `main`  
→ `??`: `if`

The only exceptions are the main statements of the program and the statements of the `if` statement.

PRM statement blocks are slightly different from other language ones:

- `do/end` differs from usual `begin/end` of others PASCAL-style languages. In fact, it is almost a 50% less characters<sup>10</sup>.
- Curly brackets from C-influenced languages do not fit with the overall PRM PASCAL style.

In a near future version, statement block will be extended to allow exception management. A potential syntax can be:

```
do
    ...
rescue e: IOException
    ...
rescue e: EmptyListException
    ...
rescue
    ...
finally
    ...
end
```

### 4.5.2 Local Variable Declaration

`let` The `let` keyword is used to declare local variables:

```
let i: Int           # i is an integer
let j, k: String    # j and k are strings
```

Initial Value and Type  
Inference

An initial value can be directly assigned with the local variable. If the initial value is present and the static type absent, the static type of the local variable is implicitly the static type of the initial value. Examples:

<sup>10</sup> `do/end` statement blocks are used in some languages (PL/1, REXX). In RUBY, `begin/end` corresponds to statement blocks, although `do/end` and curly brackets correspond to closures.

```

let j: Int := 5 + 3      # an integer with the value 8
let k := j + 1         # an integer with the value 9
let c := new Car("Blue") # a blue car

```

Default Value Without an explicit initial value, local variables are initialised at 0 for Int, '0' for Char, false for Bool, and nil for the other types.

Visibility The visibility of local variable runs from its declaration until the end of the current block.

```

do
    # 'i' is not yet known
    let i: Int
    # 'i' is known
    do
        # 'i' is still known
    end
    # 'i' is still known
end
# 'i' is no more known

```

◁ → One can declare in a same block two local variables with the same name. The last declared will mask the others. However, the compiler may produce a warning.

```

let i: String
i := "foo"
do
    let i: Int      # Warning !
    i := 5         # Correct, the Int variable
                  # masks the String variable
end
i := "bar"        # Correct, the Int variable
                  # is no more known

```

PRM encourages the use of local variables to store intermediate results. Therefore, it allows the programmer to have a liberal use of local variables:

- New local variables can be declared when they are needed. Some languages like Eiffel or Smalltalk only allow local variable declaration at begin of subprograms. Some other languages, like Ada, Modula 3, C, or Lisp, only allow them at begin of statement blocks.
- The static type is optional any can be inferred from the initial value. This feature is quite rare in statically typed languages even if it was one of the first that appears in during the PRM specification development. Modula-3 has it and it is planned for the future C# 3.0.

### 4.5.3 Assignment Statement

The assignment statement uses the quite common := and is widely used in PRM programs:

- Local variable assignment.
- Attribute assignment.

→ ??

- ??
  - ??
  - ??
  - ??
- Assignment procedure.
  - Implicit parameter declaration.
  - Exported attribute assignment.
  - Initial local variable value.

Conformance An assignment `a := b` is statically valid if the static type of the left-value `a` is a subtype of static type of the right-value `b`:

```
let x: Int
x := 4      # OK
x := 'a'    # Error
let y: Any
y := x      # OK
y := 'a'    # OK
```

Assignment Attempt The assignment attempt use the EIFFEL `?=` notation. It works exactly like the assignment, except that conformance is not checked statically but at runtime. If an assignment attempt fails, the program execution will stop<sup>11</sup>.

Example:

```
let x: Any
let y: Int
let z: Char
x := 5
y ?= x # OK
z ?= x # Error at run-time
```

→ ??: type check

In order to avoid run-time error, dynamic types can be checked before any assignment attempts:

```
let x: Any
let y := 0
...
if x isa Int then
    y := x
end
```

#### 4.5.4 Conditional Statement

Conditionals use the standard `if then elsif else` keywords. The `elsif` and `else` parts are optional, and there can be more than one `elsif` part.

```
def game(guess: Int, solution: Int): String
# A simple game
do
    if guess > solution then
        return "It's less"
    elsif guess < solution then
        return "It's more"
    else
        return "Correct"
    end
end
```

<sup>11</sup> In a future version, it will raise an exception.

```
end
```

The question about the one-liner `if` was raised but we did not find a clear and concise syntax. We find only the PERL-ish post-test `if`:

```
instr if expr
```

but it does not satisfy us.

#### 4.5.5 While Loop

The while loop is the main PRM loop structure. It is constituted with a Boolean expression and a statement block:

```
def gcd(x, y: Int)
# The greatest common divisor between x and y
# using the Euclid's algorithm
do
    while y != 0 do
        let t := y
        y := x % y
        x := t
    end
    return x
end
```

#### 4.5.6 For Loop

The `for` loops are used for collection traversal. It is quite different of the C or C++ `for`. In fact it is comparable to the PERL `foreach` or to the new JAVA 5.0 `for/in` loop

This loop can be used with any expression subtype of the built-in class `Collection`, even those defined by the programmer. Since `Array` and `Range` are subclasses of `Collection`, here are two examples:

```
let pricelist := [34.50, 21.95, 4.95, 8.45]
for price in pricelist do
    let gstprice := price * 1.1
    print("Price is ", gstprice, " including GST.\n")
end
```

```
for i in [0..10] do
    print("Value of i: ", i, "\n")
end
```

In fact, the `for` loops are no more than a while loops adaptation. The last example with a range is equivalent with:

```
do
    let x := [0..10].iterator
    while x.is_ok do
        let i := x.value
        print("Value of i: ", i, "\n")
        x.next
    end
end
```

```
end
end
```

#### 4.5.7 Return Statement

This statement has two usages according to the kind of method it is used: in a procedure or in a function. In both case, it terminate the method.

In a function, the return statement is mandatory and must provide a result value that is conform to the declared result type in the signature of the function:

```
def sign(i: Int): Int
do
  if i > 0 then
    return 1
  elsif i < 0 then
    return -1
  else
    return 0
  end
end
```

In a procedure, the return statement is optional and must not provide a value:

```
def stars(nb: Int)
do
  if nb <= 0 then
    println("I want stars.")
    return
  end

  println("A star: *")
  for nb in [2..i] do
    println("Another star: *")
  end
end
```

#### 4.5.8 Check Statement

The check statement is about correctness. It corresponds to assertions and helps to check validity of programs and identify bugs.

→ ??: Bool

A check statement is constituted by the `check` keyword, optionally an assertion label followed by a colon, then a Boolean expression.

```
def hello(name: String)
  check correct_name: name != nil and
    not name.is_empty
  print("Hello ", name)
end
```

During its execution, the program will halt on the check if the expression is evaluated to `false`<sup>12</sup>.

## 5 A PRM Conclusion

The PRM language focuses expressive, clear, simple, and coherent concepts in a statically typed object-oriented language whereas the other languages of the same family rarely focus simplicity.

Currently, the PRM specification is almost complete however some characteristics are currently instable and others are missing like constant values, enumeration types, introspection, module visibility (import/export), exceptions, contracts, regular expressions...

The standard module hierarchy needs also to be extended. Actually there is less than 6000 line of code in the base modules—that is not a lot even if PRM has a concise syntax.

The last work is about the compiler and other tools. Actually, the PRM compiler, `prmc`, is just a prototype and does not yet implement entirely the current specification. However, it produces efficient executable.

---

<sup>12</sup> In a future version, it will raise an exception.

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