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A new exact algorithm to solve the Multi-trip vehicle routing problem with time windows and limited duration

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Abstract

This article tackles the multi-trip vehicle routing problem with time windows and limited duration. A trip is a timed route such that a succession of trips can be assigned to one vehicle. We provide a two-phase exact algorithm to solve it. The first phase enumerates possible ordered lists of client matching trip maximum duration criterion. The second phase uses a Branch and Price scheme to generate and choose best set of trips to visit all customers. We propose a set covering formulation as the column generation master problem, where columns (variables) represent trips. The sub-problem selects appropriate timing for trips and has a pseudo-polynomial complexity. Computational results on Solomon's benchmarks are presented. The computational times obtained with our new algorithm are much lower than the ones obtained in the sole exact algorithm previously published on this problem.

Keywords : Vehicle routing, Time windows, Multi-trip, Column Generation, Dynamic programming, Branch and Price

| Instance | Root Solution | Solution | Total Time | Root time | Phase1 time | iter | column |
|----------|---------------|----------|------------|-----------|-------------|--------|--------|
| c201-25 | 646.51 | 659.02 | 1.561 | 0.046 | 0 | 64 | 124 |
| c202-25 | 634.772 | 653.37 | 45.819 | 0.499 | 0.031 | 564 | 555 |
| c203-25 | 626.017 | 646.4 | 247.189 | 1.795 | 0.031 | 1133 | 854 |
| c204-25 | 592.06 | 602.46 | 252.825 | 5.09 | 0.047 | 673 | 1268 |
| c205-25 | 607.913 | 636.39 | 38.325 | 0.141 | 0 | 1209 | 304 |
| c206-25 | 603.333 | 636.39 | 637.612 | 0.296 | 0.015 | 17034 | 576 |
| c207-25 | 588.783 | 603.22 | 98.273 | 0.718 | 0.015 | 1159 | 725 |
| c208-25 | 597.348 | 613.2 | 38.154 | 0.39 | 0.015 | 580 | 484 |
| r201-25 | 757.79 | 762.43 | 0.234 | 0.046 | 0 | 8 | 174 |
| r202-25 | 645.78 | 645.78 | 0.796 | 0.453 | 0.031 | 7 | 635 |
| r203-25 | 620.177 | 621.97 | 2.216 | 0.89 | 0.078 | 12 | 859 |
| r204-25 | 575.655 | 579.68 | 5.026 | 1.561 | 0.062 | 20 | 1005 |
| r205-25 | 626.48 | 634.09 | 0.827 | 0.202 | 0.015 | 21 | 375 |
| r206-25 | 596.74 | 596.74 | 0.686 | 0.515 | 0.031 | 4 | 713 |
| r207-25 | 583.658 | 585.74 | 3.496 | 0.577 | 0.046 | 18 | 783 |
| r208-25 | 575.616 | 579.68 | 6.681 | 1.53 | 0.047 | 23 | 1181 |
| r209-25 | 598.107 | 602.39 | 1.67 | 0.281 | 0.016 | 22 | 562 |
| r210-25 | 620.293 | 636.15 | 7.914 | 0.359 | 0.031 | 73 | 870 |
| r211-25 | 568.54 | 575.91 | 25.805 | 1.295 | 0.046 | 198 | 1538 |
| rc201-25 | 984.438 | 988.05 | 1.373 | 0.11 | 0 | 72 | 342 |
| rc202-25 | 837.557 | 881.49 | 25.664 | 0.483 | 0.031 | 894 | 880 |
| rc203-25 | 705.217 | 749.15 | 64.271 | 0.327 | 0.031 | 814 | 1088 |
| rc204-25 | - | - | - | - | - | - | - |
| rc205-25 | 808.579 | 840.35 | 3.746 | 0.249 | 0.015 | 137 | 598 |
| rc206-25 | 726.097 | 761.03 | 35.703 | 0.281 | 0 | 2006 | 767 |
| rc207-25 | 646.457 | 738.87 | 71807.37 | 0.859 | 0.031 | 452858 | 8212 |
| rc208-25 | - | - | - | - | - | - | - |

Table 1: Results on the Solomon’s benchmark (25 customers) with t_{max} value set (75; 220)

instances have not been solved within a limit on computing time set to 30 hours.

Tables 1 and 2 present these results. For each instance, we have the root solution cost and the solution cost of the Branch and Price scheme used in Phase 2, the computation time for both phases (Total time), the computation time of Branch and Price root in Phase 2 (Root time), the computation time of Phase 1 (Phase1 time), the number of iteration (iter) and the number of generated columns (column).

We can note, as for the VRPTW, there is great variation for time resolution between instances of the same class. We can also note a significant increase of total computation time between the instances with 25 and 50 customers.

5.4 Impact of the dominance rule in phase 1

In this section, we evaluate the impact of the dominance rule when feasible trips are generated with the elementary shortest path algorithm with resource constraints in Phase 1. The number of available vehicles (U) was set to 2 for these tests.

In Table 3, for each instance, we compare the following criteria : the number of generated structures in Phase 1 (# trips), the computation time for both phases (Total time) and the computation time of Phase 1 (Phase1 time), with (Dom) and without (No Dom) dominance

| Instance | Root Solution | Solution | Total Time | Root time | Phase1 time | iter | column |
|----------|---------------|----------|------------|-----------|-------------|-------|--------|
| c201-50 | 1309.63 | 1324.32 | 1.912 | 0.17 | 0.016 | 14 | 343 |
| c202-50 | 1280.44 | 1310.79 | 6067.16 | 1.882 | 0.25 | 12085 | 1635 |
| c203-50 | 1236.3 | 1247.77 | 386.395 | 9.503 | 0.5 | 176 | 2681 |
| c204-50 | 1181.61 | 1195.51 | 3351.04 | 63.501 | 0.796 | 620 | 6824 |
| c205-50 | 1245.19 | 1265.61 | 771.309 | 0.4 | 0.031 | 3310 | 683 |
| c206-50 | 1241.5 | 1262.47 | 6121.5 | 0.781 | 0.063 | 14776 | 1025 |
| c207-50 | 1203.8 | 1216.24 | 1675.43 | 3.434 | 0.187 | 2038 | 1624 |
| c208-50 | 1231.31 | 1249 | 4781.94 | 1.201 | 0.109 | 8830 | 1284 |
| r201-50 | 1397.07 | 1405.52 | 6.529 | 0.19 | 0.109 | 69 | 699 |
| r202-50 | 1221.82 | 1229.91 | 86.394 | 9.583 | 0.766 | 101 | 3791 |
| r203-50 | 1101.63 | 1104.51 | 67.246 | 18.646 | 2.156 | 30 | 5314 |
| r204-50 | 1010.65 | 1031.72 | 22044.9 | 45.245 | 5.266 | 9733 | 7452 |
| r205-50 | 1219.64 | 1230.26 | 63.471 | 2.032 | 0.297 | 228 | 1930 |
| r206-50 | 1150.62 | 1154.53 | 34.229 | 16.453 | 1.625 | 15 | 4887 |
| r207-50 | 1086.15 | 1094.83 | 830.624 | 26.448 | 3.016 | 481 | 6056 |
| r208-50 | 1010.65 | 1031.72 | 28145.3 | 57.142 | 5.375 | 11521 | 9702 |
| r209-50 | 1126.47 | 1143.91 | 1619.44 | 8.342 | 0.765 | 2552 | 3629 |
| r210-50 | 1152.64 | 1162.14 | 273.593 | 17.915 | 1.64 | 268 | 4871 |
| r211-50 | - | - | - | - | - | - | - |
| rc201-50 | 1814.12 | 1876.06 | 16.153 | 0.28 | 0.015 | 432 | 463 |
| rc202-50 | 1678.02 | 1763.48 | 3538.58 | 0.991 | 0.078 | 30361 | 1098 |
| rc203-50 | - | - | - | - | - | - | - |
| rc204-50 | 1406.73 | 1457.3 | 33563.8 | 7.791 | 0.281 | 68145 | 3699 |
| rc205-50 | 1698.02 | 1780.1 | 4160.9 | 0.59 | 0.047 | 39122 | 1880 |
| rc206-50 | - | - | - | - | - | - | - |
| rc207-50 | - | - | - | - | - | - | - |
| rc208-50 | - | - | - | - | - | - | - |

Table 2: Results on the Solomon's benchmark (50 customers) with t_{max} value set (75;220)