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Models, algorithms, and programs for phylogeny reconciliation

**Keywords:** phylogeny; gene duplication, loss, horizontal gene transfer; parsimony; probability; reconciliation.

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Models, algorithms, and programs for phylogeny reconciliation

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Abstract. Gene sequences contain a gold mine of phylogenetic information. But unfortunately for taxonomists, this information does not only tell the story of the species from which it was collected. Genes have their own, complex histories, which record speciation events of course, but also many other events. Among them, gene duplication, transfer and loss are especially important to identify and differentiate from speciation events. These events are not only crucial to account for when reconstructing the history of species, but they also play a fundamental role in the evolution of genomes, the diversification of organisms and the emergence of new cellular functions.

We present a state of the art of reconciliations between gene and species trees, which is a rigorous approach to identify the nature of the various events that mark the evolution of a gene family. Existing models and algorithms of reconciliations are reviewed, and difficulties to model gene transfer are discussed. We also compare different reconciliation programs together with their advantages and disadvantages and summarize reconciliation results obtained on numerous gene families.

Keywords: phylogeny; gene duplication, loss, horizontal gene transfer; parsimony; probability; reconciliation.

1 Introduction
The systematic reconstruction of gene phylogenies from a wide variety of organisms reveals an unforeseen diversity of histories, which are difficult to understand through a simple pattern of species evolution. These differences underline the importance of modelling the factors that specifically affect gene evolution. Population genetics has brought a powerful framework in which to understand the pattern of evolution of alleles through a succession of speciations. Under these multispecies coalescent models, the differences among gene phylogenies turn out to be informative for estimating population size, structure, and speciation times [1]. On another side, models of genome evolution, which account for the mechanisms of gene duplication, gene loss and Lateral Gene Transfer (LGT), are emerging. The development of such models is crucial especially to understand the evolution of unicellular organisms where LGT has played a major role [2], and more generally to clarify homology relationships among genes. The combination of duplication, transfer and loss may have been such in the history of life that no single phylogenetic marker can be considered reliable to readily infer the history of species. The following studies arguably rely on the development of models of duplication, transfer and loss [3]: reconstructing the tree of life; understanding the principles of genome evolution, the role of transfer in species adaptation, and the contribution of duplication and transfer to the evolution of new functions.
As in the multispecies coalescent model, reconciliation models consider a species tree within which a gene can evolve (Figure 1). Leaves of the species tree and the gene tree are associated and specific events are invoked to allow the gene tree to evolve within the species tree with its own phylogeny. Partial models accounting for duplication and loss alone [4,5,6], or LGT and loss alone [7,8,9] have been described. These models are realistic in particular biological cases (resp. multicellular organisms where LGT is rare and gene families for which functional redundancy can be detrimental). Here, we focus on models that account for DL (Duplication and Loss) and DTL (Duplication, Transfer and Loss) events.

Figure 1

Reconciliation is a popular approach to infer orthology relationships [10–14]. It has been shown on real datasets that such phylogeny–based methods are relatively accurate [15]. Other methods that rely solely on sequence similarity can lead to false positives and false negatives [16]. Reconciliation approach has applications in other areas such as estimating DTL rates [17,18], gene tree inference based on molecular substitution and gene evolution [19–21], and genome phylogeny reconstruction from discordant gene trees [5,22]. Reconciliation can also be used to study coevolution between parasite and their hosts (parasitology) and between organisms and their living areas (biogeography) [23–25].

Sections 2.1 and 2.1 review the main approaches of reconciliation with DL and DTL events, respectively. Section 2.1 also explains why transfers are hard to consider and describes several methods to handle them. Section 2.1 presents the pros and cons of available reconciliation software. Section 2.1 presents results of reconciliation analyses on a real dataset that covers the three domains of life.

2 Models and algorithms for reconciliation

As for phylogenetic reconstruction, parsimony and probabilistic frameworks have been developed for reconciliation inference. The former are based on explicit discrete models of gene evolution and search for a reconciliation of minimal cost given costs for individual evolutionary events. Probabilistic methods rely on continuous models and seek a reconciliation with maximum likelihood or maximum posterior probability. Parsimony methods are faster but use less realistic models than probabilistic ones.

2.1 Evolutionary scenarios with duplications and losses

Historically, DL models are the first ones to have been considered and remain relevant when only considering eukaryotes.

Parsimony models
Numerous reconciliations can exist between a gene tree $G$ and a species tree $S$. For instance, Figure 1 depicts two reconciliations $R1$ and $R2$ for the same trees, which differ by the location in $S$ of the duplication $u$ of an ancestral gene of $G$ (and by the induced losses). Several ways to represent reconciliations have been proposed. One of the most widespread is through a tree called a *Reconciled Tree* [4,6,26], here denoted RT, and defined as follows: (1) the clade of each node of RT has to be present in $S$; (2) for each internal node of RT, the clades of its children are either equal (duplication) or disjoint (speciation) and (3) $G$ has to be obtained from RT by pruning some of its subtrees (see Figure 2 below).

**Figure 2**

The model of reconciled tree (Figure 2; RT) is sometimes less intuitive than a drawing of $G$ embedded into $S$ (Figure 2; R1). A model that leads to such a representation is defined in [27], where each ancestral gene of $G$ is mapped either on a vertex (speciation) or a branch (duplication) of $S$.

Based on the reconciled tree formalism, [28] have introduced an architecture that allows to describe and explore the whole set of reconciliations between $G$ and $S$. They proved that a simple polynomial time algorithm, called LCA mapping, allows to find one of the *Most Parsimonious Reconciliations* (MPR). This well-known algorithm [4,6,29] can be implemented to run in time linear w.r.t. the number of nodes in $G$ [30]. The LCA mapping maps each gene $u$ of $G$ onto the most recent species $x$ of $S$ such that each contemporary gene that descends from $u$ belongs to a contemporary species that descends from $x$ (see Figure 3). According to this LCA mapping, a node $u$ of $G$ is a duplication if, and only if, it is mapped to the same vertex of $S$ than one of its children. Otherwise, $u$ is a speciation. Based on this rule, the so-called LCA reconciliation is computed as follows. For each node $u$ of $G$, mapped on a vertex $x$ of $S$, if $u$ is a duplication, it is located on the branch immediately above $x$, otherwise $u$ is a speciation placed on $x$. For instance, R1 of Figure 1 is the LCA reconciliation, while R2 is not.

**Figure 3**

Given a cost for each event (i.e. duplication, loss and speciation), the parsimony score of reconciliations can be defined in various ways. It can be either the sum of costs for duplications or the sum of costs for duplications and losses. The LCA reconciliation provides an MPR for the two scores, and is even the unique one in the latter case [28,31]. Given a choice for event costs, numerous reconciliations may exist with near-optimal scores. Some of them could have been optimal if slightly different event costs were used [32]. However, most reconciliation analyses focus on the LCA reconciliation and ignore such near-optimal reconciliations.

All branches of a phylogenetic tree are not equally reliable, bootstrap and posterior probabilities being usual support measures. Uncertainties can have
an important effect on reconciliation analyses. Reconciliation methods are biased when the inferred gene tree is not correct: duplications and losses tend to be placed toward the root and the leaves, respectively, of the species tree [33]. This uncertainty concerning parts of the gene tree can be taken into account by collapsing weakly supporting branches, thus creating polytomous nodes. An MPR can still be computed in polynomial time when S contains polytomies [34], but only approximate algorithms exist when G is polytomous [35,36].

Probabilistic models

On the basis of a Birth–and–Death process [37], a probabilistic model of gene evolution along a species tree S has been developed [38], where each branch has fixed length (in time) and associated duplication and loss rates. This model is used to compute the probability that the evolution of a gene which belongs to the root of S gives rise to the gene tree G and to the reconciliation R. This probability is the likelihood of R, denoted P(G,R), and can be computed in time \( O(n_G n_R) \), where \( n_G \) and \( n_R \) are the number of nodes in G and S. The reconciliation of maximum likelihood can also be computed in time \( O(n_G n_S \log n_S) \). This reconciliation model has been extended into mixed models which additionally consider the probability that G gives rise to the observed sequences. Such extensions assume either a strict molecular clock [38] or a relaxed one [19].

An important breakthrough is due to [39], who developed an efficient algorithm to compute the probability \( P(G) \), which is the sum of \( P(G,R) \) over all reconciliations R. This algorithm allows to compute the posterior probability \( P(R|G) = P(G,R)/P(G) \) in time \( O(n_G^3 n_S) \). This posterior probability is useful to evaluate the reliability of the most likely scenario with respect to other reconciliations, particularly those having close likelihood. A similar algorithm computes the probability that a given ancestral gene of G is a speciation in time \( O(n_G^3 n_S) \). By sampling duplication and loss rates, [40] developed an MCMC that uses the latter algorithm to estimate posterior probabilities of orthology relationships among sequences of a gene family.

[27] have proposed a similar algorithm as [28] to explore the space of reconciliations. The idea is to go from one reconciliation to another thanks to elementary changes in the reconciliation. This algorithm can be used in order to compute or efficiently estimate \( P(G) \) and thus the posterior probability \( P(R|G) \) [41]. Their analysis of the probabilistic landscape of the space of reconciliations shows that the immediate neighbourhood of the MPR (i.e. the LCA reconciliation) contains the most likely ones and covers most of the probability mass of \( P(G) \). This provides an efficient way to obtain very precise posterior probabilities of the most probable reconciliations. These results emphasize the strong relationship between the parsimony and probabilistic paradigms.

When considering a discrete distribution model of duplications parameterized by the branch lengths of S, a maximum likelihood
reconciliation can be computed in time $O(n_c^3 n_e)$ [42]. However, this is achieved without taking into consideration events that left no trace \(^1\) nor loss rates, in contrast with the above Birth–and–Death models. Such limitations can be problematic when losses are prevalent [43]. This approach is computationally more simple but less realistic than the above models.

2.2 Evolutionary scenarios with duplications, losses, and transfers

Computing an MPR becomes a computationally difficult problem when transfers are considered [44], although it can be solved in polynomial time with realistic constraints [45] (e.g. bounding the number of transfers, of genes per species lineage, etc.). This strong contrast in complexity is due to the chronological constraints induced by transfers. A transfer has to be locally consistent, which means that it occurs between two coexisting species. Two (or more) consecutive transfers have also to be globally consistent (see Figure 4). If these constraints are omitted, time inconsistent scenarios can ensue. [46] solves in time $O(n_s^3 n_e)$ such a variant of the MPR problem, where losses are considered a posteriori.

Figure 4

A recent promising approach to handle time constraints is to accept a dated tree $S$ as input. Time consistency can then be ensured locally by checking that donor and receiver branches of a transfer have intersecting time intervals (see Figure 4). This approach has been used in five reconciliation models: Merkle et al. [47], Hadas et Charleston [48], Tofigh [32], Doyon et al. [49] and David et Alm [18]. These approaches differ in their way to handle global consistency and in the degree of generality of their model.

Global time consistency can be ensured with the following approaches: (1) alter the position of the proposed transfers a posteriori; (2) check that all branches involved in a succession of transfers share a sub–interval of time; or (3) use a subdivision of the branches of $S$ into time slices and allow for transfers only between branches of a same time slice. The two–steps strategy of (1) does not guarantee to find an optimal reconciliation. Approach (2) only considers a subset of scenarios due to over–restrictive rules. For instance, reversing the direction of transfers $T_1$ and $T_2$ in Figure 4 leads to a globally consistent scenario that is here however rejected, since the three involved branches do not share a common time subinterval. Approach (3) ensures to find an optimal reconciliation [48,49].

Models have to be general enough to encompass the variability of all possible scenarios involving transfers. In particular, they have to consider the following two possibilities: (a) Transfers where the donor branch Looses its gene copy\(^2\) (TL event for short; R1 Figure 5); and (b) scenarios where

\(^1\) A gene duplication immediately followed by a loss of one of the copies.

\(^2\) That is the gene copy of the donor left no trace in the contemporary species (i.e. it goes extinct).
speciations and duplications are located below the LCA vertex of S. Due to transfers, constraining such events to be located on or above the LCA is no longer guaranteed to be optimal. For instance in Figure 5, R1 is more parsimonious than R2, due to points (a) and (b) above. In R1, the gene lineage (u,b) follows a TL event from (x,A) toward (z,B) and node w is a speciation located below its LCA mapping (vertex y of S).

Figure 5
According to the features introduced above, Table 1 below summarizes the pros and cons of the five reconciliation models presented in [18,32,47–49,52].

<table>
<thead>
<tr>
<th>Input Gene/Species Trees</th>
<th>Model characteristics</th>
<th>Algorithm</th>
</tr>
</thead>
<tbody>
<tr>
<td>Tree G.</td>
<td>Tree S</td>
<td>Transfer with Loss</td>
</tr>
</tbody>
</table>
| Merkle et al. [47,52] | Binary and polytomous; time interval | Binary; dated | No | Only on/above the LCA | Not guaranteed (considered a posteriori) | O(max(n_S, n_G)^
| Hadas et al. [48] | Binary | Network; dated; | No | Anywhere | Guaranteed, with time slices | Polynomial in n_S, n_G |
| Tofigh [32] | Binary | Binary; dated; | Yes | Anywhere | Guaranteed, with time slices | O(n_S^3 n_G) |
| Doyon et al. [49] | Binary | Binary; dated; | Yes | Anywhere | Guaranteed, with time slices | O(n_S^3 n_G) |
| David et Alm [18] | Binary | Binary; dated; | No | Only on/above the LCA | Guaranteed, with simple rules | O(n_S^3 n_G) |

Table 1 [Comparison of five reconciliation models accounting for duplications, losses and transfers]. The models of [32] and of [18,47–49,52] are continuous and discrete, respectively. \( \land \) and \( \lor \) are the size of S and G.

In contrast to scenario with duplication and loss, where the LCA reconciliation is the unique MPR, several optimal reconciliations are possible in the presence of transfers. This multiplicity of optimal scenarios may well lower the confidence we have in a single reconciliation drawn at random by a program from the set of all possible reconciliations. In this context, algorithms that enumerate all MPRs have been proposed [46,49,50].

3 Available programs
TreeMap [51] was the first program proposed for reconciling a gene tree G with a species tree S. A graphical interface is provided with a number of options. However, this program does not deal with dates for nodes of S, and as such cannot ensure time consistency of transfers. Notung [20] reconciles G and S according to the DL model, where at least one of the trees is binary. It has an interface that displays orthology relationships and a command line
version. It can also root or resolve polytomies of G (i.e. nodes with low support) by minimizing the parsimony score.

The reconciliation approach of [47,52] has been implemented in CoRe-Pa. This software includes a reconciliation viewer, an editor for modifying the G and S trees, as well as resampling facilities to evaluate the statistical relevancy of an MPR. It does not require costs for individual evolutionary events as it heuristically seeks those that best fit the reconstructed event frequencies.

The model of [48] has been implemented in a program called Jane [53], which also includes resampling facilities. It additionally allows a visual edition of a reconciliation (its cost is updated accordingly), and can be run from the command-line (for large-scale experiments). Reconciliations are built for a dated tree S, whose dates can be provided by the user. Alternatively, Jane uses a genetic algorithm to find optimal dates (w.r.t. reconciliation costs). Jane also allows controlling the maximal distance between two species that can exchange genes. This latter option is especially relevant when Jane is used for coevolution studies [24].

The reconciliation approach of [49] is implemented in a command-line program called Mowgli [54]. Besides computing an MPR, it also computes the number of optimal reconciliations for G and S trees. This provides an alternative (and much faster) way to measure the statistical significance of a single MPR. The method of [18] is also implemented in a command-line program called AnGST. It deals with phylogenetic uncertainties by inferring G as a combination of bootstrap subtrees that yields the “best reconciliation”.

Last, we note that some of the above software can differ from the models presented in the associated paper. For instance, on several datasets the reconciliations proposed by CoRe-Pa and AnGST have speciations located below the LCA mapping, while they are not supposed to [18,47,52].

In order to compute reconciliations, it is usually necessary to estimate duplication and loss rates. These can be estimated with Cafe [55], which takes as input a dated tree S and the number of genes per species for several gene families. For consistencies of transfers, dates for nodes of S can be obtained by relaxed molecular clock techniques working from molecular sequence [56,57].

4 Experimental results

To compare the main programs cited above, we performed some experiments on the dataset of [18]. It consists of 3983 gene family trees (with an average of 33.0 \(\pm\) 26.5 contemporary genes) globally covering 90 genomes (11 eukaryotic, 12 archaeal, and 67 bacterial) together with a dated species tree. The average cost of reconciliations computed by Mowgli on these families is 64.3 \(\pm\) 53.0, and in 42% of the cases these reconciliations are strictly more parsimonious than those provided by AnGST.

Jane was only tested on one of these gene families as this software accepts dates through its graphical interface only. The dates used by AnGST and
Mowgli were then entered manually together with the same event costs. Figure 6 displays the reconciliations proposed by the three software. For this gene family, Mowgli finds a reconciliation that is more parsimonious than those inferred by the other two software.

Figure 6

The three reconciliations differ according to the number and kind of events, the different models allowing to optimize parsimony at different degrees. When transfers are considered, the above results show the pitfalls in the design of reconciliation models and algorithms. Moreover in this case, an important aspect is that several most parsimonious reconciliations can exist. The prevalence of this effect has not yet been measured, nor the effect of the event costs in this respect. Last, we note that the probabilistic approach has been less developed here than when only considering duplications and losses.

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Key points

Reconciliation is an approach used to depict the evolution of a gene family with respect to the evolution of the species.

Several reconciliation models based on parsimony and probabilistic criteria have been proposed.

There is no agreement on a reconciliation model that deals with Horizontal Gene Transfers.

Bibliographie


