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A novel implementation of the histogram-based technique for measurement of INL of LUT-based correction of ADC

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Abstract

The histogram-based technique is commonly used for testing of Analog-to-Digital Converters (ADC). One of the parameters measured thanks to this technique is the Integral Non Linearity (INL). INL is also used as an initial data related to the ADC performances for the computation of a correction table in case of a LUT-based correction technique. In this context of embedded INL measurement and embedded computation of the table for LUT-based correction of ADC, we propose a new implementation establishing what we consider the best trade-off between silicon area overhead and computing time. We compare our solution with the state of the art: (a) with VHDL-level simulation we compare time performance, and (b) with FPGA placer we estimate the final surface head-out. © 2012 Elsevier Science. All rights reserved

Keywords: Analog-to-Digital Converter (ADC); Look-Up Table (LUT); Integral Non-linearity (INL); Histogram-based ADC correction

1. Introduction

The use of Lookup-Table (LUT) for improving Analog-to-Digital Converter (ADC) performances is common and a 20-years old research topic [1-3]. The fundamental principle consists in integrating a table that contains correction codes pre-computed during production phase, based on non-linearity measurements. In the application, the output code of the ADC is used to address the LUT that substitutes the output code with the correction code stored in the table. This method has demonstrated a great interest to avoid performance loss due to manufacturing process variability. A recent trend [3] tends to integrate resources dedicated to signal generation and computation in order to embed the correction table generation.

The advantage of this approach is that it can be used not only during the production phase to correct errors due to manufacturing process variability, but also in-situ in the application in order to correct for aging effects or specific use conditions.

In this publication we assume that the case study is in the context of LUT-based corrected ADC. According to the significant literature [4-7] on embedded analog signal generators, we assume that the generator will be accurate enough to implement the LUT-based self-correction technique. The study strictly focuses on the digital signal processing required to operate an embedded measurement of INL and to compute the correction table.

Our innovative computation technique reuses the embedded resources for measuring INL and computing the correction table. The main advantages
are: (a) the reduced silicon area overhead and (b) the fast computation time. The reduced silicon area overhead impact directly the economical viability of our solution but allows us to maintain the reliability of the circuit despite of its operation computation is not related to the function. The fast computation time is also mandatory to reduce the time for which the circuit is not operative because of the procedure for correction table computation.

In section 2 we revise the state of the art in LUT-based correction and embedded INL computation techniques. The Section 3 states our three-step method that measures the INL and computes the correction table. In section 4 we present our simulation results and compare them to the state of the art, both for silicon overhead and computing load. At last, section 5 concludes our method for INL computation for LUT-based correction of ADC, with what we consider the best trade-off between silicon overhead and computing timing.

2. State of the art

2.1. LUT computation

The intrinsic purpose of the LUT-based correction of ADC is to compensate non-linearity (NL) errors of the component. The LUT computation has been widely addressed as a research topic. Static [1] and dynamic [2, 3] characteristics of NL have been considered among different solutions. A widely used solution (e.g. [8]) consists in computing the correction table embedded in the LUT, using the measure of the Integrated Non Linearity (INL). The most common technique for INL measurement is called the histogram one [1]. The characteristics of this technique are its high accuracy and the large number of samples required for computation. This parameter is historically measured using a DSP-based measurement setting and an external Automated Test Equipment (ATE) [9, 10].

INL measurement using histogram-based technique suffers from limitations [11]. Indeed defective behaviors such as non-monotonicity or hysteresis are not detected. As a consequence the LUT-based correction technique cannot address such problems. If these non-detected defects are critical for the targeted product, a specific strategy should be defined.

2.2. Techniques for embedded INL measurement

The direct implementation of the histogram-based method for embedded INL measurement is not a viable option as it requires a huge amount of on-chip resources. Alternative implementations of the histogram-based method have been proposed in the literature for BIST purpose, based either on time decomposition [12] or space decomposition [13]. The main interested of the solution developed in [12] is that it offers a very low silicon area overhead but at the price at relatively long test time. In contrast, the solution developed in [13] presents interesting test time features but necessitates significant silicon area overhead.

In order to avoid constraints due to histogram technique, some publications presents non-histogram techniques embedded for BIST purpose [14, 15]. Authors of [14] describe a test method using a noise as stimulus. This method is very interesting as it reduces constraints on stimulus generation. However the integration of this method seems difficult mainly because it requires computing a Fast Fourier Transform (FFT), moreover the integration has not been demonstrated. [15] presents a very interesting test method allowing a concurrent procedure to the normal operation. This method apparently gives accurate measurements of INL, the main drawback is the potential significant silicon area overhead, notably because of the need of an embedded DAC or ADC dedicated to test.

For a BIST purpose, the integration of INL measurement suffers from particular constraints. Indeed dedicated silicon resources should be minimized but the INL measurement should be accurate in order to ensure the test accuracy.

A way to address the issues related to the Built-In Self-test of ADC is to relax constraints on stimulus generation and to rely the test accuracy on digital signal processing such as double-histogram methods [16 -18]. In our case study, the problem is addressed differently. Indeed the integration of efficient signal generators is research topic by itself, providing already significant results [19]. As a consequence, our interest is in the minimization of digital resources and computing time, in order to limit the impact of
the measurement process as it will be done on-chip notably for the computation of correction table.

In an in-situ LUT-based correction context, the constraints are different. Indeed important resources are available for the INL computation (the LUT), on contrary as the LUT computation is done in-situ, the INL measurement should be done as fast as possible. Like for BIST purpose, in order to ensure the accuracy of the correction the INL measurement should be accurate. The most accurate INL measurement technique is the histogram one. As a consequence, we have developed an embedded histogram computation algorithm to optimize the study case constraints: surface, time and accuracy.

3. New technique for fastening the embedded histogram-based INL measurement

In the LUT-based context, we propose an optimized technique for embedded INL measurement. This technique aims to adapt the greedy linear histogram technique according to the available resources; in this case it means the memory used to store the correction table. The purpose is to reduce the computation time. Indeed in the case of in-situ LUT computation, the processing should be done when the ADC is not functioning.

The hardware resources required to implement the LUT-based INL measurement are presented in Fig. 2.

The controller aims to manage the clock domains synchronizing the different blocks. It also commands the multiplexers (MUX #1 and #2) and counter block. The counter block has two functions: counting from 0 to $2^n$ (n the resolution of the ADC) and incrementing its input value.

The INL computation block computes the INL values from the DNL values temporarily stored for each code in the LUT. The correction block determines the correction table based on computed INL values. The correction codes are then loaded in the LUT and used in the application.

The procedure for computing correction codes and loading the LUT is made of three steps.

3.1. First step: initialization of the LUT

The first step of the method consists in initializing all the words from the LUT to -1 in two’s complement. The words are initialized to -1 considering the case of an ideal converter, for which each code should appear ones during the procedure. One increment of a word corresponding to a code during the procedure will set the final value of the word to zero. It means no non-linearity for this code. For a 10-bit ADC, the words are 10-bit long, 5 bits for the integer part and 5 bits for the decimal part. Then each word is initialized at “1111100000”. The size of the word is related to the ADC resolution because this memory will be used later to store the corrected codes of the ADC. The distribution of bits between integer and decimal parts is related to the targeted accuracy of DNL/INL measurements. The precise explanation is given in the following section.
In order to generate the whole addresses of the LUT, the counter block is set by the controller in its first mode, consisting in counting from 0 to \(2^n-1\). In order to provide the -1 value to each word of the LUT, the controller commands the multiplexer #1.

3.2. Second step: Differential Non-Linearity (DNL) computation

The second step consists in computing the DNL using the initialized LUT.

According to fundamentals of histogram-based measurement of DNL [9], the DNL measurement using a ramp stimulus is given by equation 1.

\[
DNL(i) = \frac{H_{exp}}{H_{ideal}} - 1 \quad \text{equation 1}
\]

Where \(H_{exp}\) is the real and \(H_{ideal}\) the ideal number of occurrences of code \(i\). When the expected and the real number of occurrences of code \(i\) are the same, there is obviously no error and the DNL(i) is equal to zero. The accuracy of the measurement is given by equation 2.

\[
\delta_{DNL} = \frac{1}{H_{ideal}} \quad \text{equation 2}
\]

The relation between the total number of samples and \(H_{ideal}\) is given by equation 3.

\[
H_{ideal} = N_T \cdot \frac{FS}{2^nA_{in}} \quad \text{equation 3}
\]

Where \(N_T\) is the total number of samples, \(n\) the resolution of the ADC, \(FS\) the full scale of the ADC and \(A_{in}\) the amplitude of the ramp. In the case of \(FS\) and \(A_{in}\) are equal the total number of samples is equal to \(H_{ideal}\) multiplied by \(2^n\).

Then the operations done in the second step are the following. The controller commands the multiplexers in order to provide to the data input, the counter output and to address the LUT by the ADC output code. The generated stimulus is supposed to be a ramp with a high linearity. Every time the code \(i\) gets out of the ADC, the word at the address \(i\) is incremented by 1. The 1 is added to the decimal part. The increment is done by the counter which is set in its second function by the controller. Now let’s consider the case study given in the previous section. The 10-bit word has been distributed in 5 bits of integer and 5 bits of decimal, which leads to the -1 value coded by the word 111110000 in 2’s complement. The choice of 5 bits for decimal part of word is related to the definition of \(H_{ideal}\), which is set to 32 (\(2^5\)). According to equation 2, this value is high enough to provide an accurate measurement but can be adapted to the given test conditions in order to reduce the test time. By the end of the ramp, each word contains the DNL of the code corresponding to its address. Indeed, for instance, if code \(i\), is fault-free, then \(H_{ideal}\) will be equal to \(H_{exp}\) and the corresponding word will be incremented by \(32=2^5\). By the end the word initially equal to 1111100000 will be 0000000000. As illustration, let us consider the case for which DNL(i) is equal to 1.5LSB. Then \(H_{exp}(i)\) is equal to 80 and the word corresponding to code \(i\) is incremented 80 times providing the value 0001100000. As a consequence, by defining the right initial value (-1), the right size of the decimal part according to \(H_{ideal}\) and by incrementing the word each time the code occurs, the DNL value can be directly read from the final value of the corresponding word without any expensive division as requested by equation 1. In case of a missing code, the corresponding word has not been incremented and the DNL value is -1. Despite missing codes can be identified, the considered correction technique cannot correct this kind of error. To give an explanation, let us consider the worst-case scenario. The output code is a constant value whatever the input voltage is. There are approximately \(2^n-1\) missing codes. As there is no information at the output, rather than a code, it is impossible to know what the input voltage is to correct the corresponding code.

It is important to notice that the controller generates different clock frequencies for the counter in the first and second steps. Indeed for the first step the initialization is done as fast as possible. For the second step the counter is synchronized with the ADC, i.e., operations of addressing a LUT word and incrementing its value are done during an ADC clock cycle.

3.3. Third step: INL and correction computation

The final step consists in computing INL values from DNL values stored in the LUT in the previous step and generating the correction table. In this step, the data output of the LUT is connected to the input of the INL computation block, the data input of the LUT is connected to the output of the correction block and the address input of the LUT is connected to the counter output.
The counter is set by the controller in the function counting from 0 to \(2^n-1\). The output value of the counter is used to address the LUT and to take out the DNL value corresponding to the addressing code. The INL is then computed by the INL block using the following equation

\[
INL(i) = \sum_{k=0}^{i} DNL(k) - 0.5 \times DNL(i)
\]

for \(i = 0..2^n - 1\)  

**equation 4**

The corrected codes provided to the LUT are computed using equation 5.

\[
corr_{code}(i) = i - \lfloor INL(i) \rfloor
\]

**equation 5**

where \(\lfloor \rfloor\) corresponds to the round function.

Once the correction table is loaded in the LUT, only the LUT is used during the normal operation mode: the ADC output code is provided as an address to the LUT, which substitutes this code with the corrected one.

### 4. Simulation results

The objective of the presented work is not to discuss the effectiveness of the considered correction method. This subject has been addressed in [20]. In order to evaluate the performance of the proposed implementation for INL histogram-based measurements, we have focused on two criteria: computation time and silicon area overhead.

References [12] and [13] provide equations to compute the theoretical computation time. Equation 6 is the computation time using the time-decomposition algorithm [12].

\[
T_1 = 2.25 \times \frac{H_{ideal}}{F_s} \times 2^n
\]

**equation 6**

Where \(n\) is the converter resolution, \(F_s\) the sampling frequency.

Equation 7 is the computation time using the space decomposition algorithm [12].

\[
T_2 = 3 \times \frac{H_{ideal}}{F_s} \times 2^n
\]

**equation 7**

Table 1: Comparison of computing time for different embedded histogram-based methods.

<table>
<thead>
<tr>
<th>Implementation</th>
<th>Fs=1MHz</th>
<th>Fs=20MHz</th>
<th>Fs=50MHz</th>
</tr>
</thead>
<tbody>
<tr>
<td>This time-decomposition-based method [6]</td>
<td>2.1 ms</td>
<td>102 µs</td>
<td>41 µs</td>
</tr>
<tr>
<td>Time-decomposition-based method [6]</td>
<td>294 ms</td>
<td>14.8 ms</td>
<td>5.9 ms</td>
</tr>
</tbody>
</table>

The test time using our new implementation is given by equation 8.

\[
T_3 = 2.25 \times \frac{H_{ideal}}{F_s} + \frac{4}{F_s}
\]

**equation 8**

The first part of the right side of equation 8 is the time for counting the occurrences of codes resulting in the calculation of the parameters. The second part of the right side of equation 8 is the time for the initialization. According to equation 7 and 8, the computation time of our new approach is almost three times faster than the space decomposition approach [13]. Table 1 presents the comparison between the three implementations for a 6-bit ADC and different sampling frequencies. The proposed implementation permits to drastically reduce the computation time compared to the time decomposition method [12], and is almost three times faster than the space decomposition approach [13]. Because [13] references its results in comparison with time-decomposition technique: “The hardware overhead is about four times of the hardware overhead of the time-decomposition technique”, later consideration about area overhead due to our new implementation will be only compared to the time-decomposition method.
The new approach has been implemented on a FPGA for different ADC resolutions. Table 2 provides the computation time for different ADC resolutions and sampling frequencies. Even for a high-resolution converter of 14 bits and a relatively low sampling frequency of 1MHz, the computation time remains below 1 second while the time decomposition approach necessitates more than 19,000 seconds.

Table 3 presents the silicon area overhead resulting from the FPGA implementation of our new method and from the integration of the method introduced in [12]. Both methods have been implemented for a 6-bit ADC. The LUT area is not included in the following figures as this study has been done in a context of LUT-based correction of ADC, and this area is the same for both methods.

The silicon area overhead is slightly the same for both methods. The new implementation requires 15% more cell area, but this logic area remain small compared to the LUT area. Indeed the LUT is made of $2^n$ word of n bits, where n is the ADC resolution.

Table 3:

<table>
<thead>
<tr>
<th>Implementation</th>
<th>Cells</th>
<th>Cell area (mm²)</th>
<th>Net area (mm²)</th>
</tr>
</thead>
<tbody>
<tr>
<td>This</td>
<td>257</td>
<td>0.028701</td>
<td>0.004455</td>
</tr>
<tr>
<td>Time-decomposition-based method [7]</td>
<td>223</td>
<td>0.025917</td>
<td>0.004464</td>
</tr>
</tbody>
</table>

5. Conclusion

LUT-based correction of ADC is effective. The current trend tends to integrate the procedure to compute the correction table. Our innovative computation technique has demonstrated the interest in re-using the embedded resources for measuring INL and computing the correction table. The silicon area overhead required for correction-table calculation is very small compared to the LUT area. The main advantage is the fast computation time. Indeed, as we are in the case of self-adaptive correction, the time for which the circuit is not operative because of the correction procedure should be as short as possible.

References


