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Extracting Bounded-level Modules from Deductive RDF Triplestores

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Abstract

We present a novel semantics for extracting bounded-level modules from RDF ontologies and databases augmented with safe inference rules, à la Datalog. Dealing with a recursive rule language poses challenging issues for defining the module semantics, and also makes module extraction algorithmically unsolvable in some cases. Our results include a set of module extraction algorithms compliant with the novel semantics. Experimental results show that the resulting framework is effective in extracting expressive modules from RDF datasets with formal guarantees, whilst controlling their succinctness.

Introduction

The Semantic Web consolidated a legacy of ontologies and databases today seen as reference systems for building new Semantic Web applications. To illustrate, consider a medical application for anatomy, whose goal is to showcase the structure of the human body, the most common pathologies and diseases, and the scientists that contributed to their study. A structural description of human anatomy can be drawn from FMA\(^1\) or My Corporis Fabrica (MyCF).\(^2\) A taxonomy of clinical terms about diseases can be extracted from SNOMED, \(^3\) while biographical informations about scientists implied in studies can be taken from DBPedia.\(^4\) These reference systems contain knowldege that can be reused to minimize the introduction of errors in the application. However, it is inconvenient to integrate in the application the whole datasets, as they contain complementary data and ontology axioms that are logically redundant. It is thus preferable to extract lightweight fragments of these reference systems - the modules - that are relevant for the application, and then to build on top of them.

Yet, more and more huge RDF datasets are flourishing in the Linked Data and some of them, like DBPedia or YAGO (Suchanek, Kasneci, and Weikum 2007), are increasingly reused in other more specialized datasets. RDF is a graph data model based on triples accepted as the W3C standard for Semantic Web data, with a simple ontology language. RDF Schema (RDFS). The W3C proposed OWL for writing expressive ontologies based on DL constructors. Whereas OWL is often seen as an extension of RDFS, this is not exactly the case. Both RDFS and the RDF query language (SPARQL) feature the possibility of accessing at the same time the ontology data and schema, by making variables ranging over classes or properties. This domain meta-modeling goes beyond the first-order setting typically considered in DL (De Giacomo, Lenzerini, and Rosati 2011). As a consequence, DL modularization frameworks are not applicable to popular RDF datasets like DBpedia or YAGO. Also, the clear separation between the ABox and the TBox made in DL to define the semantics of modules is not appropriate for RDF where facts and schema statements can be combined within a single RDF triplestore to accommodate heterogeneous knowledge from the Web. Another limit of the current approaches is that the existing semantics do not allow to limit the size of the extracted modules. As discussed in (Grau et al. 2008), the risk in practice is to output large portions of the initial ontologies, thus jeopardizing the gains of modularization.

Contributions. As a first contribution, we propose a unifying framework for RDF ontologies and databases we call deductive RDF triplestores, which is based on RDF triplestores equipped with safe Datalog inference rules. This rule language allows to capture in a uniform manner OWL constraints that are useful in practice, such as property transitivity or symmetry, but also domain-specific rules with practical relevance for users in many domains of interest.

The second and main contribution of the paper is a parametric semantics for bounded-level modules allowing to effectively control their size. We employ a notion of level of detail for modules in such a deductive setting. For example, a signature \(\text{subClassOf, partOf}^3\text{[eye]}\) limits the module-data extracted from a triplestore, by allowing to retrieve a description of all subclasses and subparts of the eye up to three levels.
Deductive RDF Triplestores

Data Along the lines of (Cali, Gottlob, and Pieris 2011) and (Libkin, Reutter, and Vrgoc 2013), we assume an infinite universe of constants \( \text{CONST} \), which forms the triplestore domain, and an infinite set of variables \( \text{VARS} \) used in rules. We range over constants by \( a, b, p, q \). We denote by \( x, y, z \) and \( 
\bar{x}, \bar{y}, \bar{z} \) variables and sequences of variables belonging to \( \text{VARS} \), respectively. A term, denoted by \( u, v \) is either a constant or a variable. We define RDF triplestores as relational databases over a schema restricted to a single ternary relation \( t(x, y, z) \). An atomic formula \( t \) is of the form \( t = t(u_1, u_2, u_3) \), with \( u_i \) a term. A triplestore \( D \) is a finite set of atoms of the form \( t(u_1, u_2, u_3) \) with \( u_i \in \text{CONST} \). Triplestores essentially represent RDF graph data over a relational vocabulary. A path \( p(u_0, u_n) = t(u_0, v_1, u_1), t(u_1, v_2, u_2), \ldots, t(u_{n-1}, v_n, u_n) \) is a sequence of atoms where each \( u_i, v_i \) are terms. The length of a path is the number of its atoms, here \( |p(u_0, u_n)| = n \).

Figure 1 presents an RDF triplestore, together with its graph version. The example is inspired by the MyCF ontology (Palombi et al. 2014), which classifies digital representation of human body parts, acquired by IRMs or tomographies, according to anatomical knowledge. For instance, the type edge connecting \( \text{irm142} \) with knee, corresponds to the triplestore atom \( t(\text{irm142}, \text{type}, \text{knee}) \), which is the standard RDF syntax for class membership.

Rules Deductive triplestores are equipped with safe Datalog rules. These are first-order logic formulas of the form \( \forall x, y. \phi(x, y) \rightarrow t(x) \) where \( x \) and \( y \) are sequences of variables belonging to \( \text{VARS} \), and \( t(x) \) and \( \phi(x, y) \) are an atom and a conjunction of atoms over the ternary relation \( t(x, y, z) \), constituting the \textit{head} and the body of a rule, respectively. We denote a rule by \( r \) and a set of rules by \( R \). To illustrate, the rules for class subsumption

\[
\begin{align*}
  r_1 & : t(x, \text{type}, y), t(y, \text{classOf}, z) \rightarrow t(x, \text{type}, z) \\
  r_2 & : t(x, \text{classOf}, y), t(y, \text{classOf}, z) \rightarrow t(x, \text{classOf}, z)
\end{align*}
\]

and, we exclude the presence of indirect recursion, in all cases where this involves non-recursive rules, like

\[
\begin{align*}
  r_4 & : t(x, \text{contains}, y) \rightarrow t(x, \text{partOf}, y) \\
  r_5 & : t(x, \text{partOf}, y), t(y, \text{partOf}, z) \rightarrow t(z, \text{contains}, x)
\end{align*}
\]

Datalog semantics is defined as the least fix-point of the immediate consequence operator (denoted by \( \vdash \)).

Definition 1 (Datalog Inference) The single-step Datalog inference, denoted by \( D, R \vdash_{\vdash} D' \), holds when \( D' = D \cup \{ \mu(\text{head}(r)) \} | \mu(\text{body}(r)) \in D \text{ and } r \in R \}, \) with \( \mu \) an homomorphism from variables to constants. Then, we write \( D, R \vdash D' \) when there exists a positive integer \( n \) such that \( D, R \vdash D_{n+1} \) for all \( 0 \leq n < m \), with \( D = D_0 \) and \( D' \subseteq D_n \). The triplestore \( D \) saturated with \( R \), is defined as \( \text{SAT}(D, R) = \{ t \in D' | D, R \vdash D' \} \).

We write \( D, R \vdash p(u_0, u_n) \) for the entailment of a path that holds if all path atoms are in \( \text{SAT}(D, R) \). Rule entailment, also referred as the immediate consequence operator for rules defines, by means of semantic conditions, when a Datalog rule \( r \) is entailed by a set \( R \).

Definition 2 (Rule Entailment) A rule \( r \) is entailed by a set \( R \), denoted by \( R \vdash r \), if for all triplestore \( D \) it holds that \( \text{SAT}(D, r) \subseteq \text{SAT}(D, R) \). A set \( R' \) is entailed from \( R \), denoted by \( R' \vdash r \) when \( R \vdash r \) for all \( r \in R' \).

A deductive RDF triplestore is a pair \( \langle D, R \rangle \) where \( D \) is a triplestore and \( R \) is a finite set of (possibly recursive) rules. Triplestore entailment, denoted by \( \langle D, R \rangle \vdash D' \), holds when \( D, R \vdash D' \) and \( R \vdash R' \).
Bounded-level Modules

A module is declared by means of a signature $\Sigma$ of the form $\Sigma = \langle p_1, \ldots, p_n \rangle^k[a]$ where the constants $p_1, \ldots, p_n$ represent the properties of interest of the module, the constant $a$ represents an object of interest of the module, and $k$ is a positive integer denoting the level of detail of the module. An example of module signature is $\langle \text{partOf} \rangle^3[\text{eye}]$. Intuitively, a module $M$ induced by a signature $\Sigma$ on a reference system $\langle D, R \rangle$ is a deductive triplestore $M = \langle D_M, R_M \rangle$ which is logically entailed by $\langle D, R \rangle$ and conforming to $\Sigma$, in the sense that all data and rule atoms employ the properties $p_1, \ldots, p_n$ only. Furthermore, to control the module size, the facts in $M$ are restricted to the paths rooted at the object of interest $a$, of length bounded by $k$.

We say that an atom conforms to $\Sigma$, denoted by $t(v_1, u, v_2) : \Sigma$, if $u$ is a property of $\Sigma$ or $u \in \text{VARS}$. A set of atoms $\Delta$ conforms to $\Sigma$ if all of its atoms do. Then, $\langle D, R \rangle$ conforms to $\Sigma$ if so do $D$ and $R$. In Figure 2(c) it holds that $D_3 : (\text{partOf}, \text{subClassOf})^2[\text{knee}]$. However, it does not hold that $D_3 : (\text{subClassOf})^2[\text{knee}]$.

Restricting the module paths is a way to effectively control the module size. Nevertheless, for the completeness of the module data, it is essential to guarantee that the module entails all of such bounded paths entailed by $\langle D, R \rangle$. In a deductive setting, adding new paths in the graph, defining properly $D_M$ becomes challenging.

First, we observe that to avoid incomplete modules, the paths of $D_M$ have to be drawn from $\text{SAT}(D, R)$. To see this, consider $D_2$ in Figure 2(a) and a rule inferring pairs of organs $(y, z)$ physically connected by a tendon

$$r_2 : t(x, \text{insertOn}, y), t(x, \text{insertOn}, z), t(x, \text{subClassOf}, \text{tendon}) \rightarrow t(y, \text{tendonConnected}, z)$$

A user interested in the organs directly and indirectly connected to the femur of this triplestore can declare the module signature $\Sigma_2 = (\text{tendonConnected})^2[\text{femur}]$. By restricting the module data $D_M$ to the paths in $D_2$ of length bounded by 2 that are rooted at femur and that use the property tendonConnected only, we get $D_M = \{ t(\text{femur}, \text{tendonConnected}, \text{gastroc.Muscle}) \}$. This dataset has however to be considered incomplete.

As shown in Figure 2(b), the rule $r_2$ entails on $D_2$ also the fact $t(\text{gastroc.Muscle}, \text{tendonConnected}, \text{knee})$. This forms a path of length two together with the original triple $t(\text{femur}, \text{tendonConnected}, \text{gastroc.Muscle})$, that should be included in $D_M$. The example illustrates clearly that $D_M$ depends from the rules in $R$.

However, taking into account all paths in $\text{SAT}(D, R)$ is not desirable for defining modules of bounded size. In some cases, the triples entailed by recursive rules may produce new edges in the data graph that behave like shortcuts between resources, thereby wasting the module parametricity. Consider $D_3$ in Figure 2(c) and the recursive rule $r_3$ defining the transitivity of partOf

$$r_3 : t(x, \text{partOf}, y), t(y, \text{partOf}, z) \rightarrow t(x, \text{partOf}, z)$$

The saturated triplestore $\text{SAT}(D_3, r_3)$ is depicted in Figure 2(d). It contains $t(\text{patella}, \text{partOf}, \text{knee})$ but also $t(\text{patella}, \text{partOf}, \text{leg})$ and $t(\text{patella}, \text{partOf}, \text{inferiorBody})$. More generally, it contains all triples of the form $t_b = t(\text{patella}, \text{partOf}, b)$ entailed by the transitivity of partOf. This means that if we take into account the recursive rule $r_3$ for defining the module paths, then all triples $t_b$ are likely to be part of the module induced by signature $(\text{partOf})^1[\text{knee}]$. This undermines the module parametricity because it retrieves all resources connected with knee regardless of the level of detail $k$.

Our solution to both keep into account implicit triples and make parametricity effective, is to define the module data as a subgraph of a partially-saturated triplestore obtained by applying non-recursive rules only, while fully delegating to the module the recursive rules. This leads to the following novel definition of module.

**Definition 3 (Module)** Let $\langle D, R \rangle$ be a deductive triplestore and $\Sigma = \langle p_1, \ldots, p_n \rangle^k[a]$ a signature. Then, $M = \langle D_M, R_M \rangle$ is a module for $\Sigma$ on $\langle D, R \rangle$ if

1. $\langle D_M, R_M \rangle : \Sigma$
2. $\langle D, R \rangle \vdash \langle D_M, R_M \rangle$
3. if $p_{(a,b)} : \Sigma$ and $|p_{(a,b)}| \leq k$ then
   a. $D, R, p \vdash p_{(a,b)}$ implies $D_M, R_M \vdash p_{(a,b)}$
   b. $D_M, R \vdash p_{(a,b)}$ implies $D_M, R_M \vdash p_{(a,b)}$

Point 1 and 2 of the definition state the well-formedness and the logical entailment of the modules, respectively. Point 3 is the crux of the definition. Property 3(a) says that every path rooted at $a$ of $k$-bounded length and conforming to $\Sigma$ that is entailed by the non-recursive rules of the reference system $p \vdash p_{(a,b)}$, must also be inferable by $M$. Property 3(b) enforces that the module rules $R_M$ infer the same paths conforming to $\Sigma$ as the whole set of rules $R$, but only when applied to the module data $D_M$. In contrast with the spirit of previous approaches (e.g., (Grau et al. 2008)), our definition does not enforce that every fact in the signature entailed by the reference triplestore also belongs to the module. Relaxing the module conditions in this way allows to control the module size, and cope with recursive rules.

To illustrate the definition, consider the triplestore $D_4$ of Figure 3(a) equipped with the rules below:

$$r_{3} : t(x, \text{hasFunction}, y) \rightarrow t(x, \text{participatesTo}, y)$$
$$r_{4} : t(x, \text{participatesTo}, y), t(y, \text{subClassOf}, z) \rightarrow t(x, \text{participatesTo}, z)$$
Figure 3: Triplestore and module examples

Figure 3(b) depicts \( \text{Sat}(D_4, \{r_3, r'_3\}) \). Consider now \( \Sigma_4 = (\text{participatesTo}, \text{subClassOf})^2[\text{knee}] \). A module \( M_4 \) for \( \Sigma_4 \) contains all paths rooted at knee of length at most 2, employing participatesTo and subClassOf only.

Note that if the recursive rule \( r'_4 \) is considered, then the triple \( t_1 = t(\text{knee}, \text{participatesTo}, \text{bodyPosture}) \) is included in the module dataset, which is not desirable. In contrast, \( t_2 = t(\text{knee}, \text{participatesTo}, \text{kneePosture}) \) is expected to be in a module for the signature \( \Sigma_4 \). A structure satisfying Definition 3 is \( M_4 = (D_{M_4}, R_{M_4}) \) with \( D_{M_4} \) depicted in Figure 3(c) and \( R_{M_4} = \{r'_4\} \). Note that \( t_2 \) is not explicitly in the module dataset \( D_{M_4} \) but can be inferred by \( r'_4 \) as shown in Figure 3(d).

Next, we present two algorithms for extracting module data and rules compliant with this novel semantics.

Extracting Module Data

The extraction of the module dataset can be done by leveraging on the evaluation of Datalog queries and implemented on top of existing engines. Given a module signature \( \Sigma = (p_1, \ldots, p_n)^k[a] \), the set of Datalog rules \( \Pi_\Sigma \) below computes all paths rooted at \( a \), of length bounded by \( k \), and built on the properties of interest of \( \Sigma \). It does so, in the extension of the relation \( m \), starting from a triplestore modeled with a single relation \( t \).

\[
\Pi_\Sigma = \begin{cases} 
  m_1^i(a, p_i, x) & \leftarrow t(a, p_i, x) \\
  m_2^j(x, p_i, y) & \leftarrow m_1^j(x_1, y_1, x), t(x, p_i, y) \\
  m(x, y, z) & \leftarrow m_2^j(x, y, z)
\end{cases}
\]

An instance of the rules is included for each \( i = 1 \ldots n \) and \( j = 1 \ldots k \). \( \Pi_\Sigma \) is a non-recursive set of rules of size \( O(nk) \) that can always be evaluated in at most \( k \) steps. Then, to infer all paths of bounded length entailed by non-recursive rules of a reference system, the set \( \Pi_\Sigma \) is evaluated together with \( R_{\text{NonRec}} \). As a result, the union \( \Pi_\Sigma \cup R_{\text{NonRec}} \) gives a non-recursive set of rules that can be evaluated in LOGSPACE data-complexity. The completeness of module data extraction follows from the completeness of Datalog query evaluation. Below, we write \( Q_m(D, \Pi_\Sigma \cup R_{\text{NonRec}}) \) for the answer set of the evaluation of the Datalog program.
Algorithm 1: MRE($N_{ToUnfold}$, $R_{ToApply}$, $\Sigma$)

1. for all $r_1 \in N_{ToUnfold}$ do
2.   if $r_1 : \Sigma$ then
3.     $R_M \leftarrow r_1$
4.     remove $r_1$ from $R_{ToApply}$
5.   else
6.     for all $r_2 \in R_{ToApply}$ s.t. $r_1 \neq r_2$ do
7.       for all $r \in RuleUnfolding(r_1, r_2)$ do
8.         if $r : \Sigma$ then
9.           $R_M \leftarrow r$
10.          $R_M \leftarrow$ MRE($\{r\}$, $R_{ToApply} \setminus \{r, r_2\}$, $\Sigma$)
11.       return $R_M$

not in $\Sigma$) is unfolded in a breadth-first fashion until no rule in $R_{ToApply}$ can be applied. All rules in $R$ are considered for unfolding ($R_{ToApply} = R$). Procedure RuleUnfolding($r_1, r_2$) progressively unfolds each subset of atoms in the body of $r_1$ that unify with the conclusion of $r_2$. For example, the three breadth-first unfoldings of $r_1 : t(x, p, y), t(x, p, z) \rightarrow t(x, p, y)$ with $r_2 : t(x, \text{partOf}, y) \rightarrow t(x, p, y)$ are:

- $r_3 : t(x, p, y), t(x, \text{partOf}, z) \rightarrow t(x, p, y)$
- $r_4 : t(x, \text{partOf}, y), t(x, p, z) \rightarrow t(x, p, y)$
- $r_5 : t(x, \text{partOf}, y), t(x, \text{partOf}, z) \rightarrow t(x, p, y)$

Note that a rule is never unfolded with itself by the algorithm (thus avoiding a depth-first fashion). The fact that $r_2$ used for the unfolding is discarded from $R_{ToApply}$ (line 10) ensures the termination of the extraction procedure, even in the presence of recursive rules.

Theorem 6 (Rule Extraction Algorithm) Let $R$ be a set of rules and $\Sigma$ a module signature. Algorithm MRE always terminates in $O(2^{|R|} \cdot |r|)$ and produces a set of rules $R_M$ conforming to $\Sigma$ such that for all $r : \Sigma$ it holds:

$R_M \vdash r$ implies $R \vdash r$ (Soundness)

Furthermore, when $R_{Rec} : \Sigma$ we also have:

$R \vdash r$ implies $R_M \vdash r$ (Completeness)

Algorithm MRE is sound, in the sense that it computes a set of rules entailed by $R$. Furthermore, for the case where all recursive rules in $R$ conform to $\Sigma$, the algorithm is also complete, in the sense that it produces a set of rules $R_M$ that entails all rules $R$ can entail on the properties of $\Sigma$. As a consequence, any dataset $D_M$ (computed as for Theorem 4) paired with $R_M$ constitutes a module meeting Definition 3, and in particular the point 3(b). If this condition does not hold, module extraction may be incomplete. To see this, consider again $(D, R)$ of Example 5 with $D = \{ t(a_1, p, a_2), t(a_2, p, a_3), t(a_3, p, a_2) \}$. Recall that $\Sigma = \{ p, q \} \setminus \{ a_1 \}$, and then notice that the recursive rule $r_2 : \Sigma$.

Yago2\* is the union of Yago2Taxonomy, Yago2Types and Yago2Facts datasets. We sampled classes and properties from these ontologies, and combined them to obtain a set of signatures used to run module extraction. We considered 2500 MyCF ontology classes combined with 20 subsets of its properties, of size 1-4. For the GO ontology (www.geneontology.org), we sampled 350 classes and 12 property sets (size 1-4). Since Yago knowledge is more diverse than a domain-specific ontology, to avoid empty modules we first selected three groups of properties that are frequently used together, and then subset them (size 2, 4, 6). We tested 100 Yago resources for each group. Finally, we made $k$ ranging over $\{1, 2, 3, 5, 10\}$.
Closest competitor approaches Relevant methods to our work are Traversal Views (Noy and Musen 2004) and Locality-based modules (Grau et al. 2008). Traversal Views (TV) compute a bounded-level view of an RDF database, in the same spirit as our approach. This method does not support inference rules, and it does not give any guarantee about extracted modules. In practice, in the presence of rules, a traversal view may miss relevant triples. Locality-Based (LB) module extraction computes a conservative extension of an ontology by checking logical conditions on its schema. In contrast with our method, it cannot modularize untyped RDF data and, because it enforces strong logical guarantees on a module, it cannot control a priori its size.

Results of module data extraction Figure 4(a-b) reports on the size of bounded-level modules, compared with those of TV and LB. The graphs show the average number of triples, for modules grouped by the same number of properties and $k$ value, in logarithmic scale. In Figure 4(c) we report the test on Yago2 with our approach, since LB does not support this RDF dataset.

As expected, the succinctness of bounded-level modules depends on $k$. The transitivity of the properties declared in the signature also has an impact. This is evident with Yago2 in Figure 4(c). Group 2 has properties inherently transitive (isLocatedIn, isConnectedWith) dominating for example (created, owns) in group 1 and (hasGender, isAffiliatedTo) in group 3. Hence, bounded-level modules can be very helpful to control the data succinctness with transitive properties.

Being TV unaware of rules, it may miss relevant data when implicit triples are not considered. We tested this claim, over the non-saturated MyCF ontology. Indeed, 42% (15072/35740) of the (non-empty) modules extracted by TV were missing relevant triples wrt our approach, as some sub-property rules were not evaluated. To overcome this limitation, we tested TV over the saturated MyCF. For concision, in Figure 4(a) we report only the minimal level of detail ($k = 1$). This already outlines a lower bound for the module size. As we can see, $k = 1$ already produces fairly larger modules than our approach. This is because of the MyCF rules for transitivity and property-chains. Increasing $k$ gives modules of size in the order of the saturated triplestore. The same discussion holds for GO in Figure 4(b). Hence, by having modules made of both data and rules our approach allows to retain succinctness wrt TV.

LB extraction for top-locality modules has been tested thanks to the available prototype www.cs.ox.ac.uk/isg/tools/ModuleExtractor/. For MyCF and GO, it outputs almost the whole ontology (Figures 4(a-b)). This is due to ontology axioms that cannot be ignored for the logical completeness of the method.

To conclude, the experiments outline the advantages of bounded-level modules as i) flexibility, to accomodate diverse Semantic Web datasets, and ii) succinctness, when dealing with transitive properties and rules.

Related Work and Conclusion
Module extraction has been extensively studied for DL. Related works employ basically two approaches that consist in inferring a conservative extension of an ontology (Grau et al. 2008; Konev et al. 2008), and forgetting non-interesting relations of ontology schemas (Grau and Motik 2012). Approximations and heuristics have also been devised to mitigate the complexity costs (Nortje, Britz, and Meyer 2013), and have been validated by statistical analysis (Vescovo et al. 2013). Differently from our approach, all of these techniques focus exclusively on the ontology schema, and do not permit to modularize Semantic Web datasets starting also from ontology instances or RDF facts. The closest work to our, at least in spirit, is (Noy and Musen 2004). However, it does not consider inference, which makes module extraction challenging. Answering Datalog queries over RDF has been investigated in (Libkin, Reutter, and Vrgoc 2013) and (Bry et al. 2008).

We presented a novel approach for the extraction of bounded-level modules from deductive RDF triplestores, which provides new means to reuse of Linked-Data datasets, and favorizes the development of Semantic Web applications. The key contribution of the work is a novel module semantics allowing to bound their size, and compliant module extraction algorithms. As shown by our experiments, the resulting framework allows to efficiently extract expressive modules from Semantic Web ontologies and databases with formal guarantees, whilst effectively controlling their succinctness. Future works include the study of module robustness to updates (Goasdoue and Rousset 2013), the extension towards Datalog± rules (Arenas, Gottlob, and Pieris 2014), and the connections between module extraction and recursive Datalog optimisation (Hillebrand et al. 1995).

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