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► **To cite this version:**

François Dross. Fractional Triangle Decompositions in Graphs with Large Minimum Degree. SIAM Journal on Discrete Mathematics, Society for Industrial and Applied Mathematics, 2016, 30 (1), pp.36-42. 10.1137/15M1014310 . lirmm-01336680

HAL Id: lirmm-01336680

<https://hal-lirmm.ccsd.cnrs.fr/lirmm-01336680>

Submitted on 23 Jun 2016

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Fractional triangle decompositions in graphs with large minimum degree

François Dross^a

^a*Université de Montpellier, LIRMM*
161 rue Ada, 34095 Montpellier Cedex 5, France
dross@lirmm.fr

Abstract

A triangle decomposition of a graph is a partition of its edges into triangles. A fractional triangle decomposition of a graph is an assignment of a non-negative weight to each of its triangles such that the sum of the weights of the triangles containing any given edge is one. We prove that every graph on n vertices with minimum degree at least $0.9n$ has a fractional triangle decomposition. This improves a result of Garaschuk that the same conclusion holds for graphs with minimum degree at least $0.956n$. Together with a recent result of Barber, Kühn, Lo and Osthus, this implies that for all $\epsilon > 0$, every large enough triangle divisible graph on n vertices with minimum degree at least $(0.9 + \epsilon)n$ admits a triangle decomposition.

1 Introduction

Decomposition and packing problems are central and classical problems in combinatorics, and, in particular, in design theory. Kirkman's theorem [8] from the middle of 19th century gives a necessary and sufficient condition on the existence of a Steiner triple system with a certain number of elements. In the language of graph theory, Kirkman's result asserts that every complete graph with an odd number of vertices and a number of edges divisible by three can be decomposed into triangles. Note that if a graph can be decomposed into triangles, then its vertex degrees are even and the total number of edges is divisible by three. Barber, Kühn, Lo and Osthus [2] showed that the same conclusion is true for large graphs satisfying these divisibility conditions if their minimum degree is not too far from the number of their vertices. In this short paper, we study the fractional variant of the problem and we use it, together with a result of Barber, Kühn, Lo and Osthus [2], to improve the best known bound.

Let us fix the terminology we are going to use. A *graph* is a pair of sets (V, E) such that elements of E are unordered pairs of elements of V . The elements of V are called *vertices* and the elements of E are called *edges*. We denote by uv (or vu) the edge with vertices u and v . We denote by $|G|$ the number of vertices of G . Two distinct vertices contained in the same edge are said to be *adjacent* or to be *neighbours*. Two edges that share a vertex are said to be *adjacent*. The *degree* of a vertex v is equal to the number

of neighbours of v . Let $\gcd(G)$ denote the greatest common divisor of the degrees of the vertices of G .

Two graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ are *isomorphic* if there exists a bijection b from V_1 to V_2 such that uv is an edge of G_1 if and only if $b(u)b(v)$ is an edge of G_2 for every two vertices u and v of G_1 . The *complete graph* K_k is the graph with k vertices all mutually adjacent. The graph K_3 is also called a *triangle*. A graph $G_1 = (V_1, E_1)$ is a *subgraph* of $G_2 = (V_2, E_2)$ if $V_1 \subseteq V_2$ and $E_1 \subseteq E_2$. The subgraphs of G_2 isomorphic to G_1 will be referred to as copies of G_1 .

Let H be a graph. An H -*decomposition* of a graph G is a set of subgraphs of G isomorphic to H that are edge disjoint such that each edge of G is contained in one of them. A graph is H -*decomposable* if it admits an H -decomposition. A K_3 -decomposition is also called a *triangle decomposition* and a graph is *triangle decomposable* if it is K_3 -decomposable. A graph G is H -*divisible* if $\gcd(G)$ is a multiple of $\gcd(H)$ and the number of edges of G is a multiple of the number of edges of H . It is easy to see that every H -decomposable graph is H -divisible. However, the converse is not true. For example a cycle on six vertices is K_3 -divisible, but not K_3 decomposable. As noted previously, Kirkman [8] proved that every K_3 -divisible complete graph is K_3 -decomposable. The fact that, for all H , every H -divisible complete graph is H -decomposable remained an open problem for over one hundred years before it was solved by Wilson [10].

The first generalisation to graphs that are near complete is due to Gustavsson [6]. He proved that for every graph H , there exist $n_0(H)$ and $\epsilon(H)$ such that every H -divisible graph with $n \geq n_0(H)$ vertices and minimum degree at least $(1 - \epsilon(H))n$ is H -decomposable. This has been generalised to hypergraphs in a recent result of Keevash [7]. The best that is known to date for a general graph H is due to Barber, Kühn, Lo and Osthus [2], who gave a way to turn a fractional decomposition into an exact one. They proved the following: Let F be a graph with minimum degree $\delta(F)$ and let H be a graph with chromatic number $\chi(H)$ and $e(H)$ edges. Let $C := \min\{9\chi(H)^2(\chi(H) - 1)^2/2, 10^4\chi(H)^{3/2}\}$ and let $t := \max\{C, 6e(H)\}$. Then for all $\epsilon > 0$, there exists an n_0 such that every H -divisible graph G on $n \geq n_0$ vertices with $\delta(G) \geq (1 - 1/t + \epsilon)n$ has an H -decomposition. For some particular classes of graphs, the exact asymptotic minimum degree threshold is known [2][11].

A *fractional H -decomposition* of a graph G is an assignment of non-negative weights to the copies of H in G such that for an edge e , the sum of the weights of the copies of H that contain e is equal to one. A graph is *fractionally H -decomposable* if it admits a fractional H -decomposition. A graph can be fractionally H -decomposable without being H -divisible. A fractional K_3 -decomposition is also called a *fractional triangle decomposition* and a graph is *fractionally triangle decomposable* if it is fractionally K_3 -decomposable. For all $r \geq 2$, Yuster [12] proved that every graph on n vertices with minimum degree at least $(1 - \frac{1}{9r^{10}})n$ is fractionally K_r -decomposable, and Dukes [3][4] proved that the same result holds for sufficiently large graphs on n vertices with minimum degree at least $(1 - \frac{2}{9r^2(r-1)^2})n$. Our paper already led to further research: Barber, Kühn, Lo, Montgomery and Osthus [1] proved, building on our combinatorial approach, that every graph on $n \geq 10^4r^3$ vertices with minimum degree at least $(1 - 1/10^4r^{3/2})n$ has a fractional K_r -decomposition.

In this paper we will focus on triangle decompositions of graphs with large minimum degree. The following conjecture is due to Nash-Williams [9]:

Conjecture 1 (Nash-Williams [9]). *Let G be a K_3 -divisible graph with n vertices and minimum degree at least $\frac{3}{4}n$. If n is large enough, then G is K_3 -decomposable.*

The best result towards a proof of Conjecture 1 is due to the combination of results of Garaschuk [5] and Barber, Kühn, Lo and Osthus [2].

Theorem 2 (Garaschuk [5], Barber, Kühn, Lo and Osthus [2]). *There exists an n_0 such that every K_3 -divisible graph G on $n \geq n_0$ vertices with minimum degree at least $0.956n$ is K_3 -decomposable.*

The proof of Theorem 2 relies on a result on fractional K_3 -decomposability, which we now state. The following appears as a conjecture in [5]. Note that for K_3 -divisible graphs, this is a consequence of Conjecture 1.

Conjecture 3 (Garaschuk [5]). *Let G be a graph with n vertices and minimum degree at least $\frac{3}{4}n$. If n is large enough, then G is fractionally K_3 -decomposable.*

The best known result towards proving Conjecture 1 was established by Garaschuk [5].

Theorem 4 (Garaschuk [5]). *Let G be a graph with n vertices and minimum degree at least $0.956n$. The graph G admits a fractional triangle decomposition.*

In this paper we use a different method to prove the following.

Theorem 5. *Every graph with n vertices and minimum degree at least $\frac{9}{10}n$ admits a fractional triangle decomposition.*

In [2], a particular case of Theorem 11.1 and Lemma 12.3 imply the following.

Theorem 6 (Barber, Kühn, Lo and Osthus [2]). *Suppose there exist n_0 and δ such that every graph on $n \geq n_0$ vertices with minimum degree at least δn is fractionally K_3 -decomposable. For all $\epsilon > 0$, there exists n_1 such that every K_3 -divisible graph on $n \geq n_1$ vertices with minimum degree at least $(\max(\delta, \frac{3}{4}) + \epsilon)n$ is K_3 -decomposable.*

Together with Theorem 6, our result improves Theorem 2.

Theorem 7. *Let $\epsilon > 0$. There exists an n_0 such that every K_3 -divisible graph on $n \geq n_0$ vertices with minimum degree at least $(\frac{9}{10} + \epsilon)n$ is K_3 -decomposable.*

2 Proof of Theorem 5

Let $\delta = \frac{1}{10}$. Fix a graph G with n vertices and minimum degree at least $(1 - \delta)n$. Suppose the graph G has at least one triangle with three vertices of degree at least $(1 - \delta)n + 2$. Let G' be the graph G where the edges of one such triangle are removed. Observe that G' has minimum degree at least $(1 - \delta)n$ and that if G' has a fractional triangle decomposition, then G has one too. By doing this operation several times, we can assume that G has no triangle with three vertices of degree at least $(1 - \delta)n + 2$. Let m be the number of edges of G .

Initially, we give the same weight w_Δ to every triangle such that the sum of the weights of the triangles is equal to $\frac{m}{3}$. We will modify the weights of the triangles to obtain a fractional triangle decomposition. We will do so in a way that the total sum of the weights is preserved.

We define the weight of an edge e to be the sum of the weights of the triangles that contain e . Given H a copy of K_4 in G , and two non-adjacent edges e_1 and e_2 in H , let

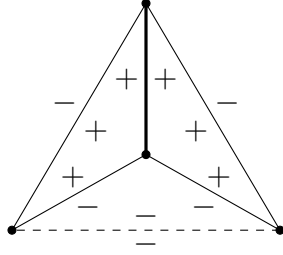


Figure 1: By removing some weight w from the two triangles containing the thick edge and adding w to the two triangles containing the dashed edge, we remove $2w$ from weight of the dashed edge and add $2w$ to the weight of the thick edge.

us call $(H, \{e_1, e_2\})$ a *rooted* K_4 of G . We will use the following procedure to modify the weights of the edges of a rooted K_4 of G :

Let $(H, \{e_1, e_2\})$ be a rooted K_4 of G . By removing a weight w from the two triangles of H that contain e_1 and adding the same weight w to each of the other two triangles (i.e. those that contain e_2), we transfer a weight of $2w$ from e_1 to e_2 . The weights of all the other edges of the graph remain unchanged (see Figure 1).

To prevent the weight of any triangle from becoming negative, we have to restrict how much weight we can transfer using the procedure above. If for some w we use the procedure to transmit a weight of $2w$ from an edge to another one, then any triangle's weight is lowered by at most w for triangles that are in the K_4 , and does not change for other triangles. Moreover, since every triangle contains a vertex with degree at most $(1 - \delta)n + 1$, any triangle is in at most $(1 - \delta)n - 1$ copies of K_4 , and thus in at most $3(1 - \delta)n - 3$ rooted K_4 (since for each K_4 there are three possible choices for the pair of edges). Since each triangle has an initial weight of w_Δ , if each rooted K_4 containing that triangle is used to transfer weight of at most $\frac{2w_\Delta}{3(1-\delta)n-3}$ between its labelled edges, then its final weight will be non-negative.

For each edge e , let T_e be the number of triangles of G containing e . We express redistributing the weights as a flow problem in an auxiliary graph, which is denoted by \widehat{G} . The vertices of \widehat{G} are the edges of G , plus two special vertices, called the *supersource* and the *supersink*. Two edges of G are adjacent as vertices in \widehat{G} if they form a pair in a rooted K_4 . The edge between them is set to have the capacity $c = \frac{2w_\Delta}{3(1-\delta)n-3}$. Let E_c be the set of these edges. If $T_e w_\Delta > 1$, then the vertex of \widehat{G} corresponding to e is joined to the supersource and the capacity of the corresponding edge of \widehat{G} is $T_e w_\Delta - 1$. Likewise, if $T_e w_\Delta < 1$, then the vertex of \widehat{G} corresponding to e is joined to the supersink and the capacity of the corresponding edge is $1 - T_e w_\Delta$. The vertices of G adjacent to the supersource are referred to as *sources* and those adjacent to the supersink as *sinks*. Let

$$M = \sum_{e \text{ source}} (T_e w_\Delta - 1) = \sum_{e \text{ sink}} (1 - T_e w_\Delta).$$

We will show that \widehat{G} has a flow of value M from the supersource to the supersink.

If \widehat{G} does not have a flow of value M , then it has a vertex cut (A_0, B_0) such that the supersource is contained in A_0 , the supersink in B_0 and the sum of the capacities of the edges from A_0 to B_0 is less than M . Let A be the edges of G corresponding to the vertices of A_0 and B the edges corresponding to the vertices of B_0 . Note that $|A| = |A_0| - 1$ and $|B| = |B_0| - 1$. Finally, let $k = |A|$, and observe that $|B| = m - k$.

Let T_A and T_B be the average T_e for e in A and in B respectively. Let $e = uv$ be an edge of G . Let W_e be the set of the vertices w such that uvw is a triangle. By the definition of T_e , $|W_e| = T_e$. Each vertex of W_e is non-adjacent to at most δn vertices of G , and thus is non-adjacent to at most δn vertices of W_e . So each vertex of W_e is adjacent to at least $T_e - \delta n$ vertices of W_e . Therefore e is in at least $\frac{T_e(T_e - \delta n)}{2}$ distinct copies of K_4 .

Let e be a vertex of A . It is adjacent to at least $\frac{T_e(T_e - \delta n)}{2} - k$ vertices of B . Therefore the cut contains at least

$$\sum_{e \in A} \left(\frac{T_e(T_e - \delta n)}{2} - k \right)$$

edges of E_c . Similarly, it contains at least

$$\sum_{e \in B} \left(\frac{T_e(T_e - \delta n)}{2} - (m - k) \right)$$

edges of E_c . Moreover, for each source e that is in B and each sink e that is in A , the cut contains the edge between e and the supersource or the supersink. Recall that the capacities of the edges of E_1 is $c = \frac{2w_\Delta}{3(1-\delta)n}$. Therefore the sum of the capacities of the edges of \hat{G} is at least

$$\sum_{e \in A} \left(\frac{T_e(T_e - \delta n)}{2} - k \right) c + \sum_{e \text{ source} \in B} (T_e w_\Delta - 1) + \sum_{e \text{ sink} \in A} (1 - T_e w_\Delta).$$

At the same time, it is also at least

$$\sum_{e \in B} \left(\frac{T_e(T_e - \delta n)}{2} - (m - k) \right) c + \sum_{e \text{ source} \in B} (T_e w_\Delta - 1) + \sum_{e \text{ sink} \in A} (1 - T_e w_\Delta).$$

Since the sum of the capacities of the edges in the considered cut is less than M , we get that

$$\sum_{e \in A} \left(\frac{T_e(T_e - \delta n)}{2} - k \right) c + \sum_{e \text{ source} \in B} (T_e w_\Delta - 1) + \sum_{e \text{ sink} \in A} (1 - T_e w_\Delta) < M \quad (1)$$

and

$$\sum_{e \in B} \left(\frac{T_e(T_e - \delta n)}{2} - (m - k) \right) c + \sum_{e \text{ source} \in B} (T_e w_\Delta - 1) + \sum_{e \text{ sink} \in A} (1 - T_e w_\Delta) < M. \quad (2)$$

The inequalities (1) and (2) can be rewritten using that

$$M = \sum_{e \text{ source}} (T_e w_\Delta - 1)$$

and

$$M = \sum_{e \text{ sink}} (1 - T_e w_\Delta)$$

respectively as follows.

$$\begin{aligned} \sum_{e \in A} (T_e(T_e - \delta n) - 2k) c - 2 \sum_{e \in A} (T_e w_\Delta - 1) &< 0 \\ \sum_{e \in B} (T_e(T_e - \delta n) - 2(m - k)) c - 2 \sum_{e \in B} (1 - T_e w_\Delta) &< 0 \end{aligned}$$

Since the summand is a convex function of T_e , we obtain the following.

$$(T_A(T_A - \delta n) - 2k)c - 2(T_A w_\Delta - 1) < 0 \quad (3)$$

$$(T_B(T_B - \delta n) - 2(m - k))c - 2(1 - T_B w_\Delta) < 0 \quad (4)$$

The inequality (3) implies that

$$T_A(T_A - \delta n) + \frac{2}{c}(1 - T_A w_\Delta) < 2k. \quad (5)$$

The inequality (4) implies that

$$2k < 2m - T_B(T_B - \delta n) + \frac{2}{c}(1 - T_B w_\Delta). \quad (6)$$

We now combine the inequalities (5) and (6) and we substitute $c = \frac{2w_\Delta}{3(1-\delta)n-3}$ to get the following.

$$T_A(T_A - \delta n) - (3n(1 - \delta) - 3)T_A < 2m - T_B(T_B - \delta n) - (3n(1 - \delta) - 3)T_B \quad (7)$$

Let e be an edge of G . Each end-vertex of e is non-adjacent to at most δn vertices of G . Hence, the edge e is contained in at least $n - 2\delta n$ triangles. Since e cannot be contained in more than n triangles, we get that $n - 2\delta n \leq T_e \leq n$. Consequently, we have $n - 2\delta n \leq T_A, T_B \leq n$.

A standard analytic argument shows that the left hand side of (7) is minimized when $T_A = n$ and the right hand side is maximized when $T_B = n - 2\delta n$. Consequently, it must hold that

$$n(n - \delta n) - 3(1 - \delta)n^2 + 3n < 2m - (n - 2\delta n)(n - 3\delta n) - 3n(1 - \delta)(n - 2\delta n) + 3n - 6\delta n.$$

Therefore

$$(2 - 12\delta + 12\delta^2)n^2 < 2m - 6\delta n. \quad (8)$$

Recall that we assumed that in G there is no triangle with three vertices of degree at least $(1 - \delta)n + 2$. Let V_b be the set of vertices of degree at least $(1 - \delta)n + 2$ in G , and let n_b be the number of vertices in V_b . Let us prove that $n_b \leq 2\delta n - 4$. Assume by contradiction that $n_b \geq 2\delta n - 3$. Since every vertex in V_b has at most $\delta n - 3$ non-neighbours in $V(G)$, every vertex in V_b has at most $\delta n - 3$ non-neighbours in V_b , and thus has at least $n_b - \delta n + 2$ neighbours in V_b . The graph induced by V_b is triangle-free, has $n_b \geq 2\delta n - 3$ vertices, and has minimum degree at least $n_b - \delta n + 2$. A triangle-free graph on k vertices has at most $\frac{k^2}{4}$ edges, thus it has minimum degree at most $\frac{k}{2}$. This implies that $n_b - \delta n + 2 \leq n_b/2$, i.e. $n_b/2 \leq \delta n - 2$, and thus $\delta n - 1.5 \leq \delta n - 2$, a contradiction. Therefore there are at most $2\delta n - 4$ vertices in V_b . We have

$$2m \leq (2\delta n - 4)n + ((1 - 2\delta)n + 4)((1 - \delta)n + 1) = (1 - \delta + 2\delta^2)n^2 + 4 + n - 6\delta n$$

and thus, with (8),

$$(2 - 12\delta + 12\delta^2)n^2 < (1 - \delta + 2\delta^2)n^2 + 4 + n - 12\delta n. \quad (9)$$

Assume $n < 20$. Since G has minimum degree at least $(1 - \delta)n = 0.9n$, the graph G is a complete graph. Then giving the same weight to every triangle leads to a fractional triangle decomposition of G .

Therefore we can assume that $n \geq 20$, and thus $4 + n - 12\delta n \leq 0$. We get from (9) that $(1 - 10\delta)(1 - \delta) = 1 - 11\delta + 10\delta^2 < 0$. Since $\delta = 0.1$, this leads to a contradiction. Therefore, there must exist a flow of value M in \widehat{G} , and hence, as described previously, the weights of the triangles can be adjusted in such a way that these weights now form a fractional decomposition of G . This finishes the proof of Theorem 5.

3 Conclusion

In this paper we proved that every graph on n vertices with minimum degree at least $\frac{9}{10}n$ is fractionally triangle decomposable. Together with a result of Barber, Kühn, Lo and Osthus [2], this implies that, for all $\epsilon > 0$, there exists a constant n_0 such that every triangle divisible graph on $n \geq n_0$ vertices with minimum degree at least $(0.9 + \epsilon)n$ is triangle decomposable.

4 Acknowledgements

This work was done during the author's visit to the group of Dan Král' at the University of Warwick; the visit was partially supported from the European Research Council under the European Union's Seventh Framework Programme (FP7/2007-2013)/ERC grant agreement no. 259385. I am deeply grateful to Tereza Klimošová for her helpful comments and to Dan Král' for his careful reading and many suggestions. Special thanks to Ben Barber and Richard Montgomery for pointing out that we can assume that every triangle has a vertex with low degree.

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