

# Scan chain encryption, a countermeasure against scan attacks

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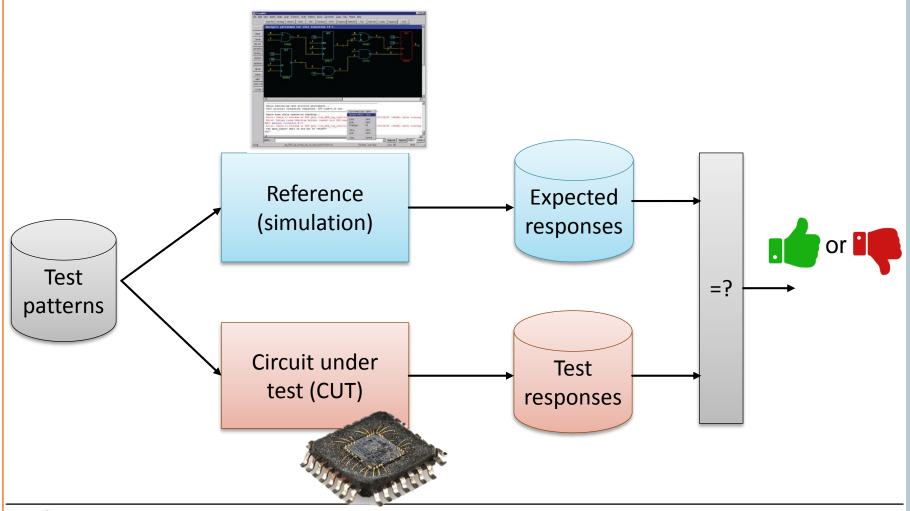




<u>Mathieu Da Silva</u>, Marie-Lise Flottes, Giorgio Di Natale, Bruno Rouzeyre

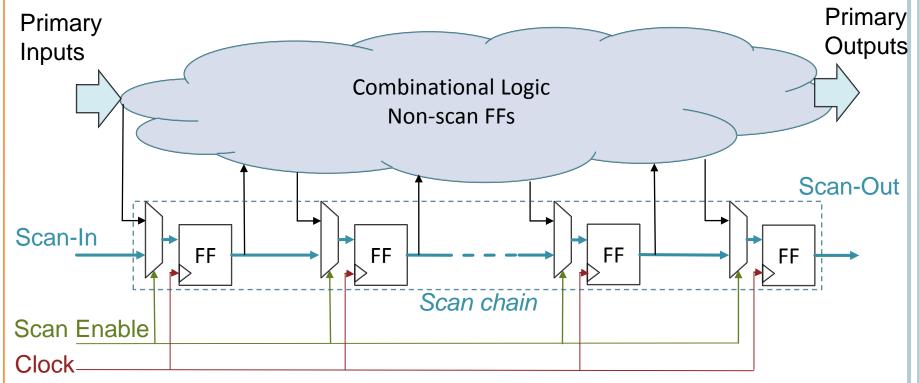
**PHISIC 2018** 

• Test of circuit is a mandatory step in IC production





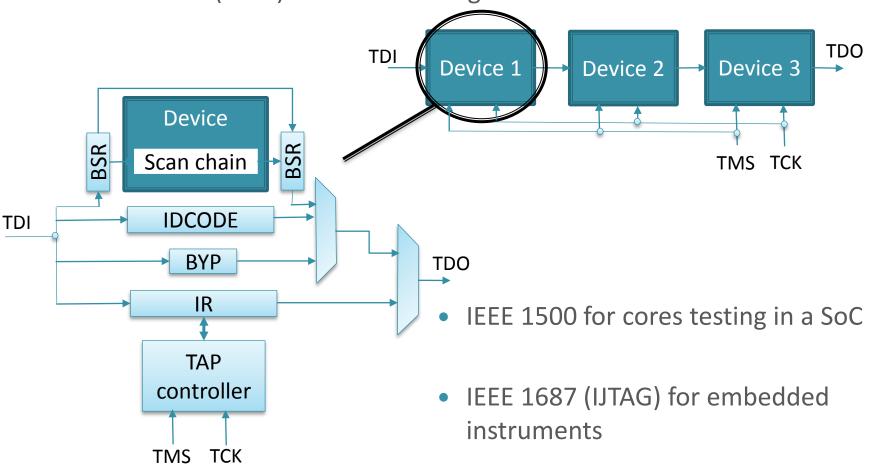
- Most popular method for Design-for-Test = Scan chains
  - Replace original FF by Scan FF connected serially together
  - Extra port « Scan-In » => controllability on internal states
  - Extra port « Scan-Out » => observability on internal states



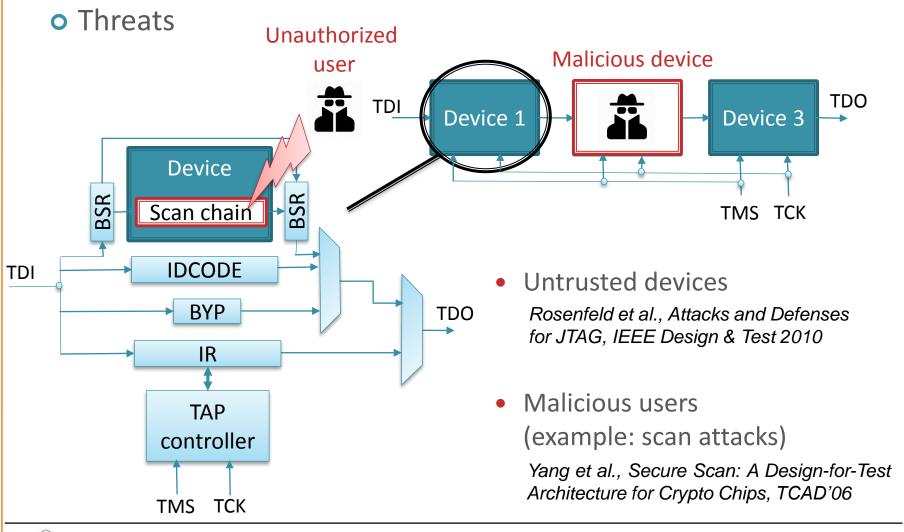


### Test standards

• IEEE 1149 (JTAG) for board testing









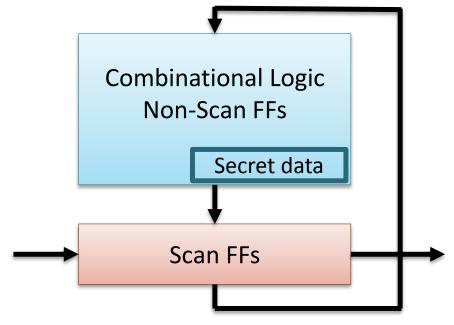
### **SUMMARY**

- 1) Scan attacks
- 2) A new countermeasure: Scan chain encryption
- 3) Implementation with block cipher
- 4) Implementation with stream cipher
- 5) Conclusion



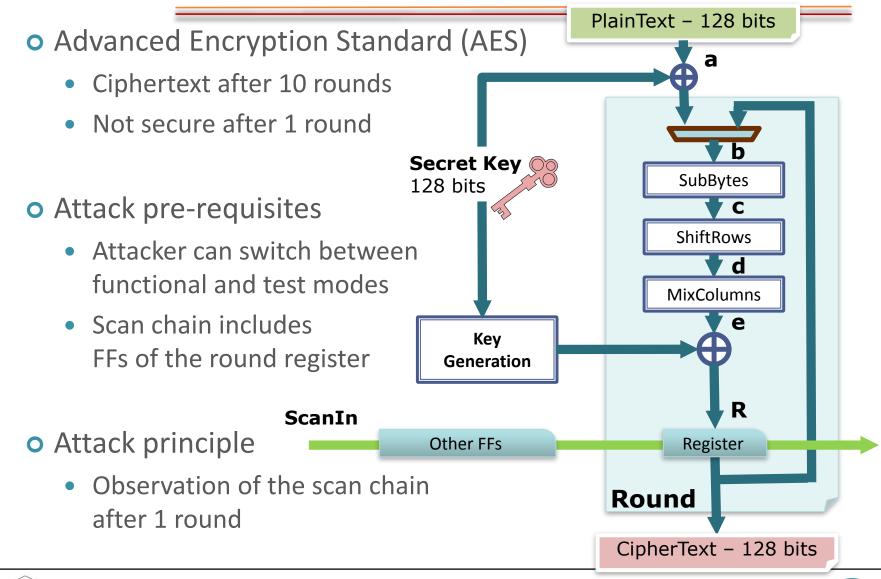
### SCAN ATTACK PRINCIPLE

- o Goal: Retrieve embedded secret data
- Exploit observability or controllability offered by scan chains
- Principle: switch between functional and scan modes
- Main target: secret key of crypto-processors (example: AES)



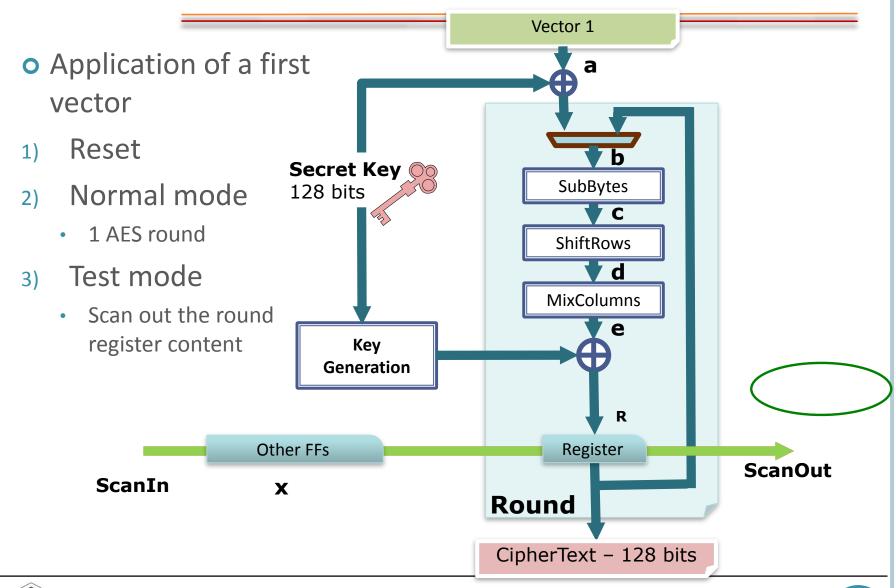


### SCAN ATTACK ON AES



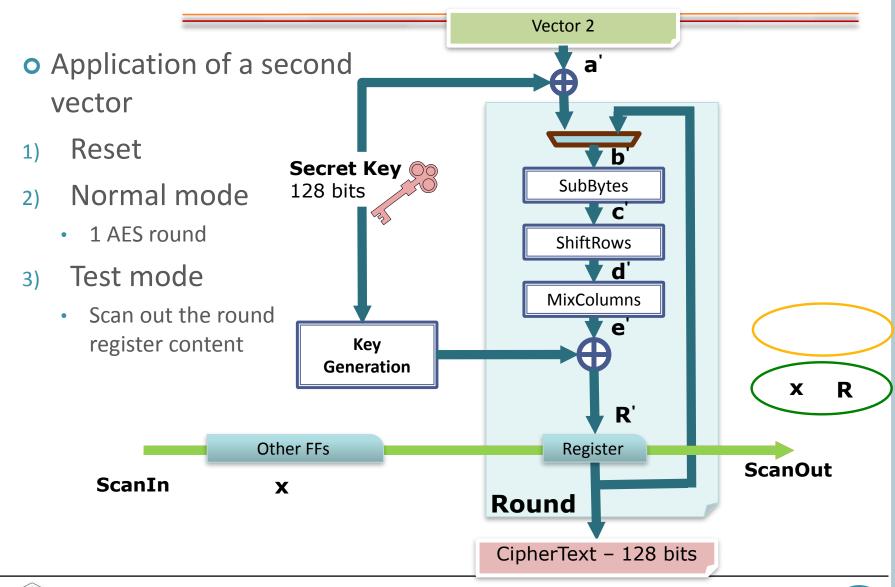


### DIFFERENTIAL ATTACK





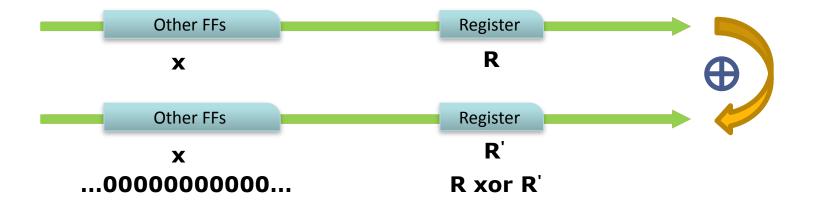
### DIFFERENTIAL ATTACK





### DIFFERENTIAL ATTACK

Hamming distance



- Attacker applies pairs of input values until hamming distance equal to specific values => key byte revealed
- On average, 32 trials
- ⇒ 512 trials to retrieve the whole 128-bit key



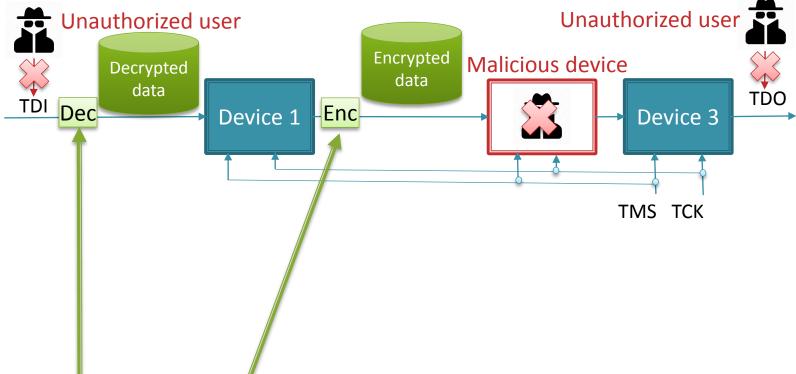
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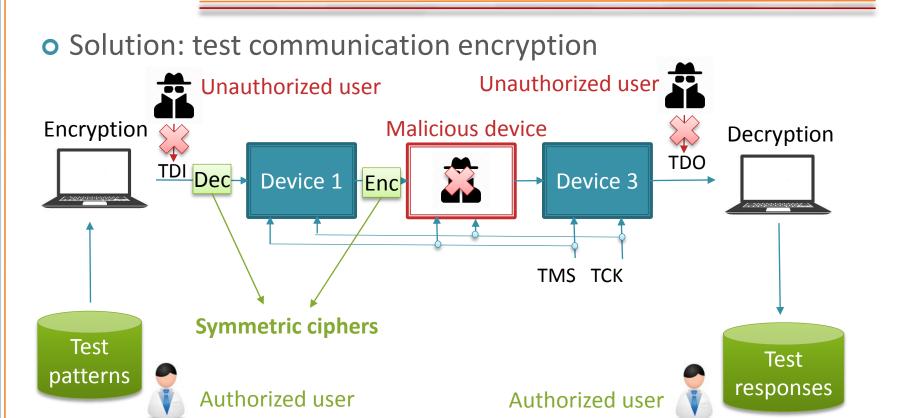
### SCAN CHAIN ENCRYPTION

Solution: test communication encryption



- Input decryption prevents sending desired test data
- Output encryption prevents reading plain test responses



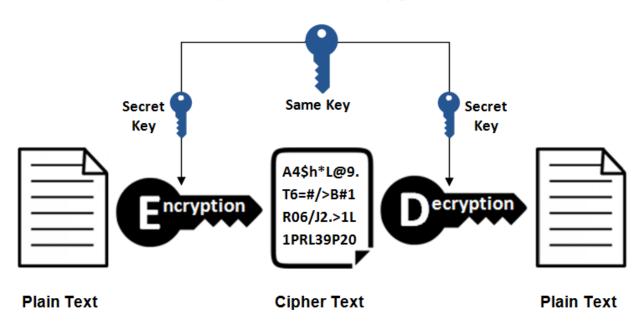


- Input decryption prevents sending desired test data
- Output encryption prevents reading plain test responses
- Test/debug only possible by authorized user knowing the secret key



### SYMMETRIC CIPHER

### **Symmetric Encryption**

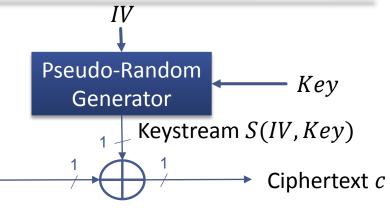


2 types of symmetric cipher: stream and block ciphers

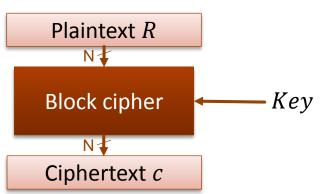


## STREAM CIPHER / BLOCK CIPHER

- Stream cipher encryption
  - Keystream XORed <u>bitwise</u>
    with the plaintext



- Block cipher encryption
  - Confusion and diffusion on a <u>block</u> of plaintext



- Preference for stream ciphers
  - "Naturally" adapted to serial test communication (JTAG, IEEE 1500, IJTAG)

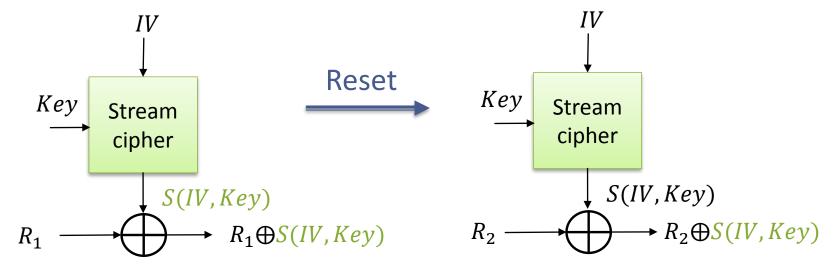
Plaintext R

- Smaller area footprint compared to block ciphers
- But ...



### TWO-TIMES PAD: STREAM CIPHER REQUIREMENT

 Two-times pad: same key and IV re-used => same keystream generated to encrypt different data



⇒ Possible to carry out attacks if requirement is not fit

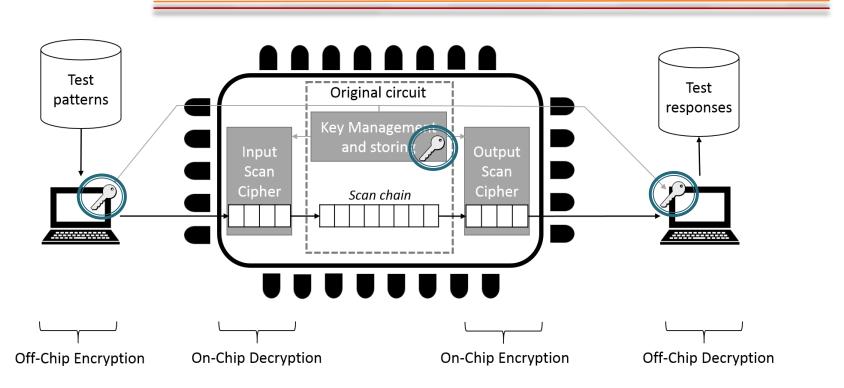
$$R1 \oplus S(W, Key) \oplus R2 \oplus S(W, Key)$$

⇒ Solution: *IV* generated randomly at each circuit reset

$$R1 \oplus S(IV_1, Key) \oplus R2 \oplus S'(IV_2, Key)$$



### **BASIC SCHEME**



- Assumption: original circuit embedded a crypto-core with its key management and storing
- Scan chain encryption solution shares the key management and storing already implemented



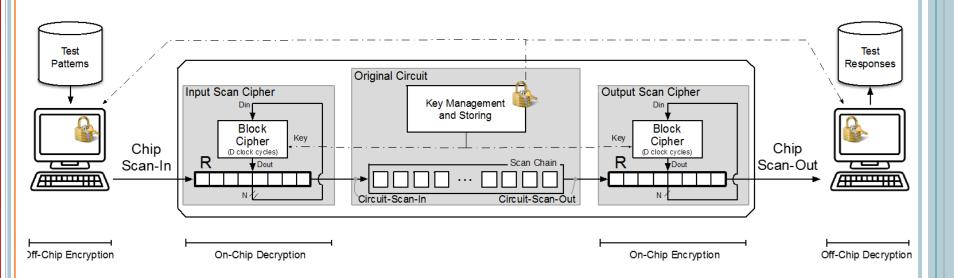
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### BLOCK CIPHER-BASED SCAN ENCRYPTION

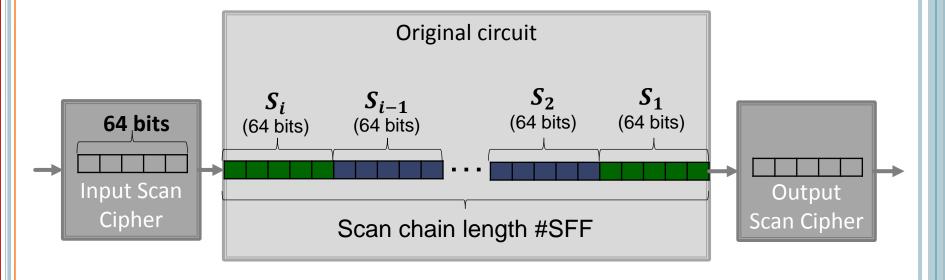
- Implementation on scan chain with 2 PRESENT block ciphers:
  - Lightweight (1 PRESENT = 2 139 GE)
  - Encryption by 64-bits block size





### MODE OF OPERATIONS

64 bits encrypted every 32 clock cycles

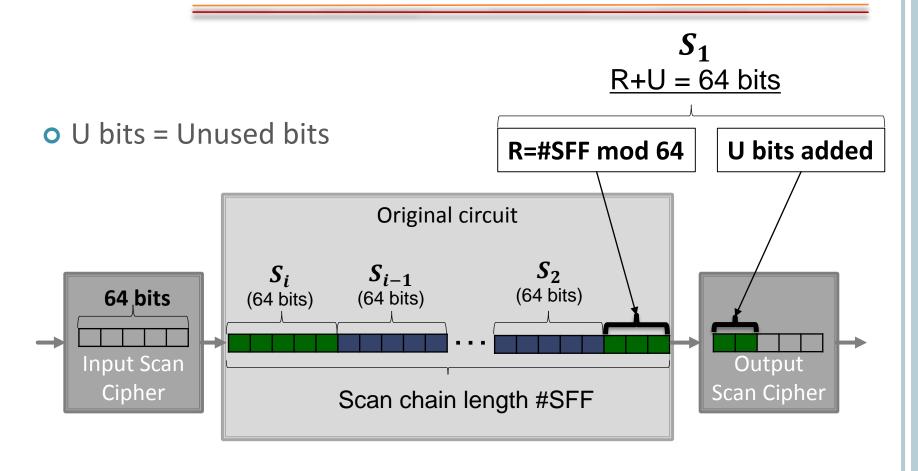


- $\Rightarrow$  #SFF = Px64
- ⇒ No test time overhead on each pattern





### Mode of operations



- $\Rightarrow$  #SFF = Px64 + R
- ⇒ Loss of U clock cycles per pattern





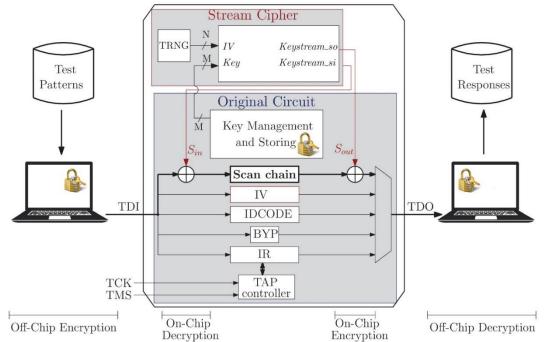
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### STREAM CIPHER-BASED SCAN ENCRYPTION

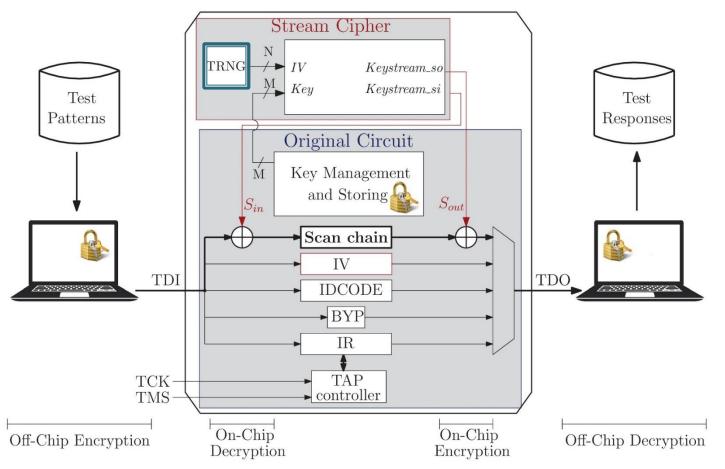
- Implementation on JTAG:
  - 1 TRIVIUM stream cipher (2 016 GE)
  - TRNG to generate random IV
  - New instruction GetIV with a test data register IV



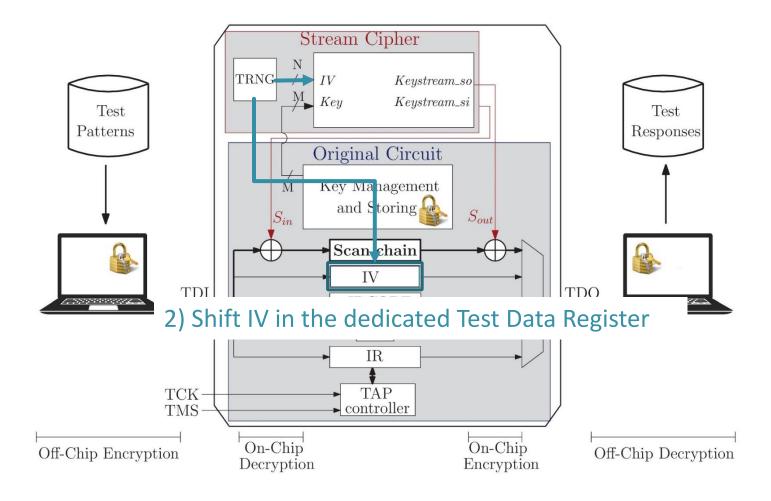
Mode of operations in 2 phases: initialization and encryption



### 1) TRNG initialization: reach sufficient entropy to generate random number

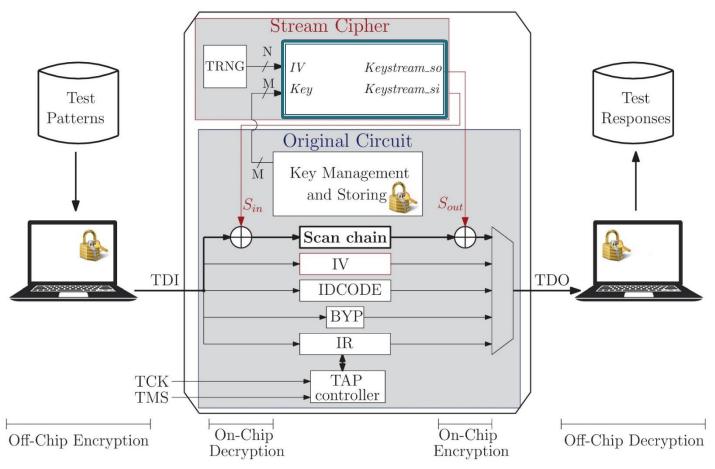




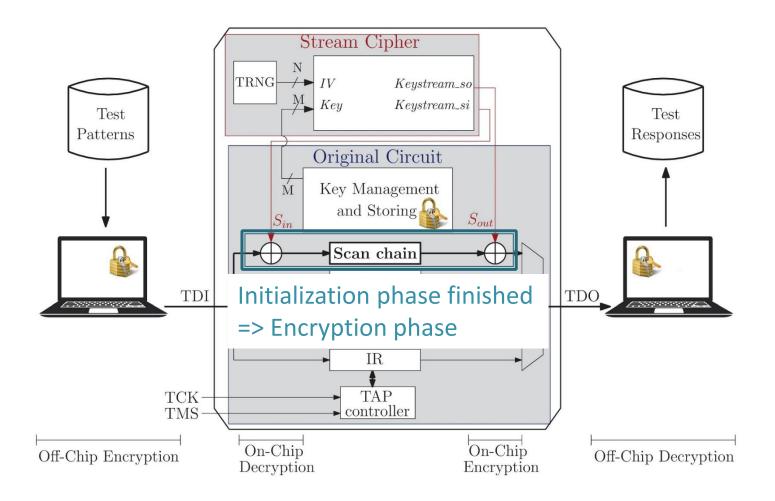




### 3) Stream cipher setup



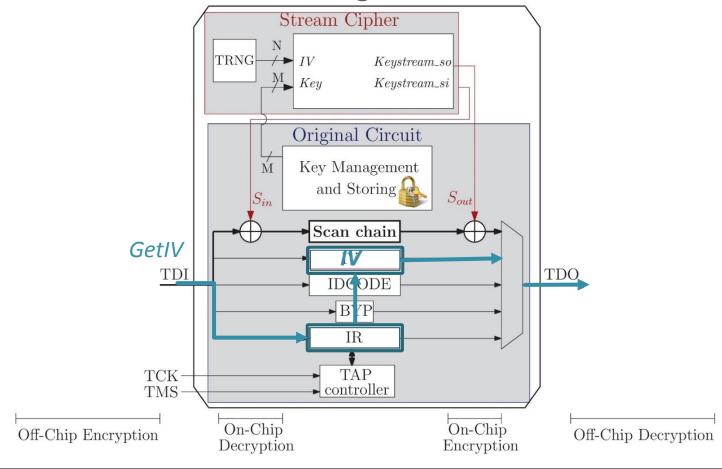






### **ENCRYPTION PHASE**

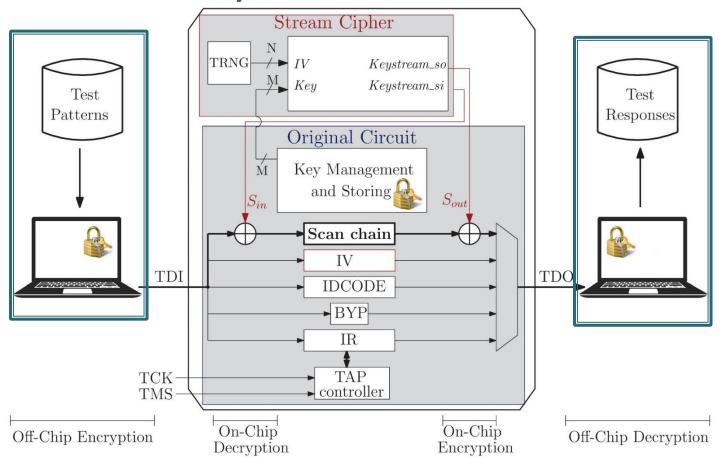
- Send GETIV instruction
- ⇒ Shift the content of the IV register out the circuit





### **ENCRYPTION PHASE**

 User can encrypt and decrypt test data with the obtained IV and the shared secret key





### TIME FOR THE INITIALIZATION PROCESS

- $\circ$   $T_{TRNG\ init}$  to initialize the TRNG
- 80 clock cycles to shift the IV in the register
- 1 152 clock cycles for the stream cipher setup

Original circuit	Triple-DES	Pipelined AES-128	Pipelined AES-256	RSA 1024	LEON3
Test time* (clock cycles)	687 101	1 944 877	4 559 845	39 405 239	11 612 051
Test time overhead					
Block-based solution (%)	+0.31	+0.81	+0.006	+0.33	+0.004
Stream-based solution (%)**	+0.18	+0.06	+0.03	+0.003	+0.01

<sup>\*:</sup> Test time considered for a fault coverage of 100%, except for LEON3 where it reaches 70%

<sup>\*\*:</sup> test time overhead without the initialization of the TRNG



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### COMPARISON BETWEEN BOTH SOLUTIONS

	Block cipher-based solution (PRESENT)	Stream cipher-based solution (TRIVIUM)				
Security						
- Scan attacks	Protected	Protected (two times pad not possible)				
- Malicious core	Protected	Protected				
Cost						
- Area	10 658.96 μm²	5 408.52 μm² (+ 31 200 μm² for TRNG)				
- Test time	Depends on the scan length (multiple or not of the block size)	Clock cycles required for the initialization phase				
Integration						
- Diagnosis & debug	Still possible in-field					
- Key management	Re-use key management already implemented					
- Integration in test daisy-chain	Possible issue with the padding of test data	No issue				



# Thankyou

### **ACKNOWLEDGEMENTS**

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Partners











