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Hitting minors on bounded treewidth graphs.

III. Lower bounds*

Julien Baste^{†‡§} Ignasi Sau[¶] Dimitrios M. Thilikos[¶]

Abstract

For a finite collection of graphs \mathcal{F} , the \mathcal{F} -M-DELETION problem consists in, given a graph G and an integer k , decide whether there exists $S \subseteq V(G)$ with $|S| \leq k$ such that $G \setminus S$ does not contain any of the graphs in \mathcal{F} as a minor. We are interested in the parameterized complexity of \mathcal{F} -M-DELETION when the parameter is the treewidth of G , denoted by tw . Our objective is to determine, for a fixed \mathcal{F} , the smallest function $f_{\mathcal{F}}$ such that \mathcal{F} -M-DELETION can be solved in time $f_{\mathcal{F}}(\text{tw}) \cdot n^{\mathcal{O}(1)}$ on n -vertex graphs. We provide lower bounds under the ETH on $f_{\mathcal{F}}$ for several collections \mathcal{F} . We first prove that for any \mathcal{F} containing connected graphs of size at least two, $f_{\mathcal{F}}(\text{tw}) = 2^{\Omega(\text{tw})}$, even if the input graph G is planar. Our main contribution consists of superexponential lower bounds for a number of collections \mathcal{F} , inspired by a reduction of Bonnet et al. [IPEC, 2017]. In particular, we prove that when \mathcal{F} contains a single connected graph H that is either P_5 or is not a minor of the banner (that is, the graph consisting of a C_4 plus a pendent edge), then $f_{\mathcal{F}}(\text{tw}) = 2^{\Omega(\text{tw} \cdot \log \text{tw})}$. This is the third of a series of articles on this topic, and the lower bounds given here together with the algorithms given in the first two articles yield a dichotomy on the complexity of $\{H\}$ -M-DELETION in terms of H , when H is a connected planar graph with $|V(H)| \geq 2$. Namely, $f_{\{H\}}(\text{tw}) = 2^{\Theta(\text{tw})}$ if H is a minor of the banner that is different from P_5 , and $f_{\{H\}}(\text{tw}) = 2^{\Theta(\text{tw} \cdot \log \text{tw})}$ otherwise.

Keywords: parameterized complexity; graph minors; treewidth; hitting minors; topological minors; dynamic programming; Exponential Time Hypothesis.

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The results of this article are permanently available at <https://arxiv.org/abs/1704.07284>. Extended abstracts containing some of the results of this article appeared in the *Proc. of the 12th International Symposium on Parameterized and Exact Computation (IPEC 2017)* [3], in the *Proc. of the 13th International Symposium on Parameterized and Exact Computation (IPEC 2018)* [4], and in the *Proc. of the 31st Annual ACM-SIAM Symposium on Discrete Algorithms (SODA 2020)* [7]. Work supported by French projects DEMOGRAPH (ANR-16-CE40-0028) and ESIGMA (ANR-17-CE23-0010).

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1 Introduction

Let \mathcal{F} be a finite non-empty collection of non-empty graphs. In the \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) problem, we are given a graph G and an integer k , and the objective is to decide whether there exists a set $S \subseteq V(G)$ with $|S| \leq k$ such that $G \setminus S$ does not contain any of the graphs in \mathcal{F} as a minor (resp. topological minor). Instantiations of these problems correspond to several well-studied problems. For instance, the cases $\mathcal{F} = \{K_2\}$, $\mathcal{F} = \{K_3\}$, and $\mathcal{F} = \{K_5, K_{3,3}\}$ of \mathcal{F} -M-DELETION (or \mathcal{F} -TM-DELETION) correspond to VERTEX COVER, FEEDBACK VERTEX SET, and VERTEX PLANARIZATION, respectively.

We are interested in the parameterized complexity of both problems when the parameter is the treewidth of the input graph. By Courcelle’s theorem [13], \mathcal{F} -M-DELETION \mathcal{F} -TM-DELETION can be solved in time $f(\text{tw}) \cdot n^{\mathcal{O}(1)}$ on n -vertex graphs with treewidth at most tw , where f is some computable function. Our objective is to determine, for a fixed collection \mathcal{F} , which is the *smallest* such function f that one can (asymptotically) hope for, subject to reasonable complexity assumptions.

This line of research has attracted some attention in the parameterized complexity community during the last years. For instance, VERTEX COVER is easily solvable in time $2^{\mathcal{O}(\text{tw})} \cdot n^{\mathcal{O}(1)}$, called *single-exponential*, by standard dynamic-programming techniques, and no algorithm with running time $2^{o(\text{tw})} \cdot n^{\mathcal{O}(1)}$ exists, unless the Exponential Time Hypothesis (ETH)¹ fails [20].

For FEEDBACK VERTEX SET, standard dynamic programming techniques give a running time of $2^{\mathcal{O}(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$, while the lower bound under the ETH [20] is again $2^{o(\text{tw})} \cdot n^{\mathcal{O}(1)}$. This gap remained open for a while, until Cygan et al. [15] presented an optimal algorithm running in time $2^{\mathcal{O}(\text{tw})} \cdot n^{\mathcal{O}(1)}$, introducing the celebrated *Cut&Count* technique. This article triggered several other techniques to obtain single-exponential algorithms for so-called *connectivity problems* on graphs of bounded treewidth, mostly based on algebraic tools [8, 18].

Concerning VERTEX PLANARIZATION, Jansen et al. [21] presented an algorithm of time $2^{\mathcal{O}(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ as a subroutine in an FPT-algorithm parameterized by k . Marcin Pilipczuk [25] proved that this running time is *optimal* under the ETH. This lower bound was achieved by using the framework introduced by Lokshtanov et al. [23, 24] for obtaining superexponential lower bounds (namely, of the form $2^{\Omega(k \cdot \log k)} \cdot n^{\mathcal{O}(1)}$, in particular for problems parameterized by treewidth), which has proved very successful in recent years [2, 15, 25]. We also use this framework in the current article.

Our results and techniques. We present lower bounds under the ETH for \mathcal{F} -M-

¹The ETH states that 3-SAT on n variables cannot be solved in time $2^{o(n)}$; see [20] for more details.

DELETION and \mathcal{F} -TM-DELETION parameterized by treewidth, several of them being tight. We first prove that for any connected² \mathcal{F} , neither \mathcal{F} -M-DELETION nor \mathcal{F} -TM-DELETION can be solved in time $2^{o(\text{tw})} \cdot n^{\mathcal{O}(1)}$, even if the input graph G is planar (cf. Theorem 2 and Corollary 1). The main contribution of this article consists of super-exponential lower bounds for a number of collections \mathcal{F} , which we proceed to describe. Let \mathcal{C} be the set of all connected graphs that contain a block (i.e., a biconnected component) with at least five edges, let \mathcal{Q} be the set containing P_5 and all connected graphs that are *not* minors of the banner (that is, the graph consisting of a C_4 plus a pendent edge), and let $\mathcal{S} = \{K_{1,i} \mid i \geq 4\}$. We prove that, assuming the ETH,

- for every finite non-empty subset $\mathcal{F} \subseteq \mathcal{C}$, neither \mathcal{F} -M-DELETION nor \mathcal{F} -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ (cf. Theorem 3),
- for every $H \in \mathcal{Q}$, $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ (cf. Theorem 4), and
- for every $H \in \mathcal{Q} \setminus \mathcal{S}$, $\{H\}$ -TM-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ (cf. Theorem 5).

The general lower bound of $2^{o(\text{tw})} \cdot n^{\mathcal{O}(1)}$ for connected collections is based on a simple reduction from (PLANAR) VERTEX COVER. The superexponential lower bounds, namely $2^{o(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$, are based on the ideas presented by Bonnet et al. [11] for generalized feedback vertex set problems. We provide two hardness results that apply to different families of collections \mathcal{F} , both based on a general framework described in Section 4.1, consisting of a reduction from the $k \times k$ PERMUTATION INDEPENDENT SET problem introduced by Lokshtanov et al. [24]. Namely, we prove, in Theorem 3 (applying to both the minor and topological minor versions), the lower bound when \mathcal{F} is any finite non-empty subset of all connected graphs that contain a block with at least five edges. We then prove, in Theorem 4 (for minors) and Theorem 5 (for topological minors), the lower bound when \mathcal{F} contains a single graph H that is either P_5 or is not a minor of the banner, with the exception of $K_{1,i}$ mentioned above for the topological minor version. The proofs of the latter two theorems are considerably longer, as we need to distinguish several cases according to certain properties of the graph H (cf. Lemma 8 up to Lemma 16).

We would like to mention that in previous versions of this article (in particular, in the conference version presented in [3]), we presented another family of reductions inspired by a reduction of Pilipczuk [25] for VERTEX PLANARIZATION, that is, for $\mathcal{F} = \{K_5, K_{3,3}\}$. Afterwards, we found a more general unifying reduction along the lines of Bonnet et

²A *connected* collection \mathcal{F} is a collection containing only connected graphs of size at least two.

al. [11], which is the one we present here. This reduction generalizes the hardness results presented in [3] and in [4] (also inspired by [11], and that can be seen as a weaker version of the current reduction), as well as the lower bound of Pilipczuk [25], which is a corollary of one of our hardness results, namely Theorem 3.

Results in other articles of the series and discussion. In the first article of this series [5], we show, among other results, that for every connected \mathcal{F} containing at least one planar graph (resp. subcubic planar graph), \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) can be solved in time $2^{\mathcal{O}(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$. In the second article of this series [6], we provide single-exponential algorithms for $\{H\}$ -M-DELETION for all the graphs H for which the superexponential lower bounds given in this article do not apply, namely those on the left of Figure 1: P_3 , P_4 , C_4 , the **claw**, the **paw**, the **chair** (sometimes also called *fork* in the literature), and the **banner**. Note that the cases $H = P_2$ [14, 20], $H = P_3$ [1, 26], and $H = C_3$ [8, 15] were already known (nevertheless, for completeness we provide a simple algorithm when $H = P_3$).

The lower bounds presented in this article, together with the algorithms given in [5, 6], cover all the cases of \mathcal{F} -M-DELETION where \mathcal{F} contains a single connected planar graph, as discussed in Section 4. Namely, we obtain the following theorem.

Theorem 1. *Let H be a connected planar graph of size at least two. Then the $\{H\}$ -M-DELETION problem can be solved in time*

- $2^{\Theta(\text{tw})} \cdot n^{\mathcal{O}(1)}$, if H is a minor of the **banner** that is different from P_5 , and
- $2^{\Theta(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$, otherwise.

Note that the graphs H described in the first item can be equivalently characterized as those that can be obtained by the **chair** or the **banner** by contracting edges. In the above statement, we use the Θ -notation to indicate that these algorithms are *optimal* under the ETH. This dichotomy is depicted in Figure 1, containing all connected planar graphs H with $2 \leq |V(H)| \leq 5$; note that if $|V(H)| \geq 6$, then H is not a minor of the **banner**, and therefore the second item above applies. See [6] for a discussion about the role played by the **banner** in this dichotomy.

The single-exponential algorithms given in [6] also apply to $\{H\}$ -TM-DELETION for which, in addition, we provide a single-exponential algorithm when $H = K_{1,i}$ for every $i \geq 1$. By Theorem 4, a single exponential-algorithm for $\{K_{1,i}\}$ -M-DELETION when $i \geq 4$ is unlikely to exist. To the best of our knowledge, this is the first example of a collection \mathcal{F} for which the complexity of \mathcal{F} -M-DELETION and \mathcal{F} -TM-DELETION differ.

Organization of the paper. In Section 2 we give some preliminaries. In Section 3 we present the single-exponential lower bound for any connected \mathcal{F} , and in Section 4 the

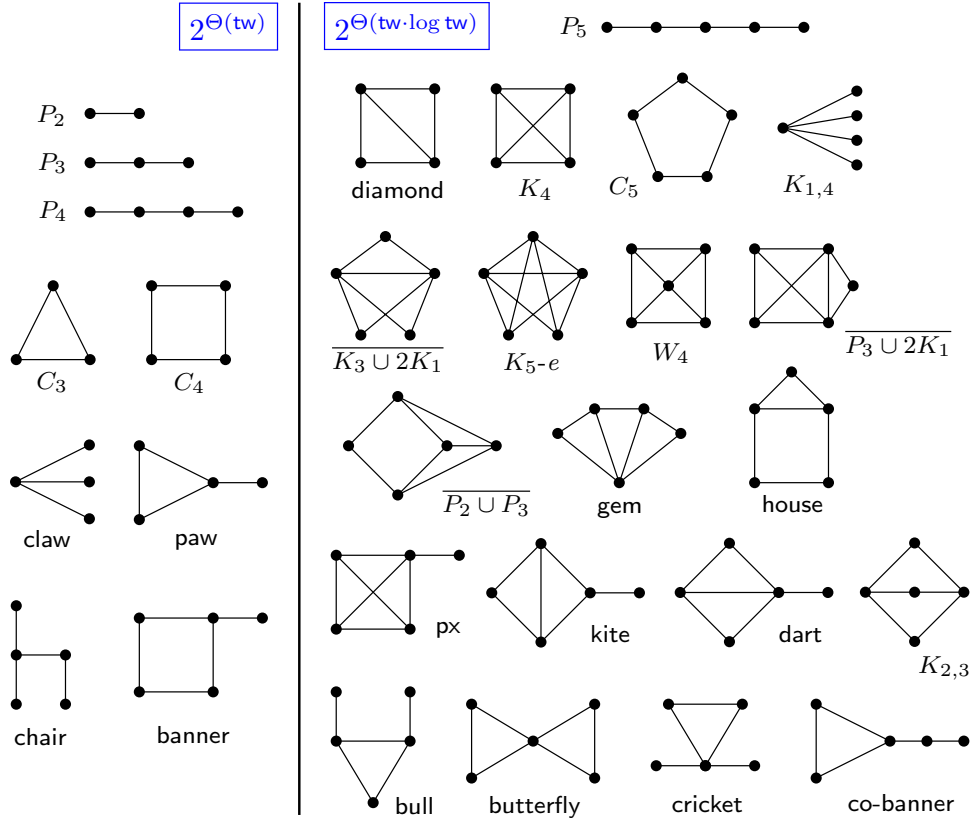


Figure 1: Classification of the complexity of $\{H\}$ -M-DELETION for all connected simple planar graphs H with $2 \leq |V(H)| \leq 5$: for the nine graphs on the left (resp. 20 graphs on the right, and all the larger ones), the problem is solvable in time $2^{\Theta(\text{tw})} \cdot n^{\mathcal{O}(1)}$ (resp. $2^{\Theta(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$). For $\{H\}$ -TM-DELETION, $K_{1,4}$ should be on the left. This figure also appears in [6].

superexponential lower bounds. We conclude the article in Section 5 with some open questions for further research.

2 Preliminaries

In this section we provide some preliminaries to be used in the following sections.

Sets, integers, and functions. We denote by \mathbb{N} the set of every non-negative integer and we set $\mathbb{N}^+ = \mathbb{N} \setminus \{0\}$. Given two integers p and q , the set $[p, q]$ refers to the set of every integer r such that $p \leq r \leq q$. Moreover, for each integer $p \geq 1$, we set

$\mathbb{N}_{\geq p} = \mathbb{N} \setminus [0, p-1]$. In the set $[1, k] \times [1, k]$, a *row* is a set $\{i\} \times [1, k]$ and a *column* is a set $[1, k] \times \{i\}$ for some $i \in [1, k]$.

We use \emptyset to denote the empty set and \emptyset to denote the empty function, i.e., the unique subset of $\emptyset \times \emptyset$. Given a function $f : A \rightarrow B$ and a set S , we define $f|_S = \{(x, f(x)) \mid x \in S \cap A\}$. Moreover if $S \subseteq A$, we set $f(S) = \bigcup_{s \in S} \{f(s)\}$. Given a set S , we denote by $\binom{S}{2}$ the set containing every subset of S that has cardinality two. We also denote by 2^S the set of all the subsets of S . If \mathcal{S} is a collection of objects where the operation \cup is defined, then we denote $\mathbf{U}\mathcal{S} = \bigcup_{X \in \mathcal{S}} X$.

Let $p \in \mathbb{N}$ with $p \geq 2$, let $f : \mathbb{N}^p \rightarrow \mathbb{N}$, and let $g : \mathbb{N}^{p-1} \rightarrow \mathbb{N}$. We say that $f(x_1, \dots, x_p) = \mathcal{O}_{x_p}(g(x_1, \dots, x_{p-1}))$ if there is a function $h : \mathbb{N} \rightarrow \mathbb{N}$ such that $f(x_1, \dots, x_p) = \mathcal{O}(h(x_p) \cdot g(x_1, \dots, x_{p-1}))$.

Graphs. All the graphs that we consider in this paper are undirected, finite, and without loops or multiple edges. We use standard graph-theoretic notation, and we refer the reader to [16] for any undefined terminology. Given a graph G , we denote by $V(G)$ the set of vertices of G and by $E(G)$ the set of the edges of G . We call $|V(G)|$ *the size* of G . A graph is *the empty* graph if its size is zero. We also denote by $L(G)$ the set of the vertices of G that have degree exactly one in the case where $|V(G)| \geq 2$, and $L(G) = V(G)$ if $|V(G)| = 1$. If G is a tree (i.e., a connected acyclic graph) then $L(G)$ is the set of the *leaves* of G . A *vertex labeling* of G is some injection $\rho : V(G) \rightarrow \mathbb{N}^+$. Given a vertex $v \in V(G)$, we define the *neighborhood* of v as $N_G(v) = \{u \mid u \in V(G), \{u, v\} \in E(G)\}$ and the *closed neighborhood* of v as $N_G[v] = N_G(v) \cup \{v\}$. If $X \subseteq V(G)$, then we write $N_G(X) = (\bigcup_{v \in X} N_G(v)) \setminus X$. The *degree* of a vertex v in G is defined as $\deg_G(v) = |N_G(v)|$. A graph is called *subcubic* if all its vertices have degree at most three.

A *subgraph* $H = (V_H, E_H)$ of a graph G is a graph such that $V_H \subseteq V(G)$ and $E_H \subseteq E(G) \cap \binom{V_H}{2}$. If $S \subseteq V(G)$, the subgraph of G *induced by* S , denoted $G[S]$, is the graph $(S, E(G) \cap \binom{S}{2})$. We also define $G \setminus S$ to be the subgraph of G induced by $V(G) \setminus S$. If $S \subseteq E(G)$, we denote by $G \setminus S$ the graph $(V(G), E(G) \setminus S)$.

If $s, t \in V(G)$ are two distinct vertices, an (s, t) -*path* of G is any connected subgraph P of G with maximum degree two and where $s, t \in L(P)$, and a *path* is an (s, t) -path for some vertices s and t . We finally denote by $\mathcal{P}(G)$ the set of all paths of G . Given $P \in \mathcal{P}(G)$, we say that $v \in V(P)$ is an *internal vertex* of P if $\deg_P(v) = 2$. Given an integer i and a graph G , we say that G is i -*connected* if for each $\{u, v\} \in \binom{V(G)}{2}$, there exists a set $\mathcal{P}' \subseteq \mathcal{P}(G)$ of (u, v) -paths of G such that $|\mathcal{P}'| = i$ and for each $P_1, P_2 \in \mathcal{P}'$ such that $P_1 \neq P_2$, $V(P_1) \cap V(P_2) = \{u, v\}$.

We denote by K_r , P_r , and C_r the complete graph, the path, and the cycle on r vertices, respectively, and by K_{r_1, r_2} the complete bipartite graph where the one part has

r_1 vertices and the other r_2 .

Block-cut trees. A connected graph G is *biconnected* if for any $v \in V(G)$, $G \setminus \{v\}$ is connected. Notice that K_2 is the only biconnected graph that it is not 2-connected and that K_1 is not biconnected. A *block* of a graph G is a maximal biconnected subgraph of G . We name $\text{block}(G)$ the set of all blocks of G and we name $\text{cut}(G)$ the set of all cut vertices of G . If G is connected, we define the *block-cut tree* of G to be the tree $\text{bct}(G) = (V, E)$ such that

- $V = \text{block}(G) \cup \text{cut}(G)$ and
- $E = \{\{B, v\} \mid B \in \text{block}(G), v \in \text{cut}(G) \cap V(B)\}$.

Note that $L(\text{bct}(G)) \subseteq \text{block}(G)$. It is worth mentioning that the block-cut tree of a graph can be computed in linear time using depth-first search [19]. Basic properties of block-cut trees can be found, for instance, in [10] (where they are called *block trees*).

Let \mathcal{F} be a set of connected graphs on at least two vertices. Given $H \in \mathcal{F}$ and $B \in L(\text{bct}(H))$, we say that (H, B) is an *essential pair* if for each $H' \in \mathcal{F}$ and each $B' \in L(\text{bct}(H'))$, $|E(B)| \leq |E(B')|$. Given an essential pair (H, B) of \mathcal{F} , we define the *first vertex* of (H, B) to be, if it exists, the only cut vertex of H contained in $V(B)$, or an arbitrarily chosen vertex of $V(B)$ otherwise. We define the *second vertex* of (H, B) to be an arbitrarily chosen vertex of $V(B)$ that is a neighbor in $H[B]$ of the first vertex of (H, B) . Note that such a vertex always exists, as a block has at least two vertices. Given an essential pair (H, B) of \mathcal{F} , we assume that the choices for the first and second vertices of (H, B) are fixed.

Moreover, given an essential pair (H, B) of \mathcal{F} , we define the *core* of (H, B) to be the graph $H \setminus (V(B) \setminus \{a\})$ where a is the first vertex of (H, B) . Note that a is a vertex of the core of (H, B) .

Minors and topological minors. Given two graphs H and G and two functions $\phi : V(H) \rightarrow V(G)$ and $\sigma : E(H) \rightarrow \mathcal{P}(G)$, we say that (ϕ, σ) is a *topological minor model of H in G* if

- for every $\{x, y\} \in E(H)$, $\sigma(\{x, y\})$ is an $(\phi(x), \phi(y))$ -path in G and
- if P_1, P_2 are two distinct paths in $\sigma(E(H))$, then none of the internal vertices of P_1 is a vertex of P_2 .

The *branch* vertices of (ϕ, σ) are the vertices in $\phi(V(H))$, while the *subdivision* vertices of (ϕ, σ) are the internal vertices of the paths in $\sigma(E(H))$.

We say that G contains H as a *topological minor*, denoted by $H \preceq_{\text{tm}} G$, if there is a topological minor model (ϕ, σ) of H in G .

Given two graphs H and G and a function $\phi : V(H) \rightarrow 2^{V(G)}$, we say that ϕ is a *minor model of H in G* if

- for every $x, y \in V(H)$ such that $x \neq y$, $\phi(x) \cap \phi(y) = \emptyset$,
- for every $x \in V(H)$, $G[\phi(x)]$ is a connected non-empty graph and
- for every $\{x, y\} \in E(H)$, there exist $x' \in \phi(x)$ and $y' \in \phi(y)$ such that $\{x', y'\} \in E(G)$.

We say that G contains H as a *minor*, denoted by $H \preceq_{\text{m}} G$, if there is a minor model ϕ of H in G .

Graph separators and (topological) minors. Let G be a graph and $S \subseteq V(G)$. Then for each connected component C of $G \setminus S$, we define the *cut-component* of the triple (C, G, S) to be the graph whose vertex set is $V(C) \cup S$ and whose edge set is $E(G[V(C) \cup S])$.

Lemma 1. *Let $i \geq 2$ be an integer, let H be an i -connected graph, let G be a graph, and let $S \subseteq V(G)$ such that $|S| \leq i - 1$. If H is a topological minor (resp. a minor) of G , then there exists a connected component G' of $G \setminus S$ such that H is a topological minor (resp. a minor) of the cut-component of (G', G, S) .*

Proof. We prove the lemma for the topological minor version, and the minor version can be proved with similar arguments. Let i, H, G , and S be defined as in the statement of the lemma. Assume that $H \preceq_{\text{tm}} G$ and let (ϕ, σ) be a topological minor model of H in G . If S is not a separator of G , then the statement is trivial, as in that case the cut-component of $(G \setminus S, G, S)$ is G . Suppose henceforth that S is a separator of G , and assume for contradiction that there exist two connected components G_1 and G_2 of $G \setminus S$ and two distinct vertices x_1 and x_2 of H such that $\phi(x_1) \in V(G_1)$ and $\phi(x_2) \in V(G_2)$. Then, as H is i -connected, there should be i internally vertex-disjoint paths from $\phi(x_1)$ to $\phi(x_2)$ in G . As S is a separator of size at most $i - 1$, this is not possible. Thus, there exists a connected component G' of $G \setminus S$ such that for each $x \in V(H)$, $\phi(x) \in V(G') \cup S$. This implies that H is a topological minor of the cut-component of (G', G, S) . \square

Lemma 2. *Let G be a connected graph, let v be a cut vertex of G , and let V be the vertex set of a connected component of $G \setminus \{v\}$. If H is a connected graph such that $H \preceq_{\text{tm}} G$ and for each leaf B of $\text{bct}(H)$, $B \not\preceq_{\text{tm}} G[V \cup \{v\}]$, then $H \preceq_{\text{tm}} G \setminus V$.*

Proof. Let G , v , V , and H be defined as in the statement of the lemma. Let $B \in \mathcal{L}(\text{bct}(H))$. If B is a single edge, then the condition $B \not\preceq_{\text{tm}} G[V \cup \{v\}]$ implies that $V = \emptyset$. But V is the vertex set of a connected component of $G \setminus \{v\}$ and so $V \neq \emptyset$. This implies that the case B is a single edge cannot occur. If B is not a simple edge, then by definition B is 2-connected and then, by Lemma 1, $B \preceq_{\text{tm}} G \setminus V$. This implies that there is a topological minor model (ϕ, σ) of H in G such that for each $B \in \mathcal{L}(\text{bct}(H))$ and for each $b \in B$, $\phi(b) \notin V$.

We show now that for each $x \in V(H)$, $\phi(x) \notin V$. If $V(H) \setminus (\bigcup_{B \in \mathcal{L}(\text{bct}(H))} V(B)) = \emptyset$ then the result is already proved. Otherwise, let $x \in V(H) \setminus (\bigcup_{B \in \mathcal{L}(\text{bct}(H))} V(B))$. By definition of the block-cut tree, there exist b_1 and b_2 in $\bigcup_{B \in \mathcal{L}(\text{bct}(H))} V(B)$ such that x lies on a (b_1, b_2) -path P of $\mathcal{P}(H)$. Let P_i be the (b_i, x) -subpath of P for each $i \in \{1, 2\}$. By definition of P , we have that $V(P_1) \cap V(P_2) = \{x\}$. This implies that there exists a $(\phi(b_1), \phi(x))$ -path P'_1 and a $(\phi(b_2), \phi(x))$ -path P'_2 in $\mathcal{P}(G)$ such that $V(P'_1) \cap V(P'_2) = \{\phi(x)\}$. Then, as v is a cut vertex of G , it follows that $\phi(x) \notin V$. Thus, for each $x \in V(H)$, $\phi(x) \notin V$. Let $\{x, y\}$ be an edge of $E(H)$. As $\sigma(\{x, y\})$ is a simple $(\phi(x), \phi(y))$ -path, both $\phi(x)$ and $\phi(y)$ are not in V and v is a cut vertex of G , we have, with the same argumentation that before that, for each $z \in V(\sigma(\{x, y\}))$, $z \notin V$. This concludes the proof. \square

Using the same kind of argumentation with minors instead of topological minors, we also obtain the following lemma.

Lemma 3. *Let G be a connected graph, let v be a cut vertex of G , and let V be the vertex set of a connected component of $G \setminus \{v\}$. If H is a graph such that $H \preceq_{\text{m}} G$ and for each leaf B of $\text{bct}(H)$, $B \not\preceq_{\text{m}} G[V \cup \{v\}]$, then $H \preceq_{\text{m}} G \setminus V$.*

In the above two lemmas, we have required graph H to be connected so that $\text{bct}(H)$ is well-defined, but we could relax this requirement, and replace in both statements “for each leaf B of $\text{bct}(H)$ ” with “for each connected component H' of H and each leaf B of $\text{bct}(H')$ ”.

Graph collections. Let \mathcal{F} be a collection of graphs. From now on instead of “collection of graphs” we use the shortcut “collection”. If \mathcal{F} is a non-empty finite collection and all its graphs are non-empty, then we say that \mathcal{F} is a *proper collection*. For any proper collection \mathcal{F} , we define $\text{size}(\mathcal{F}) = \max\{|V(H)| \mid H \in \mathcal{F}\} \cup \{|\mathcal{F}|\}$. Note that if the size of \mathcal{F} is bounded, then the size of the graphs in \mathcal{F} is also bounded. We say that \mathcal{F} is a *planar collection* (resp. *planar subcubic collection*) if it is proper and *at least one* of the graphs in \mathcal{F} is planar (resp. planar and subcubic). We say that \mathcal{F} is a *connected collection* if it is proper and *all* the graphs in \mathcal{F} are connected and of size at least two.

We say that \mathcal{F} is a (*topological*) *minor antichain* if no two of its elements are comparable via the (topological) minor relation.

Let \mathcal{F} be a proper collection. We extend the (topological) minor relation to \mathcal{F} such that, given a graph G , $\mathcal{F} \preceq_{\text{tm}} G$ (resp. $\mathcal{F} \preceq_{\text{m}} G$) if and only if there exists a graph $H \in \mathcal{F}$ such that $H \preceq_{\text{tm}} G$ (resp. $H \preceq_{\text{m}} G$).

Tree and path decompositions. A *tree decomposition* of a graph G is a pair $\mathcal{D} = (T, \mathcal{X})$, where T is a tree and $\mathcal{X} = \{X_t \mid t \in V(T)\}$ is a collection of subsets of $V(G)$ such that:

- $\bigcup_{t \in V(T)} X_t = V(G)$,
- for every edge $\{u, v\} \in E$, there is a $t \in V(T)$ such that $\{u, v\} \subseteq X_t$, and
- for each $\{x, y, z\} \subseteq V(T)$ such that z lies on the unique path between x and y in T , $X_x \cap X_y \subseteq X_z$.

We call the vertices of T *nodes* of \mathcal{D} and the sets in \mathcal{X} *bags* of \mathcal{D} . The *width* of a tree decomposition $\mathcal{D} = (T, \mathcal{X})$ is $\max_{t \in V(T)} |X_t| - 1$. The *treewidth* of a graph G , denoted by $\text{tw}(G)$, is the smallest integer w such that there exists a tree decomposition of G of width at most w . For each $t \in V(T)$, we denote by E_t the set $E(G[X_t])$.

A *path decomposition* of a graph G is a tree decomposition $\mathcal{D} = (T, \mathcal{X})$ of G such that T is a path, and the *pathwidth* of a graph G , denoted by $\text{pw}(G)$, is the smallest integer w such that there exists a path decomposition of G of width at most w . Note that, by definition, for every graph G it holds that $\text{pw}(G) \geq \text{tw}(G)$.

Parameterized complexity. We refer the reader to [14, 17] for basic background on parameterized complexity, and we recall here only some very basic definitions. A *parameterized problem* is a language $L \subseteq \Sigma^* \times \mathbb{N}$. Integer k is called the *parameter* of an instance $I = (x, k) \in \Sigma^* \times \mathbb{N}$. A parameterized problem is *fixed-parameter tractable* (FPT) if there exists an algorithm \mathcal{A} , a computable function f , and a constant c such that given an instance $I = (x, k)$, \mathcal{A} (called an *FPT algorithm*) correctly decides whether $I \in L$ in time bounded by $f(k) \cdot |I|^c$.

Definition of the problems. Let \mathcal{F} be a proper collection. We define the parameter $\mathbf{tm}_{\mathcal{F}}$ as the function that maps graphs to non-negative integers as follows:

$$\mathbf{tm}_{\mathcal{F}}(G) = \min\{|S| \mid S \subseteq V(G) \wedge \mathcal{F} \not\preceq_{\text{tm}} G \setminus S\}. \quad (1)$$

The parameter $\mathbf{m}_{\mathcal{F}}$ is defined analogously, just by replacing $\mathcal{F} \not\preceq_{\text{tm}} G \setminus S$ with $\mathcal{F} \not\preceq_{\text{m}} G \setminus S$. The main objective of this paper is to study the problem of computing the

parameters $\mathbf{tm}_{\mathcal{F}}$ and $\mathbf{m}_{\mathcal{F}}$ for graphs of bounded treewidth under several instantiations of the collection \mathcal{F} . The corresponding decision problems are formally defined as follows.

<p>\mathcal{F}-TM-DELETION</p> <p>Input: A graph G and an integer $k \in \mathbb{N}$.</p> <p>Parameter: The treewidth of G.</p> <p>Output: Is $\mathbf{tm}_{\mathcal{F}}(G) \leq k$?</p>	<p>\mathcal{F}-M-DELETION</p> <p>Input: A graph G and an integer $k \in \mathbb{N}$.</p> <p>Parameter: The treewidth of G.</p> <p>Output: Is $\mathbf{m}_{\mathcal{F}}(G) \leq k$?</p>
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Note that in both above problems, we can always assume that \mathcal{F} is an antichain with respect to the considered relation. Indeed, if \mathcal{F} contains two graphs H_1 and H_2 where $H_1 \preceq_{\mathbf{tm}} H_2$, then $\mathbf{tm}_{\mathcal{F}}(G) = \mathbf{tm}_{\mathcal{F}'}(G)$ where $\mathcal{F}' = \mathcal{F} \setminus \{H_2\}$ (similarly for the minor relation).

Throughout the article, we let n and \mathbf{tw} be the number of vertices and the treewidth of the input graph of the considered problem, respectively.

3 Single-exponential lower bound for any connected \mathcal{F}

In this section we prove the following result.

Theorem 2. *Let \mathcal{F} be a connected collection. Neither \mathcal{F} -TM-DELETION nor \mathcal{F} -M-DELETION can be solved in time $2^{o(\mathbf{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let \mathcal{F} be a connected collection. We present a reduction from VERTEX COVER to \mathcal{F} -TM-DELETION, both parameterized by the treewidth of the input graph, and then we explain the changes to be made to prove the lower bound for \mathcal{F} -M-DELETION. VERTEX COVER cannot be solved in time $2^{o(w)} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails [20] (in fact, it cannot be solved even in time $2^{o(n)}$), where w is the treewidth of the input graph. It is worth mentioning that our reduction bears some similarity with the classical reduction of Yannakakis [27] for general vertex-deletion problems.

Without loss of generality, we can assume that \mathcal{F} is a topological minor antichain. First we select an essential pair (H, B) of \mathcal{F} . Let a be the first vertex of (H, B) , b be the second vertex of (H, B) , and A be the core of (H, B) . For convenience we also refer to a and b as vertices of the copies of A .

Let G be the input graph of the VERTEX COVER problem and let $<$ be an arbitrary total order on $V(G)$. We build a graph G' starting from $G' = (V(G), \emptyset)$. For each vertex v of G , we add a copy of A , which we call A^v , and we identify the vertices v and a . For each edge $e = \{v, v'\} \in E(G)$ with $v < v'$, we remove e , we add a copy of B , which we call B^e , and we identify the vertices v and a and the vertices v' and b . This concludes

the construction of G' . Note that $|V(G')| = |V(G)| \cdot |V(A)| + |E(G)| \cdot |V(B) \setminus \{a, b\}|$ and that $\text{tw}(G') = \max\{\text{tw}(G), \text{tw}(H)\}$. For completeness, we provide a proof of the latter fact.

For each $v \in V(G)$, we define $\mathcal{D}^v = (T^v, \mathcal{X}^v)$ to be a tree decomposition of A^v and we fix $r_v \in V(T^v)$ such that $X_{r_v}^v \in \mathcal{X}^v$ contains the copy of a in A^v . For each $e \in E(G)$, we define $\mathcal{D}^e = (T^e, \mathcal{X}^e)$ to be a tree decomposition of B^e and we fix $r_e \in V(T^e)$ such that $X_{r_e}^e \in \mathcal{X}^e$ contains the copy of a and b in A^v . We know that this bag exists as $\{a, b\} \in E(H)$. Let $\mathcal{D}^G = (T^G, \mathcal{X}^G)$ be a tree decomposition of G . We can then define a tree decomposition of G' as follows. Start from $\mathcal{D} = (T, \mathcal{X})$ where T is the union of T^G , of each T^v , $v \in V(G)$, and each T^e , $e \in E(G)$, and \mathcal{X} is the union of \mathcal{X}^G , of each \mathcal{X}^v , $v \in V(G)$, and each \mathcal{X}^e , $e \in E(G)$. Then for each $v \in V(G)$, we arbitrarily choose $t_v \in V(T^G)$ such that $v \in \mathcal{X}_{t_v}^G$ and connect t_v and r_v in T . Then for each $e \in E(G)$, we arbitrarily choose $t_e \in V(T^G)$ such that $e \subseteq \mathcal{X}_{t_e}^G$ and connect t_e and r_e in T . This concludes the construction of a tree decomposition of G' . As for each $v \in V(G)$ and each $e \in E(G)$, the tree decompositions of A^v and of B^e are just smaller parts of a tree decomposition of H , we obtain that each bag of \mathcal{D} is of size at most $\text{tw}(G)$ if it comes from \mathcal{X}^G , or of size at most $\text{tw}(H)$ if it comes from $\mathcal{X} \setminus \mathcal{X}^G$. Thus $\text{tw}(G') = \max\{\text{tw}(G), \text{tw}(H)\}$.

We claim that there exists a solution of size at most k of VERTEX COVER in G if and only if there is a solution of size at most k of \mathcal{F} -TM-DELETION in G' .

In one direction, assume that S is a solution of \mathcal{F} -TM-DELETION in G' with $|S| \leq k$. By definition of the problem, for each $e = \{v, v'\} \in E(G)$ with $v < v'$, either B^e contains an element of S or A^v contains an element of S . Let $S' = \{v \in V(G) \mid \exists v' \in V(G) : v < v', e = \{v, v'\} \in E(G), (V(B^e) \setminus \{v, v'\}) \cap S \neq \emptyset\} \cup \{v \in V(G) \mid V(A^v) \cap S \neq \emptyset\}$. Then S' is a solution of VERTEX COVER in G and $|S'| \leq |S| \leq k$.

In the other direction, assume that we have a solution S of size at most k of VERTEX COVER in G . We want to prove that S is also a solution of \mathcal{F} -TM-DELETION in G' . For this, we fix an arbitrary $H' \in \mathcal{F}$ and we show that H' is not a topological minor of $G' \setminus S$. First note that the connected components of $G' \setminus S$ are either of the shape $A^v \setminus \{v\}$ if $v \in S$, $B^e \setminus e$ if $e \subseteq S$, or the union of A^v with zero, one, or more graphs $B^{\{v, v'\}} \setminus \{v'\}$ such that $\{v, v'\} \in E(G)$ if $v \in V(G) \setminus S$. As \mathcal{F} is a topological minor antichain, for any $v \in V(G)$, $H' \not\leq_{\text{tm}} A^v \setminus \{v\}$ and for any $e \in E(G)$, $H' \not\leq_{\text{tm}} B^e \setminus e$. Moreover, let $v \in V(G) \setminus S$ and let K be the connected component of $G' \setminus S$ containing v . K is the union of A^v and of every $B^{\{v, v'\}} \setminus \{v'\}$ such that $\{v, v'\} \in E(G)$. As, for each $v' \in V(G)$ such that $\{v, v'\} \in E(G)$, v' is not an isolated vertex in $B^{\{v, v'\}}$, by definition of B , for any $B' \in \text{L}(\text{bct}(H'))$, $|E(B^{\{v, v'\}} \setminus \{v'\})| < |E(B')|$. This implies that for each leaf B' of $\text{bct}(H')$ and for each $\{v, v'\} \in E(G)$, $B' \not\leq_{\text{tm}} B^{\{v, v'\}} \setminus \{v'\}$. It follows by definition of

\mathcal{F} that $H' \not\preceq_{\text{tm}} A^v$. This implies by Lemma 2 that H' is not a topological minor of K . Moreover, as H' is connected by hypothesis, it follows that H' is not a topological minor of $G' \setminus S$ either. This concludes the proof for the topological minor version.

Finally, note that the same proof applies to \mathcal{F} -M-DELETION as well, just by replacing \mathcal{F} -TM-DELETION with \mathcal{F} -M-DELETION, topological minor with minor, \preceq_{tm} with \preceq_{m} , and Lemma 2 with Lemma 3. \square

From Theorem 2 we can easily get the following corollary on planar graphs.

Corollary 1. *Let \mathcal{F} be a connected planar collection. Neither \mathcal{F} -TM-DELETION nor \mathcal{F} -M-DELETION can be solved on planar graphs in time $2^{o(\text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. We can assume that all the graphs in \mathcal{F} are planar, since an input planar graph G does not contain any nonplanar graph as a (topological) minor. We reduce from PLANAR VERTEX COVER to \mathcal{F} -TM-DELETION on planar graphs, both parameterized by the treewidth of the input graph, and the construction of G' is the same as above. Note that since all the graphs in \mathcal{F} are planar, so is the essential pair (H, B) , and therefore the graph G' is easily checked to be planar. Since PLANAR VERTEX COVER cannot be solved in time $2^{o(w)} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails [20, 22], where w is the treewidth of the input graph, the result follows. Finally, the changes to be made for the minor version are the same as those in the proof of Theorem 2. \square

4 Superexponential lower bounds

Let \mathcal{C} be the set of all connected graphs that contain a block with at least five edges, let \mathcal{Q} be the set containing P_5 and all connected graphs that are *not* minors of the banner, and let $\mathcal{S} = \{K_{1,s} \mid s \geq 4\}$. In this section, we prove the following theorems. Note that, by definition, it holds that $\mathcal{C} \subseteq \mathcal{Q}$, but we consider both sets because we will prove a stronger result for the set \mathcal{C} (Theorem 3, which applies to every subset of \mathcal{C}) than for the set \mathcal{Q} (Theorems 4 and 5, which apply to families containing a single graph H).

Theorem 3. *Let \mathcal{F} be a finite non-empty subset of \mathcal{C} . Unless the ETH fails, neither \mathcal{F} -M-DELETION nor \mathcal{F} -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$.*

Theorem 4. *Let $H \in \mathcal{Q}$. Unless the ETH fails, $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$.*

Theorem 5. *Let $H \in \mathcal{Q} \setminus \mathcal{S}$. Unless the ETH fails, $\{H\}$ -TM-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$.*

Note that if H is a connected planar graph such that $H \notin \mathcal{Q}$ (resp. $H \notin \mathcal{Q} \setminus \mathcal{S}$), then $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) can be solved in time $2^{\mathcal{O}(\text{tw})} \cdot n^{\mathcal{O}(1)}$ by the single-exponential algorithms presented in [6]. On the other hand, if H is a connected planar (resp. planar subcubic) graph, then $\{H\}$ -M-DELETION and $\{H\}$ -TM-DELETION can be solved in time $2^{\mathcal{O}(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ by the algorithms presented in [5]. In particular, note that these results altogether settle completely the asymptotic complexity of $\{H\}$ -M-DELETION when H is a connected planar graph; see Figure 1 for an illustration.

We first provide in Section 4.1 a general framework that will be used in every reduction and then we explain how to modify this framework depending on the family \mathcal{F} we are considering.

4.1 The general construction

In order to prove Theorem 3, Theorem 4, and Theorem 5, we will provide reductions from the following problem, which is closely related to the $k \times k$ PERMUTATION CLIQUE problem defined by Lokshtanov et al. [24].

$k \times k$ PERMUTATION INDEPENDENT SET

Input: An integer k and a graph G with vertex set $[1, k] \times [1, k]$.

Parameter: k .

Output: Is there an independent set of size k in G with exactly one element from each row and exactly one element from each column?

Theorem 6 (Lokshtanov et al. [24]). *The $k \times k$ PERMUTATION INDEPENDENT SET problem cannot be solved in time $2^{o(k \log k)}$ unless the ETH fails.*

Let \mathcal{F} be a finite family of non-empty graphs. The framework we are going to present follows the ideas of the construction given by Bonnet et al. [11]. This framework mostly depends on $h := \min_{H \in \mathcal{F}} |V(H)|$ but also on an integer $t_{\mathcal{F}}$ whose value will be defined later. Let (G, k) be an instance of $k \times k$ PERMUTATION INDEPENDENT SET. As we are asking for an independent set that contains exactly one vertex in each row, we will assume without loss of generality that, for each pair $(i, j), (i, j')$ in $V(G)$ with $j \neq j'$, $\{(i, j), (i, j')\} \in E(G)$. We proceed to construct a graph F that contains one gadget for each edge of the graph G . These gadgets are arranged in a cyclic way, separated by some other gadgets ensuring the consistency of the selected solution.

Formally, we first define the graph $K := K_{h-1}$. For each $e \in E(G)$, and each $(i, j) \in [1, k]^2$, we define the graph $B_{i,j}^e$ to be the disjoint union of n_h copies of K , for some integer n_h , whose value will be 2 in the minor case and $\binom{h}{2}$ in the topological minor

case, two new vertices $a_{i,j}^e$ and $b_{i,j}^e$, and $t_{\mathcal{F}}$ other new vertices called *B-extra vertices*. The graph $B_{i,j}^e$ is depicted in Figure 2.

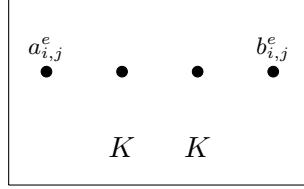


Figure 2: The graph $B_{i,j}^e$ for $e \in E(G)$ and $(i, j) \in [1, k]^2$ when $n_h = 2$ and $t_{\mathcal{F}} = 2$.

Informally, the graph $B_{i,j}^e$, for every $e \in E(G)$, will play in F the role of the vertex (i, j) in G . For each $e \in E(G)$ and each $j \in [1, k]$, we define the graph C_j^e obtained from the disjoint union of every $B_{i,j}^e$, $i \in [1, k]$, such that two graphs $B_{i_1,j}^e$ and $B_{i_2,j}^e$, $i_1 \neq i_2$, are *complete* to each other, that is, for every $i_1 \neq i_2$, if $v_1 \in V(B_{i_1,j}^e)$ and $v_2 \in V(B_{i_2,j}^e)$, then $\{v_1, v_2\} \in E(C_j^e)$. Informally, for a fixed $j \in [1, k]$, the graph C_j^e , for every $e \in E(G)$, corresponds to the column j of G . For every $e \in E(G)$, we also define the *gadget graph* D^e obtained from the disjoint union of every C_j^e , $j \in [1, k]$, by adding, if $e = \{(i, j), (i', j')\}$, every edge $\{v_1, v_2\}$ such that $v_1 \in V(B_{i,j}^e)$ and $v_2 \in V(B_{i',j'}^e)$. The graph D^e is depicted in Figure 3.

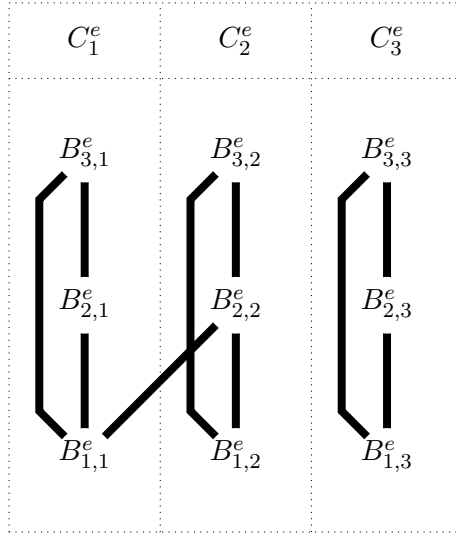


Figure 3: The gadget graph D^e for $e = \{(1, 1), (2, 2)\} \in E(G)$ where $k = 3$. A bold edge means that two graphs B are connected in a complete bipartite way.

Informally, the graph D^e , $e \in E(G)$, encodes the edge e of the graph G . For every $e \in E(G)$, we also define J^e such that $V(J^e) = \{c_j^e \mid j \in [1, k]\} \cup \{r_i^e \mid i \in [1, k]\}$ is a set

of new vertices and $E(J^e) = \emptyset$. It will be helpful to associate the c_j^e 's with “columns” and the r_i^e 's with “rows”. Note that, in the following, J^e may be enhanced, by adding vertices called *J-extra vertices*, whose number depends on the family \mathcal{F} we are working with, but will always be linear in k . The graphs J^e , $e \in E(G)$, are the *separator gadgets* that will ensure the consistency of the selected solution. Finally, the graph F is obtained from the disjoint union of every D^e , $e \in E(G)$, and every J^e , $e \in E(G)$. Moreover, we fix a cyclic permutation σ of the elements of $E(G)$, agreeing that $\sigma^{-1}(e)$ and $\sigma(e)$ is the edge before and after e , respectively, in this cyclic ordering. For each $e \in E(G)$, and each $(i, j) \in [1, k]^2$, we add to F the edges

$$\{b_{i,j}^{\sigma^{-1}(e)}, c_j^e\}, \{b_{i,j}^{\sigma^{-1}(e)}, r_i^e\}, \{r_i^e, a_{i,j}^e\}, \text{ and } \{c_j^e, a_{i,j}^e\}.$$

This concludes the definition of the framework graph F , which is depicted in Figure 4 (a similar figure appears in [11]).

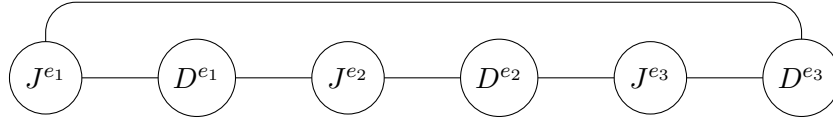


Figure 4: The shape of the framework graph F assuming that $k = 3$, G contains only the three edges e_1 , e_2 , and e_3 , and σ is the cyclic permutation (e_1, e_2, e_3) .

Note that in later constructions a gadget graph D^e , $e \in E(G)$, will only be connected to the (enhanced) separator gadgets J^e and $J^{\sigma(e)}$ in a way that will be specified later and that depends on the family \mathcal{F} .

Let $\ell := (n_h(h-1) + 2 + t_{\mathcal{F}})(k-1)km$, where $m = |E(G)|$. Note that $z = (n_h(h-1) + 2 + t_{\mathcal{F}})$ is the number of vertices of a $B_{i,j}^e$, $i, j \in [1, k]$, $e \in E(G)$ and that ℓ is the budget needed to select the vertex set of exactly $k-1$ graphs $B_{i,j}^e$, $i \in [1, k]$ in each graph C_j^e , $j \in [1, k]$, $e \in E(G)$. The pair (F, ℓ) is called the \mathcal{F} -*M-framework* of (G, k) when $n_h = 2$, and the \mathcal{F} -*TM-framework* of (G, k) when $n_h = \binom{h}{2}$. When the value n_h is not relevant, the pair (F, ℓ) is simply called the \mathcal{F} -*framework* of (G, k) . For convenience, we always assume some prespecified permutation σ associated with the graph F .

For each family \mathcal{F} , given an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, we will consider (F, ℓ) , the \mathcal{F} -framework of (G, k) , and create another pair $(F_{\mathcal{F}}, \ell)$, called the *enhanced \mathcal{F} -framework*, where $F_{\mathcal{F}}$ is a graph obtained from F by adding some new vertices and edges. The added vertices will be *B-extra vertices* or *J-extra vertices*. The added edges will be either inside some (enhanced) $B_{i,j}^e$, or from some D^e to the (enhanced) J^e and $J^{\sigma(e)}$. More formally, the additional edges will be from the set

$$\left(\bigcup_{\substack{e \in E(G), \\ i, j \in [1, k]}} V(B_{i,j}^e) \times V(B_{i,j}^e) \right) \cup \left(\bigcup_{e \in E(G)} V(D^e) \times V(J^e) \right) \cup \left(\bigcup_{e \in E(G)} V(D^e) \times V(J^{\sigma(e)}) \right),$$

by interpreting ordered pairs as edges. Note that F will always be a subgraph of $F_{\mathcal{F}}$. We will claim that there exists a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) if and only if there exists a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on $(F_{\mathcal{F}}, \ell)$. In order to do this, we first prove a generic lemma, namely Lemma 6, and then provide a property, namely Property 1 (resp. Property 2), which we will prove for each family \mathcal{F} depending on the enhanced \mathcal{F} -framework $F_{\mathcal{F}}$.

Let us now provide an upper bound on the treewidth (in fact, the pathwidth) of $F_{\mathcal{F}}$. Let $\{e_1, \dots, e_m\} = E(G)$ such that for each $i \in [1, m]$, $\sigma(e_i) = e_{i+1}$ with the convention that $e_{m+1} = e_1$. First note that, for each $e \in E(G)$, the set $V(J^e) \cup V(J^{\sigma(e)})$ disconnects the vertex set $V(D^e)$ from the rest of $F_{\mathcal{F}}$. Moreover, if $e = \{(i, j), (i', j')\}$, then the bags

$$\begin{aligned} &V(J^{e_1}) \cup V(J^e) \cup V(J^{\sigma(e)}) \cup V(C_1^e) \cup V(B_{i,j}^e) \cup V(B_{i',j'}^e), \\ &V(J^{e_1}) \cup V(J^e) \cup V(J^{\sigma(e)}) \cup V(C_2^e) \cup V(B_{i,j}^e) \cup V(B_{i',j'}^e), \dots, \\ &V(J^{e_1}) \cup V(J^e) \cup V(J^{\sigma(e)}) \cup V(C_k^e) \cup V(B_{i,j}^e) \cup V(B_{i',j'}^e) \end{aligned}$$

form a path decomposition of $G[V(J^e) \cup V(J^{\sigma(e)}) \cup V(D^e)]$ of width $(n_h(h-1) + 2 + t_{\mathcal{F}})(k+1) - 1 + 3 \cdot |V(J^{e_1})|$ (using the fact that $|V(J^e)| = |V(J^{e_1})|$ for each $e \in E(G)$). Let denote by P^e this decomposition. By concatenating the path decompositions P^{e_1}, P^{e_2}, \dots , and P^{e_m} , we obtain a path decomposition of $F_{\mathcal{F}}$ whose width, by using the fact that $|V(J^{e_1})| = \mathcal{O}(k)$, is linear in k .

We start by proving that for each column C_j^e , $e \in E(G)$, $j \in [1, k]$, containing a minimum number of vertices of the solution, all the remaining vertices belong to the same row of this column.

Lemma 4. *Let \mathcal{F} be a family of graphs and let (F, ℓ) be the \mathcal{F} -TM-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. Let S be a solution of \mathcal{F} -TM-DELETION on (F, ℓ) and let $e \in E(G)$ and $j \in [1, k]$ such that the quantity $|V(C_j^e) \setminus S|$ is maximized, i.e., for all $e' \in E(G)$ and $j' \in [1, k]$ $|V(C_j^e) \setminus S| \geq |V(C_{j'}^{e'}) \setminus S|$. Then there exists $i \in [1, k]$ such that $V(C_j^e) \setminus S \subseteq V(B_{i,j}^e)$.*

Proof. We set $z = \binom{h}{2}(h-1) + 2 + t_{\mathcal{F}}$ and observe that $|S| \leq \ell = z(k-1)km$. Let $h = \min_{H \in \mathcal{F}} |V(H)|$, and note that we can always assume that $h \geq 2$ as otherwise the problem is trivial. Choose $e \in E(G)$ and $j \in [1, k]$ so that the quantity $|V(C_j^e) \setminus S|$ is maximized. In order to prove the lemma, we show that the assumption that there exist $i_1, i_2 \in [1, k]$, with $i_1 \neq i_2$ such that $(V(C_j^e) \setminus S) \cap V(B_{i_1,j}^e) \neq \emptyset$ and $(V(C_j^e) \setminus S) \cap V(B_{i_2,j}^e) \neq \emptyset$ implies that $F \setminus S$ contains K_h as a topological minor, for any value of $t_{\mathcal{F}} \in \mathbb{N}$.

We claim that $|V(C_j^e) \setminus S| \geq z$. Indeed, if $|V(C_j^e) \setminus S| < z$, by the maximality in the choice of e and j , it follows that $|V(C_{j'}^{e'}) \setminus S| \leq z - 1$, for all $e' \in E(G)$ and $j' \in [1, k]$.

This implies that $|S \cap V(C_{j'}^{e'})| \geq |V(C_{j'}^{e'})| - (z - 1) = kz - z + 1 = z(k - 1) + 1$ for all $e' \in E(G)$ and $j' \in [1, k]$. As there are m choices for e' and k choices for i' we have that

$$|S| \geq \bigcup_{\substack{e' \in E(G) \\ i' \in [1, k]}} |S \cap V(C_{j'}^{e'})| \geq (z(k - 1) + 1)km > z(k - 1)km = \ell,$$

a contradiction. We just proved that

$$|V(C_j^e) \setminus S| \geq \binom{h}{2}(h - 1) + 2 + t_{\mathcal{F}}. \quad (2)$$

We next pick U_M as a set $V(B_{i,j}^e) \setminus S$, $i \in [1, k]$, with the maximum number of elements. We claim that, if $|U_M| \leq h - 1$, then $F \setminus S$ contains K_h as a topological minor. We first observe that at least h of the sets in $\mathcal{Z} = \{V(B_{i,j}^e) \setminus S \mid i \in [1, k]\}$ are non-empty. To verify this, suppose to the contrary that the set \mathcal{Z}' , consisting of the non-empty elements of \mathcal{Z} , has cardinality at most $h - 1$. By the maximality of the choice of U_M , we obtain that each set in \mathcal{Z}' has at most $|U_M|$ elements. We then observe that $|V(C_j^e) \setminus S| = |\bigcup \mathcal{Z}'| = (h - 1) \cdot |U_M| \leq (h - 1)^2$, a contradiction to (2), as $(h - 1)^2 < \binom{h}{2}(h - 1) + 2 + t_{\mathcal{F}}$. We just proved that $|\mathcal{Z}'| \geq h$. By picking one vertex from each set in \mathcal{Z}' , we conclude that $F \setminus S$ contains a clique of size h as a subgraph and the claim follows.

From now on, we assume that $h \leq |U_M|$. In fact, we claim that if $|U_M| \leq \binom{h}{2}(h - 1) + 2 + t_{\mathcal{F}} - \binom{h}{2}$, then $F \setminus S$ contains K_h as a topological minor. For this, let Q be any set of h vertices of U_M and let Z be a set of $\binom{h}{2}$ vertices of $V(C_j^e) \setminus (S \cup U_M)$ (this set exists because of (2)). As each vertex of Z is a neighbor of each vertex of Q , we obtain that $F[Q \cup Z]$, which is a subgraph of $F \setminus S$, contains K_h as a topological minor.

According to the previous claim, we can assume that $|U_M| > \binom{h}{2}(h - 1) + 2 + t_{\mathcal{F}} - \binom{h}{2}$ or, equivalently

$$\binom{h}{2}(h - 2) + 2 + t_{\mathcal{F}} < |U_M|. \quad (3)$$

Let a be the size of the largest clique K_a in $F[U_M]$. We claim that $a \geq h - 1$. For this, observe that if the biggest clique in $F[U_M]$ has size at most $h - 2$, then $|U_M| \leq \binom{h}{2}(h - 2) + 2 + t_{\mathcal{F}}$, contradicting (3).

We just derived that $F[U_M]$ contains a clique K_{h-1} . By our initial assumption, we have that $U_A = (V(C_j^e) \setminus S) \cap V(B_{i',j}^e) \neq \emptyset$ for some $i' \neq i$. By combining K_a with any vertex of U_A , we obtain K_h as a subgraph of $F \setminus S$, and the lemma follows. \square

Note that Lemma 4 is also valid for the minor version with the same \mathcal{F} -TM-framework. However, for our future constructions, we need this statement to hold for the (smaller) \mathcal{F} -M-framework as well, where $n_h = 2$.

Lemma 5. *Let \mathcal{F} be a family of graphs and let (F, ℓ) be the \mathcal{F} -M-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. Let S be a solution of \mathcal{F} -M-DELETION on (F, ℓ) and let $e \in E(G)$ and $j \in [1, k]$ such that the quantity $|V(C_j^e) \setminus S|$ is maximized, i.e., for all $e' \in E(G)$ and $j' \in [1, k]$ $|V(C_j^e) \setminus S| \geq |V(C_{j'}^{e'}) \setminus S|$. Then there exists $i \in [1, k]$ such that $V(C_j^e) \setminus S \subseteq V(B_{i,j}^e)$.*

Proof. Let $h = \min_{H \in \mathcal{F}} |V(H)|$. Recall that $\ell = (2(h-1) + 2 + t_{\mathcal{F}})(k-1)km$. Again, in order to prove the lemma, we show that the assumption that there exist $i_1, i_2 \in [1, k]$, with $i_1 \neq i_2$ such that $(V(C_j^e) \setminus S) \cap V(B_{i_1,j}^e) \neq \emptyset$ and $(V(C_j^e) \setminus S) \cap V(B_{i_2,j}^e) \neq \emptyset$ implies that $F \setminus S$ contains K_h as a minor, for any value of $t_{\mathcal{F}} \in \mathbb{N}$.

Let us fix the value of $t_{\mathcal{F}} \in \mathbb{N}$. Let $e \in E(G)$ and $j \in [1, k]$ be such that $|V(C_j^e) \setminus S|$ is maximized.

We claim that $|V(C_j^e) \setminus S| \geq 2h + t_{\mathcal{F}}$. Indeed, if $|V(C_j^e) \setminus S| < 2h + t_{\mathcal{F}}$, by the maximality in the choice of e and j , it follows that $|V(C_{j'}^{e'}) \setminus S| \leq 2h + t_{\mathcal{F}} - 1$, for all $e' \in E(G)$ and $j' \in [1, k]$. This implies that $|S \cap V(C_{j'}^{e'})| \geq |V(C_{j'}^{e'})| - (2h + t_{\mathcal{F}} - 1) = k(2(h-1) + t_{\mathcal{F}} + 2) - (2h + t_{\mathcal{F}} - 1) = (2h + t_{\mathcal{F}})(k-1) + 1$ for all $e' \in E(G)$ and $j' \in [1, k]$. As there are m choices for e' and k choices for i' we have that

$$|S| \geq \bigcup_{\substack{e' \in E(G) \\ i' \in [1, k]}} |S \cap V(C_{j'}^{e'})| \geq ((2h + t_{\mathcal{F}})(k-1) + 1)km > (2h + t_{\mathcal{F}})(k-1)km = \ell,$$

a contradiction. We just proved that

$$|V(C_j^e) \setminus S| \geq 2h + t_{\mathcal{F}}. \quad (4)$$

An edge $e \in E(C_j^e)$ is *transversal* if there is no $i \in [1, k]$ such that both endpoints of e belong to $B_{i,j}^e$. The important property of a transversal edge $e = \{v_1, v_2\}$ is that $N_{C_j^e}(\{v_1, v_2\}) = V(C_j^e) \setminus \{v_1, v_2\}$. A *transversal matching* of C_j^e is a matching that contains only transversal edges. Note that if there exists a transversal matching M of size h over a set of vertices $T \subseteq V(C_j^e) \setminus S$, then, by contracting every edge of M , it follows that $C_j^e[T]$, and therefore $F \setminus S$ as well, contains K_h as a minor.

Let U_M be a set $V(B_{i,j}^e) \setminus S$, $i \in [1, k]$, with the maximum number of elements. If $|U_M| \leq h + t_{\mathcal{F}}$, then, because of (4), the graph $C_j^e \setminus S$ contains a transversal matching of size at least h . This can be seen, for instance, by considering the complete k -partite graph where each part contains the vertices in $V(B_{i,j}^e) \setminus S$, for $i \in [1, k]$, and noting that it admits a perfect matching (which defines a transversal matching in $C_j^e \setminus S$ of size at least h) by applying Tutte's criterion [16] on the existence of a perfect matching in a general graph (recall that a graph G contains a perfect matching if and only if there is no set $S \subseteq V(G)$ whose removal generates more than $|S|$ odd-sized components). Thus $C_j^e \setminus S$, and therefore $F \setminus S$ as well, contains a clique of h vertices as a minor.

Assume now that $h + t_{\mathcal{F}} < |U_M|$. Let a be the maximum size of a clique in $C_j^e[U_M]$. As $h \geq 1$ we have that $|U_M| \geq t_{\mathcal{F}} + 2$, therefore

$$a \geq \left\lceil \frac{|U_M| - (t_{\mathcal{F}} + 2)}{2} \right\rceil. \quad (5)$$

We claim that, if $|U_M| < 2h + t_{\mathcal{F}}$, then $F \setminus S$ contains a clique of h vertices as a minor. For this, we set $U_A := V(C_j^e) \setminus (S \cup U_M)$ and we distinguish two cases, depending on the parity of the quantity $|U_M| - t_{\mathcal{F}}$.

Case 1: $|U_M| = t_{\mathcal{F}} + 2u$, with $h < 2u < 2h$. Then, from (5), $a \geq (2u - 2)/2$, therefore $C_j^e[U_M]$ contains a clique K^* of size $u - 1$ while the vertices of $C_j^e[U_M]$ that are not in K^* are $t_{\mathcal{F}} + u + 1$. Moreover, by (4), $|U_A| \geq (2h + t_{\mathcal{F}}) - (t_{\mathcal{F}} + 2u) = 2h - 2u$, so $C_j^e \setminus (S \cup V(K^*))$ contains a transversal matching of size at least $q := \min\{2h - 2u, t_{\mathcal{F}} + u + 1\}$, which, when contracted, creates a clique K^+ of size at least q whose vertices are connected with all $u - 1$ vertices of K^* . Also, using the inequality $h < 2u < 2h$, we obtain

$$\begin{aligned} (u - 1) + (t_{\mathcal{F}} + u + 1) &\geq 2u > h \text{ and} \\ (u - 1) + (2h - 2u) &= 2h - (u + 1) \geq h, \end{aligned}$$

therefore, in any case, $(u - 1) + q \geq h$. Thus, by taking K^* with K^+ and all the edges between them, we deduce that $C_j^e \setminus S$, and therefore $F \setminus S$ as well, contains a clique of size at least h as a minor.

Case 2: $|U_M| = t_{\mathcal{F}} + 2u + 1$, with $h < 2u + 1 < 2h$. Then $C_j^e[U_M]$ contains a clique K^* of size u and $t_{\mathcal{F}} + u + 1$ vertices outside this clique K^* . Moreover, again by (4), $|U_A| \geq 2h - 2u - 1$, so $C_j^e \setminus (S \cup V(K^*))$ contains a transversal matching of size at least $q := \min\{2h - 2u - 1, t_{\mathcal{F}} + u + 1\}$. On the other hand, using the inequality $h < 2u + 1 < 2h$, we know that

$$\begin{aligned} u + (t_{\mathcal{F}} + u + 1) &\geq 2u + 1 > h \text{ and} \\ u + (2h - 2u - 1) &= 2h - (u + 1) \geq h. \end{aligned}$$

Thus $u + q \geq h$, and, as in the previous case, we deduce that $C_j^e \setminus S$, and therefore $F \setminus S$ as well, contains a clique of size h as a minor. The claim follows.

Therefore, what remains is to examine the case where $|U_M| \geq 2h + t_{\mathcal{F}}$. In this case, because of (5), $C_j^e[U_M]$ contains a clique of size $a \geq (2h - 2)/2 = h - 1$. Combining this clique with any vertex in a set $(V(C_j^e) \setminus S) \cap V(B_{i',j}^e)$ that, because of our initial assumption, is non-empty for some $i' \neq i$, we obtain K_h as a subgraph of $F \setminus S$, and the lemma follows. \square

The purpose of Lemma 4 and Lemma 5 is to obtain Lemma 6 that states that for any solution S of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) and for any $B_{i,j}^e$, $e \in E(G)$ and $(i, j) \in [1, k]^2$, either $V(B_{i,j}^e) \cap S = \emptyset$ or $V(B_{i,j}^e) \subseteq S$. Moreover, there is exactly one $B_{i,j}^e$ such that $V(B_{i,j}^e) \cap S = \emptyset$ in each column C_j^e , $e \in E(G)$, $j \in [1, k]$.

Lemma 6. *Let \mathcal{F} be a family of graphs and let (F, ℓ) be the \mathcal{F} -M-framework (resp. \mathcal{F} -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. For every solution S of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on (F, ℓ) , for every $e \in E(G)$ and every $j \in [1, k]$, there exists $i \in [1, k]$ such that $V(C_j^e) \setminus S = V(B_{i,j}^e)$. Moreover, for every $e \in E(G)$, $V(J^e) \cap S = \emptyset$.*

Proof. Let S be a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on $(F, (n_h(h-1) + 2 + t_{\mathcal{F}})(k-1)km)$. By Lemma 5 (resp. Lemma 4), we know that for every $e \in E(G)$ and every $j \in [1, k]$, there is some i such that for every $i' \in [1, k] \setminus \{i\}$, $B_{i',j}^e \subseteq S$. As each $B_{i,j}^e$ has $n_h(h-1) + 2 + t_{\mathcal{F}}$ vertices, we obtain that $|V(C_j^e) \cap S| \geq (n_h(h-1) + 2 + t_{\mathcal{F}})(k-1)$. As there are exactly m edges and k columns, the budget is tight and we obtain that $|V(C_j^e) \cap S| = (n_h(h-1) + 2 + t_{\mathcal{F}})(k-1)$. This implies that $|V(C_j^e) \setminus S| = n_h(h-1) + 2 + t_{\mathcal{F}}$, corresponding to the size of a set $B_{i,j}^e$, for some $i \in [1, k]$. Moreover, as all the vertices of S are vertices in the sets D^e , we also have that for every $e \in E(G)$, $V(J^e) \cap S = \emptyset$. The lemma follows. \square

Using Lemma 6, for every edge $e \in E(G)$, it will be possible to make a correspondence between a permutation corresponding to a solution of $k \times k$ PERMUTATION INDEPENDENT SET and the k pairs (i, j) in $[1, k]^2$ for which $V(B_{i,j}^e) \cap S = \emptyset$. In order to ensure the consistency of the selected solution among the gadget graphs D^e , $e \in E(G)$, we need to show that given $(i, j) \in [1, k]^2$, if there exists $e \in E(G)$ such that $V(B_{i,j}^e) \cap S = \emptyset$, then for every $e' \in E(G)$, we have $V(B_{i,j}^{e'}) \cap S = \emptyset$. For this, we state two properties, namely Property 1 and Property 2, applying to the minor and topological minor version of the problem, respectively. Then we prove Lemma 7, stating that if the corresponding property holds, then we indeed have the desired consistency for the corresponding problem, which allows to find a solution of $k \times k$ PERMUTATION INDEPENDENT SET.

Property 1. *Let \mathcal{F} be a family of graphs and let $(F_{\mathcal{F}}, \ell)$ be the enhanced \mathcal{F} -M-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. Let S be a solution of \mathcal{F} -M-DELETION on $(F_{\mathcal{F}}, \ell)$. For every $e \in E(G)$, and for every $i, j \in [1, k]$, if $b_{i,j}^e \notin S$ then for every $i' \in [1, k] \setminus \{i\}$, we have $a_{i',j}^{\sigma(e)} \in S$.*

Property 2. *Let \mathcal{F} be a family of graphs and let $(F_{\mathcal{F}}, \ell)$ be the enhanced \mathcal{F} -TM-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. Let S be a solution of \mathcal{F} -TM-DELETION on $(F_{\mathcal{F}}, \ell)$. For every $e \in E(G)$, and for every $i, j \in [1, k]$, if $b_{i,j}^e \notin S$ then for every $i' \in [1, k] \setminus \{i\}$, we have $a_{i',j}^{\sigma(e)} \in S$.*

The above properties state that the choices of the vertices $a_{i,j}^e, b_{i,j}^e$ are consistent through the graph $F_{\mathcal{F}}$.

Lemma 7. *Let \mathcal{F} be a family of graphs, let $(F_{\mathcal{F}}, \ell)$ be the enhanced \mathcal{F} -M-framework (resp. enhanced \mathcal{F} -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET. If Property 1 (resp. Property 2) holds and there exists a solution S of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on $(F_{\mathcal{F}}, \ell)$, then, for any $e \in E(G)$, the set $T^e = \{(i, j) \mid V(B_{i,j}^e) \cap S = \emptyset\}$ is a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Moreover, for any $e_1, e_2 \in E(G)$, $T^{e_1} = T^{e_2}$.*

Proof. Let S be a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on $(F_{\mathcal{F}}, \ell)$. Note that this implies that S is also a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on (F, ℓ) , the \mathcal{F} -M-framework (resp. \mathcal{F} -TM-framework) of (G, k) , and so, Lemma 6 can be applied. Let σ be the cyclic permutation associated with $F_{\mathcal{F}}$. For each $e \in E(G)$, let $T^e = \{(i, j) \mid V(B_{i,j}^e) \cap S = \emptyset\}$. By Lemma 6, for each $e \in E(G)$, T^e contains exactly one element from each column. We first show that for any $e \in E(G)$, $T^e = T^{\sigma(e)}$. Let $e \in E(G)$ and let $(i, j) \in T^e$. As $(i, j) \in T^e$, we have that $b_{i,j}^e \notin S$. By Property 1 (resp. Property 2), for each $i' \in [1, k] \setminus \{i\}$ it holds that $a_{i',j}^{\sigma(e)} \in S$, and thus $(i', j) \notin T^{\sigma(e)}$. As $T^{\sigma(e)}$ contains exactly one element from each column it follows that $(i, j) \in T^{\sigma(e)}$. As both T^e and $T^{\sigma(e)}$ are of size exactly k , we obtain that $T^e = T^{\sigma(e)}$.

By repeating the above argument iterating cyclically along the permutation σ , we obtain that for any $e_1, e_2 \in E(G)$, $T^{e_1} = T^{e_2}$. Let $(i, j), (i', j') \in T^{e_1}$. The existence of an edge $e = \{(i, j), (i', j')\}$ in $E(G)$ implies that in D^e , there is a vertex in $B_{i',j'}^e$ (in fact, any vertex of $B_{i',j'}^e$) that is fully connected to a copy of $K = K_{h-1}$ that is in $B_{i,j}^e$ (in fact, to all such copies), and so D^e contains the clique K_h as a subgraph. As $h = \min_{H \in \mathcal{F}} |V(H)|$, this is not possible, and therefore T^e is an independent set in G of size k . Moreover, by the construction of G , T^e contains at most one vertex per row and by Lemma 6, it contains exactly one vertex per column. The lemma follows. \square

Given a solution P of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) , we define $S_P = \{v \in V(F_{\mathcal{F}}) \mid v \in B_{i,j}^e : e \in E(G), (i, j) \in [1, k]^2 \setminus P\}$, where $(F_{\mathcal{F}}, \ell)$ is the enhanced \mathcal{F} -framework of (G, k) . Note that $|S_P| = \ell$. In what follows we will prove that S_P is a solution of \mathcal{F} -M-DELETION (or \mathcal{F} -TM-DELETION) on $(F_{\mathcal{F}}, \ell)$ for each instantiation of $F_{\mathcal{F}}$ that we will consider.

We now proceed to describe how to complete, starting from F , the construction of the enhanced \mathcal{F} -M-framework (or enhanced \mathcal{F} -TM-framework) $F_{\mathcal{F}}$, depending on \mathcal{F} , towards proving Theorems 3, 4, and 5.

4.2 The reduction for $\mathcal{F} \subseteq \mathcal{C}$

In order to prove Theorem 3, we will need some extra definitions.

Given a finite graph H , we define the *block edge size function* $\text{besf}_H : \mathbb{N} \rightarrow \mathbb{N}$ to be such that for any $x \in \mathbb{N}$, $\text{besf}_H(x)$ equals the number of edges of H that are contained in a block with at least x edges. Note that this function is a decreasing function and, as we only deal with finite graphs, for any finite graph H , there exists $x_0 \in \mathbb{N}$ such that $\text{besf}_H(x_0) = 0$ (notice that the minimum such x_0 is one more than the maximum number of edges of a block of H). Given two block edge size functions f and g , corresponding to two graphs, we say that $f \prec g$ if there exists an x_0 , called a *witness* of the inequality, such that $f(x_0) < g(x_0)$ and for each $x \geq x_0$, $f(x) \leq g(x)$. It can be verified that \prec is a total order on the set $\{\text{besf}_H \mid H \text{ is a finite graph}\}$. Note also that given two graphs H and H' , if $\text{besf}_H \prec \text{besf}_{H'}$ then H' cannot be a minor of H . Intuitively, if one considers only the blocks of H' with at least x_0 edges, then there are too many edges to fit within the blocks of H with at least x_0 edges, where x_0 is a witness of the inequality.

Given an integer k and a graph H that contains at least one block with at least k edges, a *k -edges leaf block cut* is a tuple (X, Y, B, v) , where

- B is a block with at least k edges,
- $v \in V(B)$ is a cut vertex, or, in case H is 2-connected, any vertex of $H = B$,
- X and Y are two subsets of $V(H)$ such that $X \cup Y = V(H)$ and $X \cap Y = \{v\}$,
- $Y \setminus \{v\}$ is the vertex set of the connected component of $H \setminus \{v\}$ that contains $V(B) \setminus \{v\}$, and
- $H[Y]$ contains only one block with at least k edges (which is precisely B).

Intuitively, B would be a leaf of the block-cut tree of the graph H' obtained from H by iteratively removing every leaf block with at most $k - 1$ edges from $\text{bct}(H)$, v is the unique remaining neighbor of B in $\text{bct}(H')$, Y consists of B , and X is the rest of the remaining graph together with v . Note that as long as H contains at least one block with at least k edges, there exists a k -edges leaf block cut.

We are now ready to prove Theorem 3.

Proof of Theorem 3. Let $\mathcal{F} \subseteq \mathcal{C}$ and let (F, ℓ) be the \mathcal{F} -M-framework (resp. \mathcal{F} -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}} = 0$. Let $H \in \mathcal{F}$ be a graph that minimizes besf_H with regard to the relation \prec over all the graphs of \mathcal{F} . Let (X, Y, B, v) be a 5-edge leaf block cut of H . Let $b = |E(B)|$, let $H_X = H[X]$, and let $H_Y = H[Y]$. Let v' be a neighbor of v in B and H_Y^- be the graph

obtained from H_Y by removing $\{v, v'\}$. Note that H_X and H_Y^- are connected and H_Y contains only one block with at least five edges, and this block is precisely B .

We are now ready to describe the graph $F_{\mathcal{F}}$. All the new vertices are J -extra vertices. Namely, starting from F , for each $e \in E(G)$, we add a copy of H_X , and we denote by q^e the copy of v . For each $e \in E(G)$ and for each $i \in [1, k]$, we add a copy of H_Y^- where we identify v and q^e , and u and r_i^e . This completes the definition of $F_{\mathcal{F}}$. We stress that in this construction there are no B -extra vertices, i.e., $t_{\mathcal{F}} = 0$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\mathcal{F}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$ (recall that $h = \min_{H \in \mathcal{F}} |V(H)|$), or the graph Z depicted in Figure 5. We claim that $\text{besf}_Z \prec \text{besf}_H$ with witness b . Indeed, since $|E(H_Y^-)| < |E(H_Y)|$, $|E(B)| \geq 5$, and the blocks of Z that are not copies of H_X or H_Y^- have four edges, it follows that $\text{besf}_Z(b) < \text{besf}_H(b)$, and $\text{besf}_Z(b') \leq \text{besf}_H(b')$ for all $b' > b$. Therefore, $\text{besf}_Z \prec \text{besf}_H$ with witness b . This in turn implies, because of the choice of H , that $\text{besf}_Z \prec \text{besf}'_H$ for each $H' \in \mathcal{F}$. Therefore, no $H' \in \mathcal{F}$ is a minor of the graph Z . Thus S_P is a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) of size ℓ .

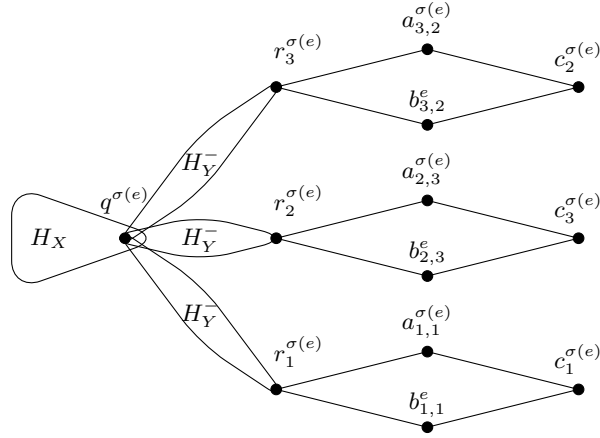


Figure 5: A connected component Z of $F_H \setminus S$ that is not a copy of K , with $T^e = \{(1, 1), (2, 3), (3, 2)\}$.

Assume now that S is a solution of \mathcal{F} -M-DELETION (resp. \mathcal{F} -TM-DELETION) on $F_{\mathcal{F}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$, then, since by Lemma 6 it holds that $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the copy of H_X , the copy of H_Y^- between $q^{\sigma(e)}$ and $r_{i'}^{\sigma(e)}$, and the path that starts at $q^{\sigma(e)}$, goes through the corresponding copy of H_Y^- until $r_i^{\sigma(e)}$, and continues with the vertices $b_{i,j}^e, c_j^{\sigma(e)}, a_{i',j}^{\sigma(e)}$, and $r_{i'}^{\sigma(e)}$, induce a graph that contains H as a topological minor.

This implies that if $a_{i',j}^{\sigma(e)} \notin S$, H is a topological minor of $F_{\mathcal{F}} \setminus S$. As this is not possible by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the theorem follows from Lemma 7. \square

4.3 The reduction for $\{H\}$ -DELETION

This section is dedicated to the proofs of Theorem 4 and Theorem 5, which we restate here for better readability.

Theorem 4. *Let $H \in \mathcal{Q}$. Unless the ETH fails, $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$.*

Theorem 5. *Let $H \in \mathcal{Q} \setminus \mathcal{S}$. Unless the ETH fails, $\{H\}$ -TM-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$.*

Thus, we will focus on cases where the family \mathcal{F} contains only one graph H . We start with a number of lemmas, namely Lemma 8 up to Lemma 15, in which we distinguish several cases according to properties of H such that its number of cut vertices and the presence of certain cycles and vertices of degree one. Altogether, these cases will cover all the possible graphs H considered in Theorem 4 and Theorem 5. The proofs of each of these lemmas are quite similar and follow the same structure. Namely, we first describe the graph $F_{\{H\}}$, and then we prove the equivalence between the existence of solutions of $k \times k$ PERMUTATION INDEPENDENT SET and $\{H\}$ -M-DELETION (or $\{H\}$ -TM-DELETION). In the reverse direction, we will prove that Property 1 and Property 2 hold, and therefore we can apply Lemma 7.

Thanks to Theorem 3, we can assume that each block of H contains at most four edges, i.e., each block of H is an edge, a C_3 , or a C_4 . In this setting, we have $h = |V(H)|$.

Lemma 8. *Let H be a connected graph such that the number of cycles (of size three or four) in H with at least two cut vertices is exactly one. Neither $\{H\}$ -M-DELETION nor $\{H\}$ -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework (resp. $\{H\}$ -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}} = 0$. Let B be the block of H with at least three edges and two cut vertices, and let $\{v, v'\}$ be an edge of B . Let H^- be the graph H where the edge $\{v, v'\}$ has been removed.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , for each $e \in E(G)$, we introduce a new vertex q^e . For each $e \in E(G)$ and each $i \in [1, k]$, we add a copy of H^- where we identify v and r_i^e , and v' and q^e . This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of

size $h - 1$, or the graph Z depicted in Figure 6. As Z does not contain any cycle that contains at least two cut vertices, we obtain that H is not a minor of Z . Thus S_P is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) of size ℓ .

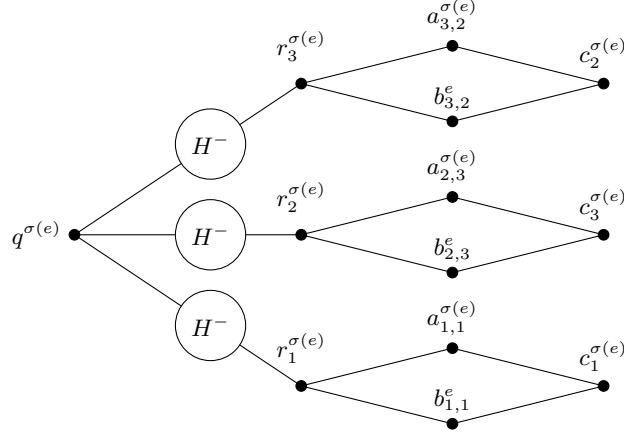


Figure 6: A connected component Z of $F_{\{H\}} \setminus S_P$ that is not a copy of K , with $T^e = \{(1, 1), (2, 3), (3, 2)\}$.

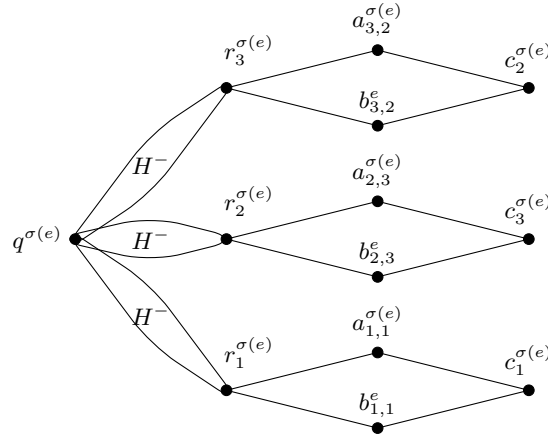


Figure 7: A connected component Z of $F_{\{H\}} \setminus S_P$ that is not a copy of K , with $T^e = \{(1, 1), (2, 3), (3, 2)\}$.

Assume now that S is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$, then, since by Lemma 6 it holds that $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the path $r_i^{\sigma(e)}, b_{i,j}^e, c_j^{\sigma(e)}, a_{i',j}^{\sigma(e)}, r_{i'}^{\sigma(e)}$ together with the two copies of H^- attached to

$q^{\sigma(e)}$ and $r_i^{\sigma(e)}$ and to $q^{\sigma(e)}$ and $r_{i'}^{\sigma(e)}$ induce a graph that contains H as a topological minor. This implies that if $a_{i',j}^{\sigma(e)} \notin S$, H is a topological minor of $F_{\{H\}} \setminus S$. As this is not possible by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the lemma follows from Lemma 7. \square

We now assume that H contains at least three cut vertices. In particular, this applies to the case where $H = P_5$.

Lemma 9. *Let H be a connected graph that contains at least three cut vertices. Neither $\{H\}$ -M-DELETION nor $\{H\}$ -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework (resp. $\{H\}$ -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be precised later. By Lemma 8 and the fact that H has at least three cut vertices, we can assume that H contains at least three cut vertices that do not belong to the same block. Therefore, we can find three cut vertices a, c, b and four blocks $B_a, B_{a,c}, B_{c,b}, B_b$ such that B_a is a leaf of the block-cut tree of H and $B_a, a, B_{a,c}, c, B_{c,b}, b, B_b$ is a path in this block-cut tree. Let a' be a vertex of $V(B_a) \setminus \{a\}$ and r be a vertex of $V(B_b) \setminus \{b\}$. We define R_a to be the connected component of $H \setminus \{a', c\}$ that contains a , R_b to be the connected component of $H \setminus \{c, r\}$ that contains b , R_c to be the connected component of $H \setminus (R_a \cup R_b)$ that contains c , and R_r to be the connected component of $H \setminus R_b$ that contains r . Note that $\{a'\}, V(R_a), V(R_c), V(R_b),$ and $V(R_r)$ form a partition of $V(H)$. This decomposition of H is depicted in Figure 8.

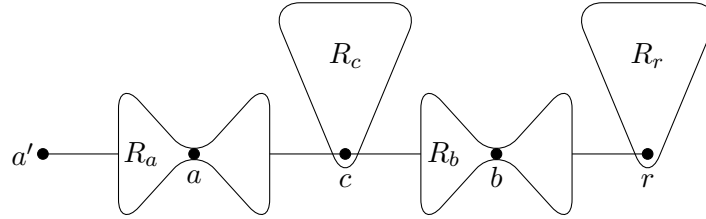


Figure 8: The decomposition of the graph H where $a, c,$ and b are three cut vertices.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , for each $e \in E(G)$ and each $i \in [1, k]$, we add a copy of R_c where we identify c and c_i^e , and a copy of R_r where we identify r and r_i^e . Moreover, for each $e \in E(G)$ and each $i, j \in [1, k]$ we add a copy of R_a where we identify a and $a_{i,j}^e$ and we connect the vertices $N_H(a') \cap R_a$ to r_i^e and the vertices $N_H(c) \cap R_a$ to c_j^e . We also add a copy of R_b where we identify b and $b_{i,j}^e$ and we connect the vertices $N_H(c) \cap R_b$ to $c_j^{\sigma(e)}$ and the vertices $N_H(r) \cap R_b$ to $r_i^{\sigma(e)}$. Note

that the vertices of the copies of R_a and R_b are B -extra vertices, and the vertices of the copies of R_c and R_r are J -extra vertices. In particular, we have $t_{\mathcal{F}} = |V(R_a) \cup V(R_b)| - 2$ and, for each $e \in E(G)$, $|V(J^e)| = |V(R_c) \cup V(R_r)| \cdot k$. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 9. As Z contains $h - 1$ vertices (both vertices a' and r of H are mapped to r_i^e), we obtain that H is not a minor of $F_{\{H\}} \setminus S_P$. Thus S_P is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) of size ℓ .

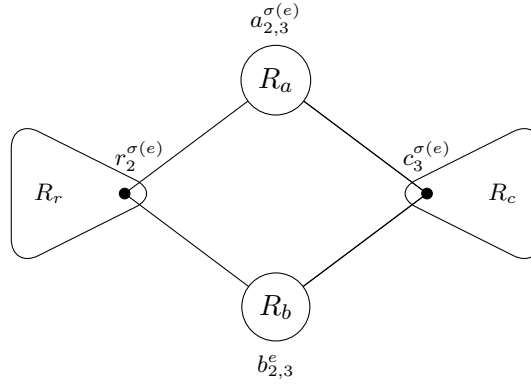


Figure 9: A connected component Z of $F_{\{H\}} \setminus S$ that contains $b_{2,3}^e$ with $(2, 3) \in T^e$.

Assume now that S is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$, then, since by Lemma 6 it holds that $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the vertex $r_{i'}^{\sigma(e)}$, the copy of R_a attached to $a_{i',j}^{\sigma(e)}$, the copy of R_c attached to $c_j^{\sigma(e)}$, the copy of R_b attached to $b_{i,j}^e$, and the copy of R_r attached to $r_i^{\sigma(e)}$ induce the graph H . This implies that if $a_{i',j}^{\sigma(e)} \notin S$, H is a subgraph of $F_{\{H\}} \setminus S$. As this is not possible by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the lemma follows from Lemma 7. \square

In the next lemma, we consider the case where H is a particular type of tree that covers the case where $H = K_{1,4}$ for the $\{H\}$ -M-DELETION problem.

Lemma 10. *Let H be a tree with at most two cut vertices and at least four vertices of degree one. $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be specified later. If H has two cut vertices x and y , then we set s_x (resp. s_y) to be the number of vertices pendent to x (resp. y). If H has only one cut vertex, we set $s_x = p - 2$ and $s_y = 2$, where p is the number of vertices of degree one.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , for each $e \in E(G)$ and each $i, j \in [1, k]$, we add $s_x - 1$ (resp. $s_y - 1$) pendent vertices to $a_{i,j}^e$ (resp. $b_{i,j}^e$). Note that the pendent vertices are B -extra vertices. In particular we have $t_{\mathcal{F}} = p - 2$ and, for each $e \in E(G)$, $|V(J^e)| = 2k$, i.e., there are no J -extra vertices. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_T$ is either a copy of the graph K , which is of size $h - 1$, or the subgraph induced by $a_{i,j}^e, r_i^e, b_{i,j}^{\sigma^{-1}(e)}, c_j^e$, and the vertices that are pendent to $a_{i,j}^e$ and $b_{i,j}^{\sigma^{-1}(e)}$, for every $(i, j) \in T$. It can be easily verified that, as by hypothesis, $s_x + s_y \geq 4$, this latter subgraph, depicted in Figure 10, does not contain H as a minor. Thus $F_{\{H\}} \setminus S_T$ does not contain H as a minor and S_P is a solution of $\{H\}$ -M-DELETION of size ℓ .

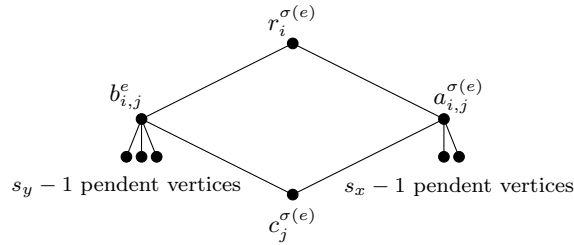


Figure 10: A connected component of $F_{\{H\}} \setminus S$ that is not a copy of K , with $s_x = 3$ and $s_y = 4$.

Assume now that S is a solution of $\{H\}$ -M-DELETION on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$, then, as by Lemma 6 it holds that $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the path $r_{i'}^{\sigma(e)}, a_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, b_{i,j}^e, r_i^e$ combined with the $s_x - 1$ vertices pendent to $a_{i',j}^{\sigma(e)}$ and the $s_y - 1$ vertices pendent to $b_{i,j}^e$ induce a graph Z that contains H as a minor. As, by definition of S , $F_{\{H\}} \setminus S$ does not contain H as a minor, we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 holds and the lemma follows from Lemma 7. \square

Observe that in the end of the above proof, if H contains two cut vertices, then Z also contains H as a topological minor, but this is not true if H is a star; this is consistent

with the single-exponential algorithms given in [6]. Therefore, we obtain the following lemma for topological minors.

Lemma 11. *Let H be a tree with exactly two cut vertices and at least four vertices of degree one. $\{H\}$ -TM-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Lemma 12. *Let H be a connected graph that contains exactly two cut vertices and each cut vertex is part of a cycle. Neither $\{H\}$ -M-DELETION nor $\{H\}$ -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework (resp. $\{H\}$ -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}} = 0$. Thanks to Lemma 8, we can assume that the block containing both cut vertices is not a cycle, hence it is an edge. Let v and v' be the two cut vertices and let H^- be the graph obtained from H by contracting the edge $\{v, v'\}$. We denote by w the new vertex.

We are now ready to describe the graph $F_{\{H\}}$. We set $t_{\mathcal{F}} = 0$. Starting from F , for each $e \in E(G)$ and each $i \in [1, k]$ we add a copy of H^- where we identify w and r_i^e . In particular, for each $e \in E(G)$, $|V(J^e)| = (|V(H^-)| + 1) \cdot k$. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 11. As Z has only one cut vertex and every block of this graph is a minor of C_4 , while $H \not\leq_m C_4$, we obtain that H is not a minor of it. Thus S_P is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) of size ℓ .

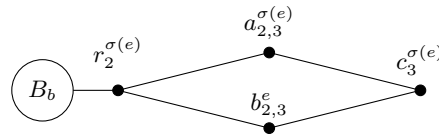


Figure 11: A connected component Z of $F_{\{H\}} \setminus S$ that is not a copy of K , that contains $b_{2,3}^e$ with $(2, 3) \in P$, where B_b means that we have attached $b \geq 2$ cycles to the vertex $r_2^{\sigma(e)}$.

Assume now that S is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$, then, since by Lemma 6 it holds that $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the path $r_i^{\sigma(e)}, b_{i,j}^e, c_j^{\sigma(e)}, a_{i',j}^{\sigma(e)}, r_{i'}^{\sigma(e)}$ together with the copies of H^- attached to $r_i^{\sigma(e)}$ and $r_{i'}^{\sigma(e)}$ induce a graph that contains H as a minor. This implies that if $a_{i',j}^{\sigma(e)} \notin S$, H

is a topological minor of $F_{\{H\}} \setminus S$. As this is not possible by definition of S , we have that $a_{i,j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the lemma follows from Lemma 7. \square

Lemma 13. *Let H be a connected graph with exactly two cut vertices such that exactly one of the two cut vertices is part of a cycle. Neither $\{H\}$ -M-DELETION nor $\{H\}$ -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework (resp. $\{H\}$ -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be defined later. By assumption, the block containing both cut vertices is an edge. Let x and y be the two cut vertices. Let C be a block that is a cycle, and without loss of generality we may assume that $x \in V(C)$. Let H_x (resp. H_y) be the connected component of $H \setminus ((V(C) \cup \{y\}) \setminus \{x\})$ (resp. $H \setminus \{x\}$) that contains x (resp. y) (see Figure 12).

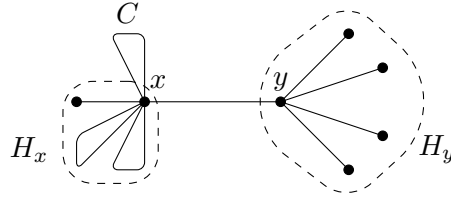


Figure 12: A visualization of C , H_x , and H_y in H .

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , we add, for each $e \in E(G)$ and each $i, j \in [1, k]$, a vertex $\bar{a}_{i,j}^e$ and the edges $\{\bar{a}_{i,j}^e, r_i^e\}$ and $\{c_j^e, \bar{a}_{i,j}^e\}$. Moreover, for each $e \in E(G)$ and each $j \in [1, k]$, we add a copy of H_x where we identify x and c_j^e , and a copy of H_y where we identify y and r_j^e . The vertices in the copies of H_x and the copies of H_y are J -extra vertices and the vertices $\bar{a}_{i,j}^e$, $e \in E(G)$ and $i, j \in [1, k]$ are B -extra vertices. In particular, we have $t_{\mathcal{F}} = 1$ and, for each $e \in E(G)$, $|V(J^e)| = (|V(H_x)| + |V(H_y)|) \cdot k$. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . The connected components of $F_{\{H\}} \setminus S_P$ are either copies of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 13. Since H_y has no cycles and H_x has one cycle less than H , it follows that if Z contains H as a minor, then there is a cycle of this minor that contains both $c_j^{\sigma(e)}$ and $r_i^{\sigma(e)}$. But in that case we cannot find in Z a block consisting of one edge whose both endpoints are cut vertices, corresponding to the edge $\{x, y\}$. We obtain that H is not a minor of the depicted graph. Thus $F_{\{H\}} \setminus S_P$ does not contain H as a minor and S_P is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) of size ℓ .

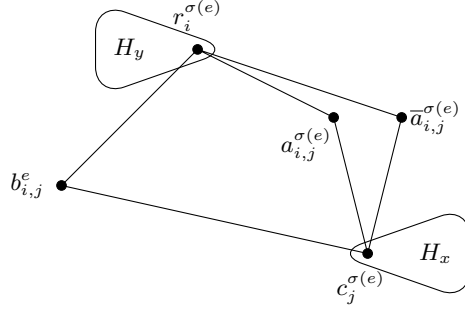


Figure 13: A connected component Z of $F_{\{H\}} \setminus S$ that is not a copy of K .

Assume now that S is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$ then, as by Lemma 6 $S \cap V(J^{\sigma(e)}) = \emptyset$, it follows that the C_4 induced by $r_{i'}^{\sigma(e)}, a_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, \bar{a}_{i',j}^{\sigma(e)}$, the path $c_j^{\sigma(e)}, b_{i,j}^e, r_i^{\sigma(e)}$, together with the copy of H_x attached to $c_j^{\sigma(e)}$ and the copy of H_y attached to $r_i^{\sigma(e)}$ induce a subgraph of $F_{\{H\}} \setminus S$ that contains H as a topological minor. This subgraph, together with an extra copy of H_y attached to $r_{i'}^{\sigma(e)}$, is depicted in Figure 14. As this is forbidden by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the lemma follows from Lemma 7. \square

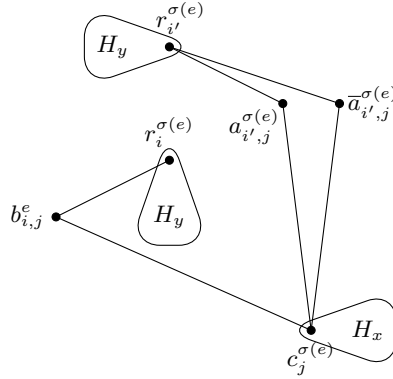


Figure 14: A connected component Z of $F_{\{H\}} \setminus S$ if $b_{i,j}^e \notin S$ and $a_{i',j}^{\sigma(e)} \notin S$ for some $e \in E(G)$, $i, i', j \in [1, k]$, $i \neq i'$.

Lemma 14. *Let H be a connected graph with exactly one cut vertex and at least two cycles. Neither $\{H\}$ -M-DELETION nor $\{H\}$ -TM-DELETION can be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework (resp. $\{H\}$ -TM-framework) of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be specified later. Let x be the cut vertex of H , and let B_1 and B_2 be two blocks that are cycles. We define the graph H_x to be $H \setminus (V(B_1 \cup B_2) \setminus \{x\})$.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , we add, for each $e \in E(G)$ and each $i, j \in [1, k]$, two new vertices $\bar{a}_{i,j}^e$ and $\bar{b}_{i,j}^e$ and the edges $\{\bar{a}_{i,j}^e, r_i^e\}$, $\{r_i^e, \bar{b}_{i,j}^e\}$, $\{\bar{b}_{i,j}^e, c_j^e\}$, and $\{c_j^e, \bar{a}_{i,j}^e\}$. Then for each $e \in E(G)$ and each $j \in [1, k]$ we add a copy of H_x where we identify x with c_j^e . The vertices in the copies of H_x are J -extra vertices and the vertices $\bar{a}_{i,j}^e$ and $\bar{b}_{i,j}^e$, $i, j \in [1, k]$ and $e \in E(G)$, are B -extra vertices. In particular we have $t_{\mathcal{F}} = 2$ and, for each $e \in E(G)$, $|V(J^e)| = (|V(H_x)| + 1) \cdot k$. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 15. Since Z has only one block more than H_x , it follows that Z does not contain H as a minor. Thus $F \setminus S_T$ does not contain H as a minor, and S_P is a solution of $\{H\}$ -M-DELETION (resp. $\{H\}$ -TM-DELETION) of size ℓ .

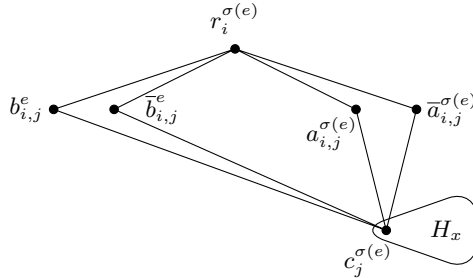


Figure 15: A connected component Z of $F_{\{H\}} \setminus S$ that is not a copy of K .

Assume now that S is a solution of $\{H\}$ -M-DELETION on $F_{\{H\}}$ of size ℓ . Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$ then, as by Lemma 6 $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the graph induced by the two paths $r_{i'}^{\sigma(e)}, a_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, b_{i,j}^e, r_i^{\sigma(e)}$ and $r_{i'}^{\sigma(e)}, \bar{a}_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, \bar{b}_{i,j}^e, r_i^{\sigma(e)}$, together with the copy of H_x attached to $c_j^{\sigma(e)}$, depicted in Figure 16, is a subgraph of $F \setminus S$ containing H as a topological minor. As this is forbidden by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 (resp. Property 2) holds and the lemma follows from Lemma 7. \square

In the next two lemmas, namely Lemma 15 and Lemma 16, we deal separately with the minor and topological minor versions, respectively.

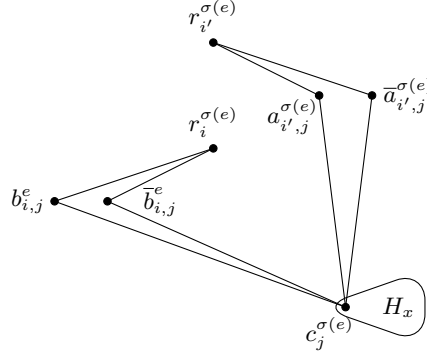


Figure 16: A connected component Z of $F_{\{H\}} \setminus S$, if $b_{i,j}^e \notin S$ and $a_{i',j}^{\sigma(e)} \notin S$ for some $e \in E(G)$, $i, i', j \in [1, k]$, $i \neq i'$.

Lemma 15. *Let H be a connected graph with exactly one cut vertex and exactly one cycle such that H is not a minor of the banner. $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -M-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be specified later. Let s be the number of vertices of degree one in H . As H is not a minor of the banner and (because of Theorem 3) we can assume that each block of H contains at most four edges, we have that $s \geq 2$.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , we add, for each $e \in E(G)$ and each $j \in [1, k]$, three new vertices d_j^e , f_j^e , and g_j^e and the edges $\{d_j^e, f_j^e\}$ and $\{f_j^e, g_j^e\}$. Moreover, for each $e \in E(G)$ and each $i, j \in [1, k]$, we add the edges $\{b_{i,j}^{\sigma^{-1}(e)}, d_j^e\}$ and $\{g_j^e, a_{i,j}^{\sigma(e)}\}$, and $s - 2$ vertices pendent to $b_{i,j}^e$. The vertices d_j^e , f_j^e , and g_j^e , $j \in [1, k]$ and $e \in E(G)$, are J -extra vertices, and the pendent vertices are B -extra vertices. In particular we have $t_{\mathcal{F}} = s - 2$ and, for each $e \in E(G)$, $|V(J^e)| = 5k$. This completes the definition of $F_{\{H\}}$.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 17. Note that Z contains three different cycles, but for each of them some pendent edge is missing in order to find H as a minor. Thus $F \setminus S_P$ does not contain H as a minor, and S_P is a solution of $\{H\}$ -M-DELETION.

Assume now that S is a solution of $\{H\}$ -M-DELETION. Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$ then, as by Lemma 6 $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the graph Z induced by the two paths $r_{i'}^{\sigma(e)}, a_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, b_{i,j}^e, r_i^{\sigma(e)}$ and $a_{i',j}^{\sigma(e)}, g_j^{\sigma(e)}, f_j^{\sigma(e)}, d_j^{\sigma(e)}, b_{i,j}^e$, and the $s - 2$ vertices pendent to $b_{i,j}^e$, depicted in Figure 18, is a subgraph of $F_{\{H\}} \setminus S$ containing H as a minor. As this

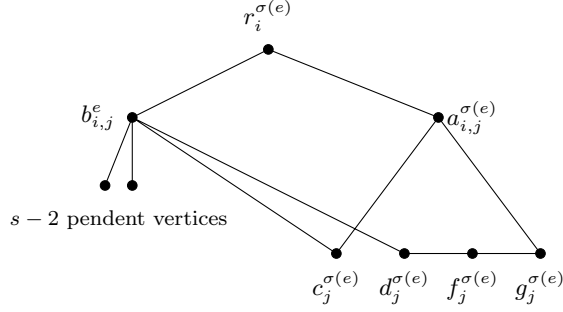


Figure 17: A connected component Z of $F_{\{H\}} \setminus S$ that is not a copy of K with $s = 4$.

is forbidden by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 1 holds and the lemma follows from Lemma 7. \square

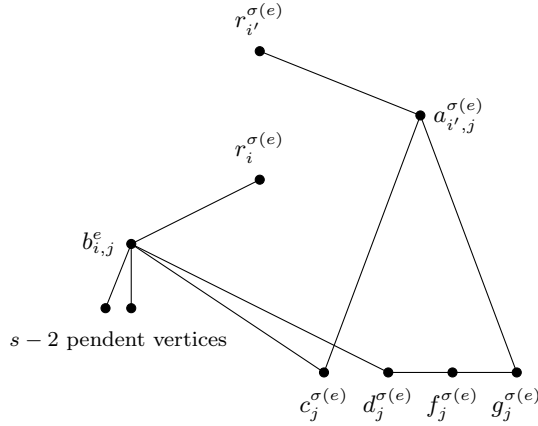


Figure 18: A connected component Z of $F_{\{H\}} \setminus S$, if $b_{i,j}^e \notin S$ and $a_{i',j}^{\sigma(e)} \notin S$ for some $e \in E(G)$, $i, i', j \in [1, k]$, $i \neq i'$ where $s = 4$.

Lemma 16. *Let H be a connected graph with exactly one cut vertex and exactly one cycle such that H is not a minor of the banner. $\{H\}$ -TM-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails.*

Proof. Let (F, ℓ) be the $\{H\}$ -TM-framework of an input (G, k) of $k \times k$ PERMUTATION INDEPENDENT SET, where $t_{\mathcal{F}}$ will be specified later. Let s be the number of vertices of degree one in H . As in Lemma 15, since H is not a minor of the banner and (because of Theorem 3) we can assume that each block of H contains at most four edges, we have that $s \geq 2$.

We are now ready to describe the graph $F_{\{H\}}$. Starting from F , we add, for each $e \in E(G)$, a vertex q^e and s vertices pendent to q^e , and for each $e \in E(G)$ and each $j \in [1, k]$, the edge $\{q^e, r_j^e\}$. The vertices q^e , $e \in E(G)$, and the pendent vertices are J -extra vertices. In particular we have $t_{\mathcal{F}} = 0$ and, for each $e \in E(G)$, $|V(J^e)| = 2k + s + 1$. This completes the definition of $F_{\{H\}}$. Note that this construction is similar to the construction provided in Section 4.2 with H_x being a star with s leaves, x the non-leaf vertex, and $H_{\bar{y}}$ an edge.

Let P be a solution of $k \times k$ PERMUTATION INDEPENDENT SET on (G, k) . Then every connected component of $F_{\{H\}} \setminus S_P$ is either a copy of the graph K , which is of size $h - 1$, or the graph Z depicted in Figure 19. Note that each vertex of Z contained in a cycle is of degree at most three. Since $s \geq 2$, there is a vertex in H of degree at least four contained in a cycle. Thus $F \setminus S_P$ does not contain H as a topological minor, and S_P is a solution of $\{H\}$ -TM-DELETION.

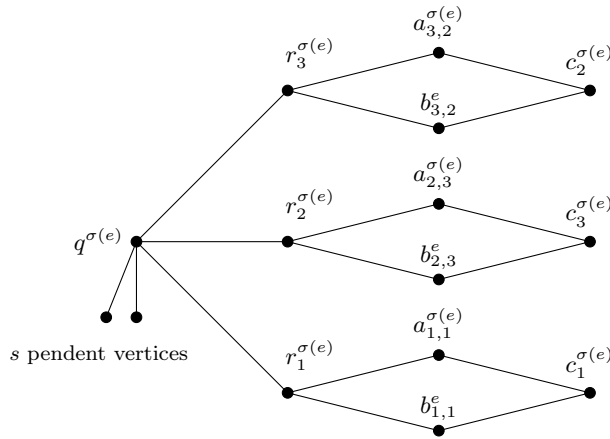


Figure 19: A connected component Z of $F_{\{H\}} \setminus S_P$ that is not a copy of K , with $T^e = \{(1, 1), (2, 3), (3, 2)\}$ and $s = 2$.

Assume now that S is a solution of $\{H\}$ -TM-DELETION. Let $e \in E(G)$ and let $i, j \in [1, k]$ such that $b_{i,j}^e \notin S$. Let $i' \in [1, k]$ such that $i \neq i'$. If $a_{i',j}^{\sigma(e)} \notin S$ then, as by Lemma 6 $S \cap V(J^{\sigma(e)}) = \emptyset$, we have that the graph Z induced, by the cycle $q^{\sigma(e)}, r_{i'}^{\sigma(e)}, a_{i',j}^{\sigma(e)}, c_j^{\sigma(e)}, b_{i,j}^e, r_i^{\sigma(e)}, q^{\sigma(e)}$ and the s vertices pendent to $q^{\sigma(e)}$, is a subgraph of $F_{\{H\}} \setminus S$ containing H as a topological minor. This situation is depicted in Figure 20. As this is forbidden by definition of S , we have that $a_{i',j}^{\sigma(e)} \in S$. Thus Property 2 holds and the lemma follows from Lemma 7. \square

We are now ready to prove Theorem 4 and Theorem 5.

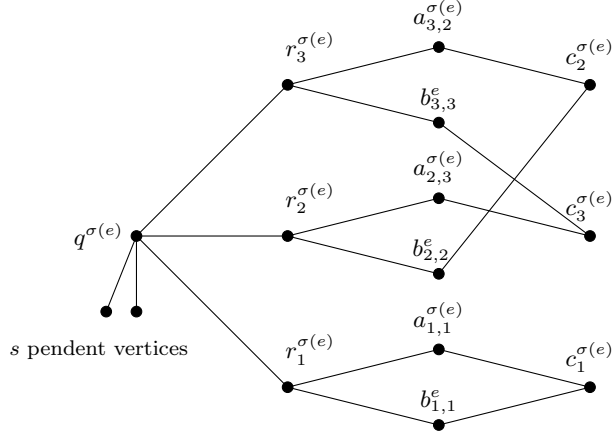


Figure 20: A connected component Z of $F_{\{H\}} \setminus S_P$ that is not a copy of K , with $T_e = \{(1, 1), (2, 2), (3, 3)\}$, $T_{\sigma(e)} = \{(1, 1), (2, 3), (3, 2)\}$, and $s = 2$.

Proof of Theorem 4 and Theorem 5. Let H be a graph in \mathcal{Q} . If H is a star with at least four leaves, then by Lemma 10 $\{H\}$ -M-DELETION cannot be solved in time $2^{o(\text{tw} \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ unless the ETH fails. In the following we assume that H is not a star. This permits us to proceed with the proofs of both theorems in a unified way.

If H contains at least one block with at least five vertices, then such a block would have at least five edges as well (by definition of a block), hence by Theorem 3, the theorems hold. We can now assume that H does not contain any block with at least five vertices. Therefore, every block of H is either an edge, a C_3 , or a C_4 .

If H contains at least three cut vertices that do not belong to the same block, then Lemma 9 can be applied. We can now assume that H contains at most two cut vertices.

Assume now that H contains exactly two cut vertices and let B the block that contains both of them. If B is not an edge, then Lemma 8 can be applied. Otherwise, we distinguish cases depending on the shape of the two connected components of H after removing the only edge of B . If both connected components contain a cycle then Lemma 12 can be applied, if only one of them contains a cycle then Lemma 13 can be applied, and if none of them contains a cycle then, as H is not a minor of the banner, Lemma 10 can be applied.

Assume now that H contains exactly one cut vertex. As H is not a star, then either H contains at least two cycles, and so Lemma 14 can be applied, or H contains exactly one cycle (and is not a minor of the banner) and therefore Lemma 15 or Lemma 16 can be applied. The theorems follow. \square

5 Conclusions and further research

We provided lower bounds for \mathcal{F} -M-DELETION and \mathcal{F} -TM-DELETION parameterized by the treewidth of the input graph, several of them being tight. In particular, the results of this article together with those of [5, 6] settle completely the complexity of $\{H\}$ -M-DELETION when H is planar and connected.

Concerning the topological minor version, in order to establish a dichotomy for $\{H\}$ -TM-DELETION when H is planar and connected, it remains to obtain algorithms in time $2^{\mathcal{O}(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ for the graphs H with maximum degree at least four, like the **gem** or the **dart** (see Figure 1), as for those graphs the algorithm in time $2^{\mathcal{O}(\text{tw} \cdot \log \text{tw})} \cdot n^{\mathcal{O}(1)}$ given in [5] cannot be applied.

It is easy to check that the lower bounds presented in this article also hold for *treedepth* (as it is the case in [25]) which is a parameter more restrictive than treewidth [12]. Also, it is worth mentioning that \mathcal{F} -M-DELETION and \mathcal{F} -TM-DELETION are unlikely to admit polynomial kernels parameterized by treewidth for essentially any collection \mathcal{F} , by using the framework introduced by Bodlaender et al. [9] (see [12] for an explicit proof for any problem satisfying a generic condition).

Finally, let us mention that Bonnet et al. [11] recently studied generalized feedback vertex set problems parameterized by treewidth, and showed that excluding C_4 plays a fundamental role in the existence of single-exponential algorithms. This is related to our dichotomy for cycles illustrated in Figure 1 (which we proved independently in [3] building on the work of Pilipczuk [25]), namely that $\{C_i\}$ -DELETION can be solved in single-exponential time if and only if $i \leq 4$.

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